

RFID A1 Module User Manual

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1 Introduction

1.1 Device Overview

Features

- Low cost RFID Reader with MIFARE Classic®, MIFARE Ultralight® and NTAG2xx® support
- Multiple Tag support
- Polling functionality (Stand-alone mode)
- Command interface via I2C and SPI
- I2C speed up to 400 kHz
- SPI speed up to 500 kHz
- Compact form factor
- Castellated SMT pads for easy and reliable PCB mounting
- High transponder read and write speed
- Low power design
- Single operating voltage: 2.5V to 3.6V
- -25°C to 85°C operating range
- AES-128 encryption engine
- RoHS compliant

Applications

- Access control
- Monitoring goods
- Approval and monitoring consumables
- Pre-payment systems
- Managing resources
- Connection-less data storage systems

Description

The RFID A1 module is the first in an evolving family of 13.56MHz sub assemblies from Eccel Technology Ltd (IB Technology). The product is designed with embedded applications in mind. This product is an ideal design choice if the user wishes to add RFID capability to their design quickly and without requiring extensive RFID and embedded software expertise and time. An on board low power ARM microcontroller handles the RFID configuration setup and provides the user with a powerful yet simple command interface to facilitate fast and easy read/write access to the memory and features of the various transponders supported by this module.

A polling option with configurable TPI outputs that are dependent upon the UID of the tag detected, makes it simple to build a standalone device that can control door locks, etc. and additionally send the signal to the host controller.

The module simply requires a single power and GND connection from the user PCB, along with two connections to an antenna. Eccel Technology Ltd (IB Technology) provide a range of suitable antennas designed for use with this module.

1.2 Pinout

Table 1-1

1.3 Application

The RFID A1 module is specifically designed for embedded applications, where a fast or low pin count connection to a host microcontroller is required. Fastest data exchange speed can be obtained by using the Serial Peripheral Interface Bus. The lowest pin count connection can be achieved by using the I2C Bus. For the module to be a fully functional RFID reader/writer it only requires connection to a power supply and antenna. Eccel Technology Ltd manufactures a wide range of such antennas designed to give optimal performance with this module. Please contact us for further details of our range of antennas.

Figure 1.4-1 Typical application with SPI interface

2 Electrical Characteristics

2.1 Test Conditions

Typical device parameters have been measured at ambient temperature 22°C ±3°C and a power supply of 3.3V ±5%.

2.2 Absolute Maximum Ratings

Table 2-1

2.3 Operating Conditions

Table 2-2

2.4 Current Consumption

Table 2-3

Current measurement was taken place with all digital input pins connected to Vdd. Pins number 16 and 17 was left floating. Antenna was not connected except for the Power Up Current measurement and Maximum Current measurement.

2.5 GPIO

Table 2-4

2.6 Antenna Output

Table 2-5

2.7 Communication Buses

Table 2-6

2.8 Flash

Table 2-7

3 System

3.1 Overview

The general overview of system components is shown in [Figure 3.1-1.](#page-10-0) The system internally consists of four main parts:

- CORE the main processing part of the microcontroller firmware responsible for managing all system tasks, parsing and the execution of commands received from the user's master controller.
- RFID dedicated RFID IC together with its firmware drivers responsible for communication with the RFID tag.
- COMMUNICATION SYSTEM microcontroller's peripherals (I2C and SPI) and its firmware drivers providing communication functionality to the user.
- POWER MANAGER a subsystem responsible for managing power states and clocks in the module to minimise power consumption during operation.

Figure 3.1-1

3.2 Modules

3.2.1 Core

The core of the system shown above in figure 3.1, is the part of the modules' hardware and firmware responsible for managing tasks, prioritizing them and allocating resources to other components of the system. When the user sends a command, the core is the part of the system which ensures that the command is executed properly, that the outcome is saved and that a response is sent back to the user's master controller.

3.2.2 RFID

The RFID component is the part of the hardware and module firmware responsible for the complete communication interaction with the transponder via RF. In the RFID A1 module this part currently provides communication capabilities with MIFARE Classic®, MIFARE Ultralight® and NTAG2xx® transponder types. The key features of this RFID section are both fastest speed of transmission and system responsiveness, and optimal low power consumption.

3.2.3 Communication System

The Communication System component of the RFID A1 module is responsible for the entire communication with the user master controller, together with parsing commands for subsequent execution by the module core. Two communication interfaces are available: SPI and I2C. Both have equivalent access to the Communication System. It is possible to communicate with the module using all two interfaces at the same time, but it is strongly recommended to select and use only one at a time.

Communication peripherals are always active unless the nPWRDN signal is asserted. Doing so will put the module into Deep Sleep Mode and turn off entire Communication System.

3.2.4 Power Manager

The Power Manager is the part of the firmware, together with hardware support, that manages the power states of the system and tries to minimize power consumption as much as possible. The user has the option to put the module into Power Down Mode by pulling low the nPWRDN pin.

When the system enters Power Down Mode, all clocks and submodules are disabled and no response to user communication commands will occur until the system is re-enabled by the nPWDN being taken high by the user's system.

3.3 Memory Map

The device memory layout is shown below in Table 3.1.

In the RFID A1 module there are 728 bytes of user accessible memory. Each byte of the memory has a defined factory default value and these values can be recovered by using the 'Reset to Factory Defaults' command. The first 288 bytes are volatile memory which also have a default reset state that is the same as the factory default value. The other memory is buffered non-volatile memory and can be modified and stored using 'Unlock' and 'Lock' commands. The RFID A1 module is automatically returned to factory default state if after power up there is no valid configuration stored in non-volatile memory. The first 288 bytes of this user accessible memory are not stored in non-volatile memory and are always reset to factory defaults after a power-up sequence or after exit from the Power Down Mode.

Reading and writing to the registers can be done via one of two interfaces available on the module – I2C or SPI. The rest of the memory is accessible only when the module is unlocked. The default state of the module after power up is locked.

3.3.1 Result Register

The Result Register is 1-byte long, with located at address 0x0000, with both read and write access. Writing to this register has no effect. The register contains the result (error code) of the last executed command. The list of all possible results is shown in [Table 3-2.](#page-14-1)

Table 3-2

3.3.2 Command Register

The Command Register is 1-byte, located at address 0x0001, with both read and write access. Writing to this register is recognized by the module as a command execution request. Depending upon the command (value written to the register) the Command Parameters Register is parsed to extract arguments for the command. Whilst commands are executing, the Result Register value is set to 0xFF and the nBUSY line is driven low. When command execution is complete the memory is updated and the nBUSY line returns to the high state indicating that the module is ready to receive another command.

3.3.3 Command Parameters Register

Table 3-3

The Command Parameters Register is 18-bytes long, located at addresses from 0x0002 to 0x0013, with both read and write access. This is the place from where the system parses the arguments necessary to perform requested operation when a command is executed. Depending upon the command, this register is parsed and interpreted in different ways. The details of interpretation of the data stored in this register can be found in chapter [4.2.](#page-20-2) The module never changes the values inside this register except after power-up or exit from Deep Sleep Mode.

3.3.4 Tag UID Register

Table 3-4

Tag UID Register is 10-bytes long, located at addresses from 0x0014 to 0x001D, with read only access. In most cases the tag UIDs are 4 bytes, 7 bytes or 10 bytes long, thus this register can cover the longest UID but not necessarily all bytes available in the register will be used. Bytes within the register are ordered from the least significant byte to the

most significant byte. This register is updated whenever Get UID and Type commands are used. The Tag UID Size Register contains information detailing how many bytes of the ten available represent the tag UID.

3.3.5 Tag Type Register

The Tag Type Register is 1-byte long located at address 0x001E, with read only access. This register contains information about the type of tag which was last seen in the field. Possible tag types are shown in [Table 3-5.](#page-16-3)

Table 3-5

3.3.6 Tag UID Size Register

The Tag UID Size Register is 1-byte long, located at address 0x001F, with read only access. It contains the information of what the UID size in bytes was of the last tag in the field.

3.3.7 Data Buffer

The Data Buffer is a 256-byte long, located at address 0x0020 to 0x011F, with both read and write access. This buffer is used for data transfers between the tag and the user of the module.

3.3.8 Password Register

Table 3-6

The Password Register is 8-bytes long, located at address 0x0120 to 0x0127, with both read and write access but only after first unlocking the device. This register is inaccessible when the module is locked. It contains an 8-byte long password which must be used with the Unlock Command to unlock protected memory. This password can be changed when the device is unlocked. The password change will only be updated and become valid after executing the Lock command.

3.3.9 AES Initialization Vector Register

Table 3-7

The AES Initialization Vector Registers are two 16-bytes long registers, located at address 0x0128 to 0x0147, with both read and write access but only after first unlocking the device. These registers are inaccessible when the module is locked. They can be used as an initialization vector for the first encrypted data block. The byte order in the memory is from the least significant byte. Their role during data encryption and decryption is described in detail in chapte[r 4.2.8](#page-29-0) and [4.2.9.](#page-31-0)

3.3.10 AES Key Register

Table 3-8

The AES Key Registers are two 16-bytes long registers, located at address 0x0148 to 0x0167, with both read and write access but only after first unlocking the device. These registers are inaccessible when the module is locked. Both registers contain an AES encryption key which can be used for encryption of the Data Buffer. Their role during encryption and decryption of the data in the buffer is described in detail in chapter [4.2.8](#page-29-0) and [4.2.9.](#page-31-0)

3.3.11 Authentication Key / Password Register

The Authentication Key and Password Registers are forty 6-bytes long registers, located at address 0x0168 to 0x0257, with both read and write access but only after first unlocking the device. These registers are inaccessible when the module is locked. When working with MIFARE Classic® tags these registers contain the password keys used for block authentication in the tag. When working with Ultralight and NTAG transponders, these registers contain 4-byte passwords. There are forty Authentication Key and Password registers numbered from 0 to 39. The number of the key register to be used is passed as an argument in some commands.

3.3.12 User Memory

There are 128 bytes of memory available for the user as a protected memory space from address 0x0258 to 0x02D7. This memory is inaccessible when the device is locked. This data is stored into non-volatile memory when the Lock command is executed.

4 Functional Description

4.1 Overview

The RFID A1 module is designed to provide a simple interface to communicate and manage RFID tags from the MIFARE Classic®, MIFARE Ultralight® and NTAG2xx® families. The interface is provided via SPI and I2C busses. From the user perspective, the module is a memory region where he can read from and write to. Depending upon the memory address, the role of registers located in the memory differ as per the Memory Map description given in Chapter [3.3.](#page-12-0) By writing to and reading from these registers, the user can achieve the desired result.

4.2 Commands

The RFID A1 Module has a special register, the Command Register. Writing to this register is interpreted as a command execution request from the user. The command type corresponds to the value written into the register. Execution of a command starts when the communication is finished. Depending upon the bus used, definition of the communication endpoint varies. When using the I2C bus, communication is finished when a STOP condition is sent. In case of the SPI bus, communication is finished on the rising edge of the CS signal.

The nBUSY line goes from the high state to the low state and indicates that the module is busy (it parses data stored in the registers and executes commands, if any command was sent). A more detailed description of the signal behavior is given in chapter [5.1.](#page-60-1)

The full RFID A1 command list is shown below i[n Table 4-1.](#page-21-0)

Table 4-1

Depending upon the command, different arguments are taken as parameters and various registers are modified. After each command execution, the Result Register is updated.

4.2.1 Get UID and Type (0x01)

The 'Get UID and Type' command takes no arguments. After receiving this command, the RFID A1 Module checks for any tag presence in the field. If there is no tag in the field, it returns 'No Tag' value in the 'Tag Type Register'. If there is a tag in the field, it reads its UID and type and writes it to Tag UID and Tag Type registers. The 'UID Size Register' together with the 'Result Register' are updated also. This command must always be executed first before any other command to turn on and initialize the tag. Additionally, this command must be executed after any error when doing any operations on a tag to reset the tag and to enable it to be able to respond to further command requests. Additionally, via this command one can discover serial numbers from MIFARE Plus®, MIFARE® DESfire® and any transponder compatible with the ISO14443A standard.

4.2.2 Read Block (0x02)

Table 4-2

The Read Block command takes as arguments block number of the first block to read (Absolute Block Number), the number of blocks to read (Read Length), the byte offset in the Data Buffer (Data Buffer Offset), the Authentication Key Number and (optionally) the Authentication Key. If the Authentication Key is set-up, the command first tries to authenticate the tag memory sector which contains the first block. If authentication fails, the 'Result Register' is updated with the appropriate error value and command execution finishes. After a successful authentication, the first block is read and copied to the Data Buffer. The command continues to read and copy following blocks if the Read Length is higher than 1. If the next blocks fall into another tag memory sector, this sector is authenticated using the same key. If everything goes well all block bytes from all blocks are copied to the Data Buffer starting from Data Buffer Offset, and the Result Register is updated with the 'No Error' value. If the RFID Module is unable to read tag memory, an error value is stored in the Result Register.

4.2.3 Write Block (0x03)

Table 4-3

The Write Block command takes as arguments block number of the first block to write (Absolute Block Number), the number of blocks to write (Write Length), the byte offset in the Data Buffer (Data Buffer Offset), the Authentication Key Number and (optionally) the Authentication Key. If the Authentication Key is set-up, the command first tries to authenticate the tag memory sector which contains the first block. If authentication fails, the Result Register is updated with the appropriate error code and command execution finishes. After a successful authentication, the data is copied from the Data Buffer to the first block. The command continues to copy data from the Data Buffer to the following blocks if the Write Length is higher than 1. If subsequent blocks fall into another tag memory sector, this sector is authenticated using the same key. If everything goes well all block bytes are copied from the Data Buffer starting from Data Buffer Offset, and the Result Register is updated with the 'No Error' code. If the RFID Module is unable to write to the blocks an appropriate error code is stored in the Result Register.

4.2.4 Read Data Block (0x04)

Table 4-4

The Read Data Block command takes as arguments block number of the first block to read (Absolute Block Number), the number of blocks to read (Read Length), the byte offset in the Data Buffer (Data Buffer Offset), the Authentication Key Number and (optionally) the Authentication Key. If the Authentication Key is set-up, the command first tries to authenticate the tag memory sector which contains the first block. If authentication fails, the Result Register is updated with the appropriate error value and command execution finishes. After successful authentication, the first block is read and copied to the Data Buffer. The command continues to read and copy following blocks if the Read Length is higher than 1. The difference between the Read Data Block command and the Read Block command is that the first one omits blocks containing authentication keys and lock bits. It only reads the data blocks from tag memory. If the next blocks fall into another tag memory sector, this sector is authenticated using the same key. If everything goes well all block bytes from all blocks are copied to the Data Buffer starting from the Data Buffer Offset, and the Result Register is updated with the 'No Error' value. If the RFID Module is unable to read tag memory, an error value is stored in the Result Register.

4.2.5 Write Data Block (0x05)

Table 4-5

The Write Data Block command takes as arguments block number of the first block to write (Block Address), the number of blocks to write (Write Length), the byte offset in the Data Buffer (Data Buffer Offset), the Authentication Key Number and (optionally) the Authentication Key. If the Authentication Key is set-up, the command first tries to authenticate the tag memory sector which contains the first block. If authentication fails, the 'Result Register' is updated with the appropriate error value and command execution finishes. After a successful authentication, the data is copied from the Data Buffer to the first block. The command continues to copy data from the Data Buffer to the following blocks if the Write Length is higher than 1. The difference between Write Data Block command and Write Block command is that the first one omits blocks containing authentication keys and lock bits. It only writes to the data blocks in tag memory. If subsequent blocks fall into another tag memory sector, this sector is authenticated using the same key. If everything goes well all block bytes are copied from the Data Buffer starting from the Data Buffer Offset, and the Result Register is updated with a 'No Error' value. If the RFID Module is unable to write to the blocks, an appropriate error code is stored in the Result Register..

4.2.6 Read Page (0x06)

Table 4-6

The Read Page command takes as arguments page number of the first page to read (Absolute Page Number), the number of pages to read (Read Length) and the byte offset in the Data Buffer (Data Buffer Offset). If any of the pages to be read are password protected a Password Authentication is necessary before doing a read operation. If read fails, the Result Register is updated with an error value and command execution finishes. After successful communication, the first page is read and copied to the Data Buffer. Command continues to read and copy subsequent pages if the Read Length is higher than 1. If everything goes well all bytes from all pages are copied to the Data Buffer starting from Data Buffer Offset, and the Result Register is updated with a 'No Error' value.

4.2.7 Write Page (0x07)

Table 4-7

The Write Page command takes as arguments page number of the first page to be written (Absolute Page Number), the number of pages to be written (Write Length) and the byte offset in the Data Buffer (Data Buffer Offset). If any of the pages to be written are password protected, then a Password Authentication is necessary before doing a write operation. If write fails, the Result Register is updated with an error value and the command execution finishes. After successful communication, the data is copied from the Data Buffer to the first page. This Command continues to copy to subsequent pages if the Write Length is higher than 1. If everything goes well all bytes are copied from the Data Buffer starting from the Data Buffer Offset, and the Result Register is updated with a 'No Error' value.

4.2.8 Encrypt Data (0x08)

Table 4-8

The Encrypt Data command takes as command arguments the Encryption Key Number (0x00 or 0x01), the Initialization Vector Number (0x00 or 0x01), the Data Buffer Offset (16-bytes blocks) and the Data Length (16-bytes blocks). This command encrypts the 'Data Length' number of 16-byte blocks using the AES128 CBC encryption methodology. It operates only within the Data Buffer. This encryption method is shown in [Figure 4.2-1.](#page-30-0)

Cipher Block Chaining (CBC) mode encryption

Figure 4.2-1

The Initialization Vector is used at the beginning to encrypt the first plaintext. For each subsequent 16-byte block instead of the Initialization Vector, the algorithm uses the ciphertext output of the previous 16-byte block encryption. The same AES Key is used for all blocks. When working with the Data Buffer during encryption the plaintext is substituted with the corresponding ciphertext.

4.2.9 Decrypt Data (0x09)

Table 4-9

The Decrypt Data command takes as arguments the Decryption Key Number (0x00 or 0x01), the Initialization Vector Number (0x00 or 0x01), the Data Buffer Offset (16-bytes) and the Data Length (16-bytes). This command decrypts the 'Data Length' number of 16-bytes blocks using the AES128 CBC decryption methodology. It operates only within the Data Buffer. This decryption method is shown in [Figure 4.2-2.](#page-32-0)

Cipher Block Chaining (CBC) mode decryption

Figure 4.2-2

The Initialization Vector is used at the beginning to decrypt the first ciphertext. For each subsequent 16-byte block, instead of the Initialization Vector, the algorithm uses the ciphertext output of the previous 16-byte block decryption. The same AES Key is used for all blocks. When working with the Data Buffer during decryption the ciphertext is substituted with the corresponding plaintext.

4.2.10 Read Value (0x0A)

Table 4-10

The Read Value command takes as arguments block number of the first block to read (Absolute Block Number), the offset of the Data Buffer where the value and read address will be stored (Data Buffer Offset), the Authentication Key Number and (optionally) the Authentication Key. This command has the same functionality as the Read Block command except it treats the block as a special Value type defined by the MIFARE® standard. It tries to parse the block content to this 4-byte signed value and to read a-byte address stored in last four bytes of the block. If read was successful, the value is stored in the Data Buffer at the pointed offset, with the address byte stored straight after the value. The value is stored as a signed 32-bit integer, with least significant byte first, and the address is stored as an unsigned 8-bit value.

4.2.11 Write Value (0x0B)

Table 4-11

The Write Value command takes as arguments block number of the first block to write (Absolute Block Number), the signed 32-bit integer value to be stored (Value), the block address value to store in the memory block (Stored Block Address), the Authentication Key Number and (optionally) the Authentication Key. This command has the same functionality as the Write Block command, except that it treats the block as a special Value type as defined by the MIFARE® standard. It tries to format the block content to this 4-byte signed value together with the address in the last 4 bytes of the memory.

4.2.12 Increment Value (0x0C)

Table 4-12

The Increment Value command is an operation on a value-type block as defined by the MIFARE® standard. It takes as arguments block number of the block where the value is stored (Absolute Block Number), the signed 32-bit integer number which will be added to the value (Delta Value), the Authentication Key Number and (optionally) the Authentication Key. The command reads the value from the block to the volatile memory on the tag and increments it by the selected Delta Value. There is no further operation done. To store the value in the same or another block (copy to non-volatile memory on the tag), the user must execute the Transfer Value command.

4.2.13 Decrement Value (0x0D)

Table 4-13

The Decrement Value command is an operation on a value-type block, as defined by the MIFARE® standard. It takes as arguments block number of the block where the value is stored (Absolute Block Number), the signed 32-bit integer number which will be subtracted from the value (Delta Value), the Authentication Key Number and (optionally) the Authentication Key. The command reads the value from the block into the volatile memory on the tag and decrements it by the delta value. There is no further operation done. To store the value in the same or another block (copy to nonvolatile memory on the tag), the user must execute the Transfer Value command.

4.2.14 Restore Value (0x0E)

Table 4-14

The Restore Value command is an operation on a value-type block, as defined by the MIFARE® standard. It takes as arguments block number of the block where the value is stored (Absolute Block Number), the Authentication Key Number and (optionally) the Authentication Key. If the block is properly formatted, then the value from the block is copied to a volatile memory register on the tag. There is no further operation done. To store the value in the same or another block (copy to non-volatile memory on the tag), the user must execute the Transfer Value command.

4.2.15 Transfer Value (0x0F)

Table 4-15

The Transfer Value command is an operation on a value-type block as defined by the MIFARE® standard. It takes as arguments block number of the block where the value is stored (Absolute Block Number), the Authentication Key Number and (optionally) the Authentication Key. The value from the volatile register on the tag is copied to the pointed block. The block is formatted as a value-type block during this operation.

4.2.16 Recover Value (0x10)

Table 4-16

The Recover Value command takes as arguments block number of the block where the value is stored (Absolute Block Number), the offset of the Data Buffer where the value and read address will be stored (Data Buffer Offset), the Authentication Key Number and (optionally) the Authentication Key. This command has the same functionality as the Read Value command, except that it can be used on a block which is corrupted – it tries to recover data from a corrupted block. The format of a value-type block allows for some bits to be corrupted and it still be possible to read and recover the proper value. The 4-byte signed value followed by 1-byte address value is stored in the Data Buffer under the Data Buffer Offset index.

4.2.17 Get Version (0x11)

The Get Version command doesn't take any arguments. This command can be used with MIFARE Ultralight® EV1, NTAG213®, NTAG215® and NTAG216® tags. After successful reading, the first 8 bytes in the data buffer are filled with data defined by the NXP standard. Example are shown i[n Table 4-17 \(](#page-40-1)MIFARE Ultralight® EV1) an[d Table 4-18](#page-40-2) (NTAG®).

| Byte No | Description | MF0UL11/MF0ULH11 | MF0UL21/MF0ULH21 | Interpretation | |
|----------------|---------------------|------------------|------------------|---------------------------|--|
| Ω | Fixed header | 0x00 | 0x00 | | |
| | Vendor ID | 0x04 | 0x04 | NXP Semiconductors | |
| $\overline{2}$ | Product type | 0x03 0x03 | | MIFARE Ultralight | |
| 3 | Product subtype | 0x01/0x02 | 0x01/0x02 | | |
| 4 | Major product | 0x01 | 0x01 | EV1 | |
| | version | | | | |
| 5 | Minor product | 0x00 | 0x00 | V ₀ | |
| | version | | | | |
| 6 | Storage size | 0x0B | 0x0E | | |
| | Protocol type | 0x03 | 0x03 | ISO/IEC 14443-3 | |
| | | | | compliant | |

Table 4-17

Table 4-18

Please refer to the NXP documentation for more information.

4.2.18 Read Signature (0x12)

The Read Signature command doesn't take any arguments. This command can be used with MIFARE Ultralight® EV1, NTAG213®, NTAG215® and NTAG216® tags. After successful reading, the first 32 bytes in the data buffer are filled with the ECC signature defined by the NXP standard. Please refer to the NXP documentation for more information.

4.2.19 Configure UID (0x13)

Table 4-19

The Configure UID command takes as arguments the UID configuration byte (UID Type), the Authentication Key Number and (optionally) the Authentication Key. This command changes the configuration of the UID on some MIFARE Classic® tags. The UID Type parameter is explained in [Table 4-20.](#page-42-0)

Table 4-20

4.2.20 Read Counter (0x14)

Table 4-21

The Read Counter command takes as arguments the tag Counter Number and the Data Buffer Offset in bytes. This command reads the counter value of the counter pointed to by the Counter Number and stores it in the Data Buffer at the Data Buffer Offset index as an unsigned 24-bit integer with the least significant byte first.

4.2.21 Increment Counter (0x15)

Table 4-22

The Increment Counter command takes as arguments the tag Counter Number and the Increment Value to be added to the counter value. This command increments the counter value of the counter pointed to by the Counter Number. The Increment Value is a 24-bit unsigned number with the least significant byte first.

4.2.22 Check Tearing Event (0x16)

Table 4-23

The Check Tearing Event command takes as arguments the tag Counter Number and the Data Buffer Offset in bytes. This command checks whether there was a tearing event in the counter and stores the flag value in the Data Buffer at the Data Buffer Offset index. The value '0x00' is stored if there has been no tearing event, and '0x01' is stored if a tearing event occurred.

4.2.23 Password Authentication (0x17)

Table 4-24

The Password Authentication command takes as arguments the Data Buffer Offset where the PACK result will be stored, the Password Number and (optionally) the Password. This command tries to authenticate the tag using the chosen (pointed to) password. The 2-byte PACK result is stored in the Data Buffer under Data Buffer Offset index.

4.2.24 Halt (0x18)

The Halt command takes no arguments. It halts the tag and turns off the RF field. It must be executed at the end of each operation on a tag to disable the antenna and reduce the power consumption.

4.2.25 Calculate CRC (0x19)

Table 4-25

The Calculate CRC command takes as arguments the Memory Address, the Data Length in bytes and the Data Buffer Offset. The CRC calculation starts at the byte pointed at by the memory address and is done on the 'Data Length' number of bytes in the volatile memory. The result is stored in the Data Buffer at the Data Buffer Offset index. The result is a 16-bit unsigned value with least significant byte first. The CRC calculation code is shown below.

static const uint16 t CCITTCRCTable $[256] = \{$ 0x0000, 0x1021, 0x2042, 0x3063, 0x4084, 0x50a5, 0x60c6, 0x70e7, 0x8108, 0x9129, 0xa14a, 0xb16b, 0xc18c, 0xd1ad, 0xe1ce, 0xf1ef, 0x1231, 0x0210, 0x3273, 0x2252, 0x52b5, 0x4294, 0x72f7, 0x62d6, 0x9339, 0x8318, 0xb37b, 0xa35a, 0xd3bd, 0xc39c, 0xf3ff, 0xe3de, 0x2462, 0x3443, 0x0420, 0x1401, 0x64e6, 0x74c7, 0x44a4, 0x5485, 0xa56a, 0xb54b, 0x8528, 0x9509, 0xe5ee, 0xf5cf, 0xc5ac, 0xd58d, 0x3653, 0x2672, 0x1611, 0x0630, 0x76d7, 0x66f6, 0x5695, 0x46b4, 0xb75b, 0xa77a, 0x9719, 0x8738, 0xf7df, 0xe7fe, 0xd79d, 0xc7bc, 0x48c4, 0x58e5, 0x6886, 0x78a7, 0x0840, 0x1861, 0x2802, 0x3823, 0xc9cc, 0xd9ed, 0xe98e, 0xf9af, 0x8948, 0x9969, 0xa90a, 0xb92b, 0x5af5, 0x4ad4, 0x7ab7, 0x6a96, 0x1a71, 0x0a50, 0x3a33, 0x2a12, 0xdbfd, 0xcbdc, 0xfbbf, 0xeb9e, 0x9b79, 0x8b58, 0xbb3b, 0xab1a, 0x6ca6, 0x7c87, 0x4ce4, 0x5cc5, 0x2c22, 0x3c03, 0x0c60, 0x1c41, 0xedae, 0xfd8f, 0xcdec, 0xddcd, 0xad2a, 0xbd0b, 0x8d68, 0x9d49, 0x7e97, 0x6eb6,


```
0x5ed5, 0x4ef4, 0x3e13, 0x2e32, 0x1e51, 0x0e70, 
0xff9f, 0xefbe, 0xdfdd, 0xcffc, 0xbf1b, 0xaf3a, 
0x9f59, 0x8f78, 0x9188, 0x81a9, 0xb1ca, 0xa1eb, 
0xd10c, 0xc12d, 0xf14e, 0xe16f, 0x1080, 0x00a1, 
0x30c2, 0x20e3, 0x5004, 0x4025, 0x7046, 0x6067, 
0x83b9, 0x9398, 0xa3fb, 0xb3da, 0xc33d, 0xd31c, 
0xe37f, 0xf35e, 0x02b1, 0x1290, 0x22f3, 0x32d2, 
0x4235, 0x5214, 0x6277, 0x7256, 0xb5ea, 0xa5cb, 
0x95a8, 0x8589, 0xf56e, 0xe54f, 0xd52c, 0xc50d, 
0x34e2, 0x24c3, 0x14a0, 0x0481, 0x7466, 0x6447, 
0x5424, 0x4405, 0xa7db, 0xb7fa, 0x8799, 0x97b8, 
0xe75f, 0xf77e, 0xc71d, 0xd73c, 0x26d3, 0x36f2, 
0x0691, 0x16b0, 0x6657, 0x7676, 0x4615, 0x5634, 
0xd94c, 0xc96d, 0xf90e, 0xe92f, 0x99c8, 0x89e9, 
0xb98a, 0xa9ab, 0x5844, 0x4865, 0x7806, 0x6827, 
0x18c0, 0x08e1, 0x3882, 0x28a3, 0xcb7d, 0xdb5c, 
0xeb3f, 0xfb1e, 0x8bf9, 0x9bd8, 0xabbb, 0xbb9a, 
0x4a75, 0x5a54, 0x6a37, 0x7a16, 0x0af1, 0x1ad0, 
0x2ab3, 0x3a92, 0xfd2e, 0xed0f, 0xdd6c, 0xcd4d, 
0xbdaa, 0xad8b, 0x9de8, 0x8dc9, 0x7c26, 0x6c07, 
0x5c64, 0x4c45, 0x3ca2, 0x2c83, 0x1ce0, 0x0cc1, 
0xef1f, 0xff3e, 0xcf5d, 0xdf7c, 0xaf9b, 0xbfba, 
0x8fd9, 0x9ff8, 0x6e17, 0x7e36, 0x4e55, 0x5e74, 
0x2e93, 0x3eb2, 0x0ed1, 0x1ef0 };
```

```
static uint16_t GetCCITTCRC(const uint8_t* Data, uint32_t Size) {
uint16_t CRC; 
uint16_t Temp; 
uint32_t Index; 
if (Size == 0) {
return 0;
```

```
}
```

```
CRC = 0xFFF;
```

```
for (Index = 0; Index < Size; Index + +\}Temp = (uint16_t)( (CRC >> 8) ^ Data[Index] ) & 0x00FF; 
CRC = CCITTCRCTable[Temp] ^ (CRC << 8); 
} 
return CRC;
```
}

4.2.26 Copy Data (0x1A)

Table 4-26

The Copy Data command copies data around inside the RFID Module memory. This command takes as arguments the Destination Address, the Source Address and the Data Length in bytes.

4.2.27 Unlock (0x1B)

Table 4-27

The Unlock command takes as an argument an 8-bytes long password. If the password matches the actual password in the module the memory which contains:

- AES Initialization Vectors
- AES Encryption Keys

- Authentication Keys and Passwords
- User Memory
- Password

will be unlocked and made accessible for read and write operations.

4.2.28 Lock (0x1C)

This Function takes no arguments. If the device is unlocked and the user executes the Lock command, protected memory which contains:

- AES Initialization Vectors
- AES Encryption Keys
- Authentication Keys and Passwords
- User Memory
- Password

will be saved to non-volatile memory and locked (direct access will be blocked).

4.2.29 Get Module Version (0x1D)

The Get Module Version command takes no arguments. After execution, the Data Buffer at index 0x00 is filled with a NULL-ended ASCII string which describes the hardware and firmware version of the module.

4.2.30 Reset to Default (0x1E)

The Reset to Default command doesn't take any arguments. It resets all memory including protected memory to factory default settings. All locked information is lost.

4.2.31 Enumerate Tags UID (0x1F)

The Enumerate Tags UID command doesn't take any arguments. This command scans for all tags within the range of the antenna and enumerates their UIDs in the Data Buffer. The first byte in the Data Buffer says how many tags were discovered. If there are no tags in the range, then the Result Register is updated with 0x08 (Tag is not present) value. After enumeration, if there is at least one tag in the field, the Result Register is updated with 0x00 (No Error) value. Information about all enumerated tags is stored in the Data Buffer, starting from the second byte of the buffer. There are as many records in the Data Buffer as there were tags found. Each record contains the Tag UID Size as the first byte and the Tag UID is stored in the rest of the bytes in the record. The record length is variable and equal to the UID size plus one byte.

4.2.32 Enumerate Tags UID and Type (0x20)

The Enumerate Tags UID and Type command doesn't take any arguments. This command scans for all tags within the range of the antenna and enumerates their UID and type in the Data Buffer. The first byte in the Data Buffer says how many tags were discovered. If there are no tags in the range, then the Result Register is updated with 0x08 (Tag is not present) value. After enumeration, if there is at least one tag in the field the Result Register is updated with 0x00 (No Error) value. Information about all enumerated tags is stored in the Data Buffer starting from the second byte of the buffer. There are as many records in the Data Buffer as there were tags found. Each record contains the Tag Type as the first byte, the Tag UID Size as the second byte and the Tag UID is stored in the rest of the bytes in the record. The record length is variable and equal to the UID size plus two bytes.

4.2.33 Select Tag (0x21)

Table 4.28

The Select Tag command takes as arguments the UID Size in bytes and the UID bytes themselves. If the command execution is successful, then the addressed tag is selected and the Result Register is updated with 0x00 (No Error) value.

4.2.34 Polling (0x22)

Table 4.29

Polling Period – Delay between polling attempts for detection of tag presence.

Number Of Defined Tags – States how many UIDs are stored in user memory.

Defined Tag: TPI Config – Bits 4 to 5 defines which TPI is used. Bits 0 to 2 defines state on defined Tag presence. See Table 4.30 for details.

Defined Tag: Tag Timeout – T = Value x 100mS. Delay between Tag detection and polling restart.

Undefined Tag: TPI Config – Bits 4 to 5 defines which TPI is used. Bits 0 to 2 defines the state on defined Tag presence. See Table 4.30 for details.

Undefined Tag: Tag Timeout – T = Value x 100mS. Delay between Tag detection and polling restart.

Table 4.30

The Polling command takes arguments listed in the above tables. If the command execution is successful, then the Result Register is updated with 0x00 (No Error) value. For more information about the Polling command, please see chapte[r 4.3 Polling Mode.](#page-54-0)

4.3 Polling Mode

This command enables Tag UID polling. In this mode the module executes the continuous repeated Enumerate Tags UID command, writes the Command Polling value to the Command Register and parameters to the Parameter Registers to start polling. Parameters are described i[n 4.2.34.](#page-52-1) Polling mode can be stopped by writing any data to the Module memory.

4.3.1 Polling internal procedure

Figure 4.3

4.3.2 Polling auto start

Polling can be automatically started after power-up or reset. This functionality requires that a command and parameters be written to User Memory from index 0. Polling mode can be stopped by writing any data to the Module memory. Access to User Memory must be unlocked by the Unlock command before writing data. Data in the User Memory is saved after a Lock command.

The command followed by parameters is stored in User Memory from index 0x00.

Table 4.31

4.3.3 Polling use simple example – Undefined Tags

This example of polling command use demonstrates how to configure simple signalling of any Tag presence. Command Polling followed by parameters:

Table 4.32

The module will check for Tag presence every 100ms. If a tag is detected, then TPI1 is set to the High state. After 5 seconds TPI1 is set to the Low state and the module checks for Tag presence once again.

4.3.4 Polling use simple example – Defined Tags

This example of a polling command use shows how to configure simple signalling of only defined Tag presence. Command Polling followed by parameters:

Table 4.33

The module will check for Tag presence every 100ms. If a tag is detected and this tag is on the defined list in user memory then TPI0 is set to the Low state. After 10 seconds TPI0 is set to the Low state and the module checks for Tag presence once again. The module does not react (only waits 100mS) if the detected Tag is not on list. Information of how to prepare a Defined Tag List can be found in chapter 4.3.6.

4.3.5 Polling use simple example – Undefined and Defined Tags

This example of polling command use shows how to configure simple signalling of defined and undefined Tag presence. Command Polling followed by parameters:

Table 4.34

The module will test for Tag presence every 100ms. If a tag is detected and this tag is on the defined list in the user memory, then TPI0 is set to the high state. After 10 seconds TPI0 is set to the low state and the module checks for tag presence once again. If the tag is not on the defined List, then TPI1 is set to the low state. After 3 seconds TPI1 is set to the high state and the module checks for Tag presence once again. Information of how to prepare a defined Tag List can be found in chapter 4.3.6.

4.3.6 Defined Tag List

The defined tag list is stored in User Memory from index 0x07. UIDSize followed by UID must be written for every listed tag.

If any of UIDSize contains an incorrect value (allowed values are 4, 7 and 10) then the module stops comparing the detected UID to the list. The number of Tags on the list is defined by the 'Number of Defined Tags' polling command parameter. Access to User Memory must be unlocked by the Unlock command before writing data. Data in User Memory is saved after a Lock command.

4.3.7 Polling stop and communication with tag example.

To communicate with one of the detected tags user must send Select Tag command followed by UID.

Figure 4.4 Polling stop and Tag r/w example.

4.4 Power Down Mode

The Power Down functionality is provided via the nPWRDN pin. This pin is configured as an input without any pull-up or pull-down resistors, thus it must be driven by the user. For normal operation, this pin should be connected to VCC or driven high. After the user drives this pin low the system prepares to go into Power Down Mode and drives nBUSY line low just before entering this state. To wake up module the nPWRDN pin must be driven high. The module will restart and drive the nBUSY line low to inform the user that it is starting the system again. When start-up is completed, the nBUSY line is driven high. During the restart, none of the volatile memory content is retained.

The reaction time of the system will vary depending upon the active configuration and the following:

- Communication if the line is pulled low whilst there is ongoing communication, then the Power Manager will wait until communication is finished before going into the Power Down state.
- Command execution if the line is pulled low whilst the system is executing a command, then the Power Manager will wait until command execution is finished before going into the Power Down state.

When the module is in the Power Down state, all systems are turned-off and all IOs are in a high impedance state. The module won't respond to any signals or communication coming from outside.

4.5 Memory Locking

The module has an option to lock its memory using an 8-byte long password. The default state of the module after power-up is locked and the default password is all bytes equal to 0x00. The user cannot see the memory content when the module is locked, and any write to the memory will be discarded. During the unlock procedure, the content of the non-volatile memory is copied to the volatile memory and is available for reading and writing. Changes made in volatile memory are updated in non-volatile memory during the lock procedure. Commands using data from this memory range use the volatile memory when the module is unlocked and non-volatile memory when the module is locked

5 Communication Interface

5.1 nBUSY line

The RFID A1 Module provides two communication buses: SPI Slave and I2C Slave. Each of these buses has an equal level of access to the module internal registers. These buses operate independently but the user cannot use both at the same time. The module has no safety system and doing so can make the system work in unpredictable ways.

The RFID A1 Module provides nBUSY signal line which shares the functionality between both communication buses. This signal line is configured as an output. The nBUSY signal is driven low by the module when anything is written to the memory. At the same time when nBUSY line is driven low the result register is set to 0xFF value. This inform the user that module is busy and is parsing data or executing command and nothing else should be written to the module now. The user must wait for nBUSY line going high or Result Register value changing from 0xFF to perform any operation on the module or to read any data from the module which is a result of current processing.

5.2 SPI Bus

The RFID A1 Module provides a SPI slave bus. This is one of two buses provided by the module. It is the fastest communication option in the module. The high clock frequency together with simultaneous communication in both directions gives the fastest data exchange rate, but requires the most number of IO lines.

5.2.1 Bus signals and timings requirements

There are four signals necessary for communication via SPI bus. These are CLK, MOSI, MISO and CS. CLK (Clock), CS (Chip Select) and MOSI (Master Out Slave In) signals are controlled by the master on every transmission and these pins on the RFID A1 module are configured as inputs without any pull-up or pull-down resistors. MISO (Master In Slave Out) signal is driven by the module to Vcc and to ground when there is ongoing transmission. During communication via the SPI bus the user must pay attention to the nBUSY signal or must be sure that the Result Register has other value than 0xFF when writing to the module.

The SPI frame timing together with clock polarisation and phase is shown in [Figure 5.2-1.](#page-61-1)

As [Figure 5.2-1](#page-61-1) shows the clock idle state is high. On a falling edge of the clock the module switches the state of the data line (MISO) and on a rising edge of the clock the module samples the data line (MOSI). During transmission, the most significant bit is transmitted first. One frame is eight bits long (one byte). The CS signal is active low thus the master must drive the CS signal low before clocking the clock signal. The maximum clock frequency for the system is 500 kHz. An additional restriction is also placed on the minimum allowed time between the CS line going low and the first clock falling edge and the minimum allowed time for CS being high between transmission (please see chapte[r 2.7\)](#page-8-2). This is the result of low power consumption as the system must jump from sleep state to run state at various points during communication.

5.2.2 Exchanging data (reading or reading and writing)

The SPI bus is a type of synchronous bus where both sides (master and slave) synchronised by the master clock send data at the same time in both directions. Each transmission is initiated when the master drives the CS line low and ends when the CS line is driven high. All bytes transmitted in each direction between falling and rising edge of the CS line are considered as one packet and bytes within those packets are ordered. The First byte is transmitted always after the falling edge of the CS signal. A typical packet exchange is shown in [Table 5-1.](#page-61-2)

When the CS signal goes low the byte index resets. The first three bytes transmitted from the module (MISO line) are trash and should be discarded by the master. The first three bytes transmitted from the master (MOSI line) are used to set up the transmission starting from byte four onwards. The first byte is the least significant byte of the address to where to write to or from where to read the data. The second byte is the most significant byte of the address. The address is represented as an unsigned 16-bit number. The third byte is the write indicator flag. If the master wants the data sent via the MOSI line to be written to the module's registers, then the third byte in the packet should have the value '0x01'. If the master sets this byte to value '0x00', then the RFID A1 module will discard incoming data (will not store them) and just stream the data on the MISO line from the RFID A1 memory.

5.2.3 Command processing

The RFID A1 module after CS goes high (end of communication) starts to process the data. Even if the master didn't write anything to the memory and just read data from the module, it will still drive the nBUSY line low for a while to do internal housekeeping. A typical communication timing is shown below i[n Figure 5.2-2.](#page-62-1)

Figure 5.2-2

5.3 I2C Bus

The RFID A1 Module provides a I2C slave bus. The advantages of the I2C are that it requires only two lines to communicate and that multiple devices can be connected to the 2-wire same bus.

The implementation of I2C in the module complies with the 7-bit addressing I2C standard, thus the timing diagrams for bit and frame representation will be omitted in this document. Each transmission on the I2C line begins with an I2C START condition and ends with an I2C STOP condition. In between these events there could be as many repeated STARTs as the user wants. The behaviour of the nBUSY line is the same for all buses, thus as well for I2C. The module starts processing the command after receiving a stop condition on I2C bus.

5.3.1 I2C Address Construction

After reset or power-up when the module starts, it samples three IO lines (I2C AD0, I2C AD1 and I2C AD2) to construct the I2C address to which it will respond. Thanks to this solution there is the possibility to connect to eight such modules on one I2C bus. Byte wise the address sent on the I2C bus after START is constructed as follows:

Table 5-2

5.3.2 Writing to the memory

Writing to the RFID A1 Module memory via the I2C bus is very similar to other memory type devices. A typical transmission for writing to the registers is shown below in [Table 5-3.](#page-63-3)

Table 5-3

The master must initiate the transmission on the I2C bus by sending a START condition followed by the I2C address of the slave device with R/W byte set to 0. The second and third byte sent by the master will set up the address index under which the data will be stored. Starting from the fourth byte the data will be stored under the passed address value and the memory index will be automatically increased by one each time. When the RFID A1 module recognizes its I2C address it drives the nBUSY line low to inform the master that it is receiving the data.

5.3.3 Reading from the memory

The master can read from the module in two possible ways presented below in [Table 5-4](#page-64-1) an[d Table 5-5.](#page-64-2)

Table 5-5

As shown i[n Table 5-4,](#page-64-1) to read from a certain address the master first must write the address index by sending a START condition followed by the I2C address to be written and two bytes representing the index. Later and without sending a STOP condition the master must send a repeated START signal, this time followed by the I2C address to be read. Bytes after the I2C read address will be transmitted from the module and will contain memory data starting from the address index set up by the master at the beginning of the transmission. The index is automatically increased and stored between transmissions. If the user wants to continue reading data starting from the index value left since last transmission, he can use the simplified version of the transmission shown i[n Table 5-5.](#page-64-2)

6 Mechanical

6.1 Dimensions

Drawing 1

| UNIT | | A1 | | | | |
|-------------|-------|-------|-------|-------|-------|-------|
| mm | | | 20.06 | 20.32 | 1 2 7 | 3.68 |
| inches | 0.079 | 27.56 | -70 | 0.8 | 0.05 | 0.145 |

Table 6-1

6.2 Recommended Footprint

Drawing 2

MIFARE, MIFARE Ultralight, MIFARE Plus, MIFARE Classic, NTAG and MIFARE DESFire are trademarks of NXP B.V.

No responsibility is taken for the method of integration or final use of the A1 module

More information about the A1 module and other products can be found at the Internet site:

http://www.eccel.co.uk

or alternatively contact ECCEL Technology (IB Technology) by e-mail at:

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