

The revision list can be viewed directly by clicking the title page. The revision list summarizes the locations of revisions and additions. Details should always be checked by referring to the relevant text.

SH7261 Group

User's Manual: Hardware

Renesas 32-Bit RISC Microcomputer SuperH[™] RISC engine Family / SH7260 Series

> R5S72611 R5S72612 R5S72613

Notice

- 1. All information included in this document is current as of the date this document is issued. Such information, however, is subject to change without any prior notice. Before purchasing or using any Renesas Electronics products listed herein, please confirm the latest product information with a Renesas Electronics sales office. Also, please pay regular and careful attention to additional and different information to be disclosed by Renesas Electronics such as that disclosed through our website.
- 2. Renesas Electronics does not assume any liability for infringement of patents, copyrights, or other intellectual property rights of third parties by or arising from the use of Renesas Electronics products or technical information described in this document. No license, express, implied or otherwise, is granted hereby under any patents, copyrights or other intellectual property rights of Renesas Electronics or others.
- 3. You should not alter, modify, copy, or otherwise misappropriate any Renesas Electronics product, whether in whole or in part.
- 4. Descriptions of circuits, software and other related information in this document are provided only to illustrate the operation of semiconductor products and application examples. You are fully responsible for the incorporation of these circuits, software, and information in the design of your equipment. Renesas Electronics assumes no responsibility for any losses incurred by you or third parties arising from the use of these circuits, software, or information.
- 5. When exporting the products or technology described in this document, you should comply with the applicable export control laws and regulations and follow the procedures required by such laws and regulations. You should not use Renesas Electronics products or the technology described in this document for any purpose relating to military applications or use by the military, including but not limited to the development of weapons of mass destruction. Renesas Electronics products and technology may not be used for or incorporated into any products or systems whose manufacture, use, or sale is prohibited under any applicable domestic or foreign laws or regulations.
- 6. Renesas Electronics has used reasonable care in preparing the information included in this document, but Renesas Electronics does not warrant that such information is error free. Renesas Electronics assumes no liability whatsoever for any damages incurred by you resulting from errors in or omissions from the information included herein.
- 7. Renesas Electronics products are classified according to the following three quality grades: "Standard", "High Quality", and "Specific". The recommended applications for each Renesas Electronics product depends on the product's quality grade, as indicated below. You must check the quality grade of each Renesas Electronics product before using it in a particular application. You may not use any Renesas Electronics product for any application categorized as "Specific" without the prior written consent of Renesas Electronics. Further, you may not use any Renesas Electronics product for any application for which it is not intended without the prior written consent of Renesas Electronics. Renesas Electronics shall not be in any way liable for any damages or losses incurred by you or third parties arising from the use of any Renesas Electronics product for an application categorized as "Specific" or for which the product is not intended where you have failed to obtain the prior written consent of Renesas Electronics. The quality grade of each Renesas Electronics product is "Standard" unless otherwise expressly specified in a Renesas Electronics data sheets or data books, etc.
	- "Standard": Computers; office equipment; communications equipment; test and measurement equipment; audio and visual equipment; home electronic appliances; machine tools; personal electronic equipment; and industrial robots.
	- "High Quality": Transportation equipment (automobiles, trains, ships, etc.); traffic control systems; anti-disaster systems; anti crime systems; safety equipment; and medical equipment not specifically designed for life support.
	- "Specific": Aircraft; aerospace equipment; submersible repeaters; nuclear reactor control systems; medical equipment or systems for life support (e.g. artificial life support devices or systems), surgical implantations, or healthcare intervention (e.g. excision, etc.), and any other applications or purposes that pose a direct threat to human life.
- 8. You should use the Renesas Electronics products described in this document within the range specified by Renesas Electronics, especially with respect to the maximum rating, operating supply voltage range, movement power voltage range, heat radiation characteristics, installation and other product characteristics. Renesas Electronics shall have no liability for malfunctions or damages arising out of the use of Renesas Electronics products beyond such specified ranges.
- 9. Although Renesas Electronics endeavors to improve the quality and reliability of its products, semiconductor products have specific characteristics such as the occurrence of failure at a certain rate and malfunctions under certain use conditions. Further, Renesas Electronics products are not subject to radiation resistance design. Please be sure to implement safety measures to guard them against the possibility of physical injury, and injury or damage caused by fire in the event of the failure of a Renesas Electronics product, such as safety design for hardware and software including but not limited to redundancy, fire control and malfunction prevention, appropriate treatment for aging degradation or any other appropriate measures. Because the evaluation of microcomputer software alone is very difficult, please evaluate the safety of the final products or system manufactured by you.
- 10. Please contact a Renesas Electronics sales office for details as to environmental matters such as the environmental compatibility of each Renesas Electronics product. Please use Renesas Electronics products in compliance with all applicable laws and regulations that regulate the inclusion or use of controlled substances, including without limitation, the EU RoHS Directive. Renesas Electronics assumes no liability for damages or losses occurring as a result of your noncompliance with applicable laws and regulations.
- 11. This document may not be reproduced or duplicated, in any form, in whole or in part, without prior written consent of Renesas Electronics.
- 12. Please contact a Renesas Electronics sales office if you have any questions regarding the information contained in this document or Renesas Electronics products, or if you have any other inquiries.
- (Note 1) "Renesas Electronics" as used in this document means Renesas Electronics Corporation and also includes its majority owned subsidiaries.
- (Note 2) "Renesas Electronics product(s)" means any product developed or manufactured by or for Renesas Electronics.

General Precautions in the Handling of MPU/MCU Products

The following usage notes are applicable to all MPU/MCU products from Renesas. For detailed usage notes on the products covered by this manual, refer to the relevant sections of the manual. If the descriptions under General Precautions in the Handling of MPU/MCU Products and in the body of the manual differ from each other, the description in the body of the manual takes precedence.

1. Handling of Unused Pins

Handle unused pins in accord with the directions given under Handling of Unused Pins in the manual.

- ⎯ The input pins of CMOS products are generally in the high-impedance state. In operation with an unused pin in the open-circuit state, extra electromagnetic noise is induced in the vicinity of LSI, an associated shoot-through current flows internally, and malfunctions may occur due to the false recognition of the pin state as an input signal. Unused pins should be handled as described under Handling of Unused Pins in the manual.
- 2. Processing at Power-on

The state of the product is undefined at the moment when power is supplied.

⎯ The states of internal circuits in the LSI are indeterminate and the states of register settings and pins are undefined at the moment when power is supplied.

In a finished product where the reset signal is applied to the external reset pin, the states of pins are not guaranteed from the moment when power is supplied until the reset process is completed.

In a similar way, the states of pins in a product that is reset by an on-chip power-on reset function are not guaranteed from the moment when power is supplied until the power reaches the level at which resetting has been specified.

3. Prohibition of Access to Reserved Addresses

Access to reserved addresses is prohibited.

- ⎯ The reserved addresses are provided for the possible future expansion of functions. Do not access these addresses; the correct operation of LSI is not guaranteed if they are accessed.
- 4. Clock Signals

After applying a reset, only release the reset line after the operating clock signal has become stable. When switching the clock signal during program execution, wait until the target clock signal has stabilized.

- ⎯ When the clock signal is generated with an external resonator (or from an external oscillator) during a reset, ensure that the reset line is only released after full stabilization of the clock signal. Moreover, when switching to a clock signal produced with an external resonator (or by an external oscillator) while program execution is in progress, wait until the target clock signal is stable.
- 5. Differences between Products

Before changing from one product to another, i.e. to one with a different type number, confirm that the change will not lead to problems.

⎯ The characteristics of MPU/MCU in the same group but having different type numbers may differ because of the differences in internal memory capacity and layout pattern. When changing to products of different type numbers, implement a system-evaluation test for each of the products.

Configuration of This Manual

This manual comprises the following items:

- 1. General Precautions in the Handling of MPU/MCU Products
- 2. Configuration of This Manual
- 3. Preface
- 4. Contents
- 5. Overview
- 6. Description of Functional Modules
	- CPU and System-Control Modules
	- On-Chip Peripheral Modules The configuration of the functional description of each module differs according to the module. However, the generic style includes the following items:
	- i) Feature
	- ii) Input/Output Pin
	- iii) Register Description
	- iv) Operation
	- v) Usage Note

When designing an application system that includes this LSI, take notes into account. Each section includes notes in relation to the descriptions given, and usage notes are given, as required, as the final part of each section.

- 7. List of Registers
- 8. Electrical Characteristics
- 9. Appendix
	- Product Type, Package Dimensions, etc.

10. Main Revisions for This Edition (only for revised versions)

The list of revisions is a summary of points that have been revised or added to earlier versions. This does not include all of the revised contents. For details, see the actual locations in this manual.

11. Index

Preface

This LSI is an RISC (Reduced Instruction Set Computer) microcomputer that includes a Renesas original RISC CPU as its core, and the peripheral functions required to configure a system.

- Target Users: This manual was written for users who will be using this LSI in the design of application systems. Target users are expected to understand the fundamentals of electrical circuits, logical circuits, and microcomputers.
- Objective: This manual was written to explain the hardware functions and electrical characteristics of this LSI to the target users.

Notes on reading this manual:

- In order to understand the overall functions of the chip Read the manual according to the contents. This manual can be roughly categorized into parts on the CPU, system control functions, peripheral functions and electrical characteristics.
- In order to understand the details of the CPU's functions Read the SH-2A, SH2A-FPU Software Manual.
- In order to understand the details of a register when its name is known Read the index that is the final part of the manual to find the page number of the entry on the register. The addresses, bits, and initial values of the registers are summarized in section 30, List of Registers.

• Description of Numbers and Symbols

Aspects of the notations for register names, bit names, numbers, and symbolic names in this manual are explained below.

• Description of Registers

Each register description includes a bit chart, illustrating the arrangement of bits, and a table of bits, describing the meanings of the bit settings. The standard format and notation for bit charts and tables are described below.

All trademarks and registered trademarks are the property of their respective owners.

Contents

Section 1 Overview

1.1 SH7261 Group Features

This LSI is a single-chip RISC (Reduced Instruction Set Computer) microprocessor that integrates a Renesas original RISC CPU core with peripheral functions required for system configuration.

The CPU incorporated in this LSI is the SH-2A CPU, which features upward compatibility on the object code level with the SH-1, SH-2, and SH-2E microcomputers. The CPU has a RISC-type instruction set and employs a superscalar architecture and the Harvard architecture, which greatly improves instruction execution speed. In addition, the 32-bit internal-bus architecture independent of the bus for the direct memory access controller (DMAC) enhances data processing power. This CPU realizes low-cost, high-performance, and high-functioning systems for applications such as high-speed realtime control, which has been next to impossible with the conventional microcomputers.

This LSI has a floating-point unit and a cache.

In addition, this LSI includes on-chip peripheral functions necessary for system configuration, such as, 32-Kbyte RAM for high-speed operation, a controller area network (RCAN-ET) $*$ ¹, $IEBus^{TM*2}$ controller (IEB)^{*3}, CD-ROM decoder (ROM-DEC), a serial sound interface (SSI), a serial communication interface with FIFO (SCIF), 1^2C bus interface 3 (IIC3), a multi-function timer pulse unit 2 (MTU2), an 8-bit timer (TMR), a realtime clock (RTC), an A/D converter, a D/A converter, an interrupt controller (INTC), I/O ports, and advanced user debugger II (AUD-II).

This LSI also provides an external memory access support function to enable direct connection to various memory devices or peripheral LSIs. These on-chip functions significantly reduce costs of designing and manufacturing application systems.

The features of this LSI are listed in table 1.1.

Notes: 1. R5S72611 and R5S72613 support this function.

- 2. The IEBusTM (Inter Equipment BusTM) is a trademark of Renesas Electronics Corporation.
- 3. R5S72612 and R5S72613 support this function.

Table 1.1 SH7261 Group Features

1.2 Product Lineup

Table 1.2 Product Lineup

1.3 Block Diagram

The block diagram of this LSI is shown in figure 1.1.

Figure 1.1 Block Diagram

1.4 Pin Assignments

Figure 1.2 Pin Assignments of R5S72611

Figure 1.3 Pin Assignments of R5S72612

Figure 1.4 Pin Assignments of R5S72613

1.5 Pin Functions

Table 1.3 lists the pin functions.

Table 1.3 Pin Functions

Note: $*$ The pin with the pull-up function.

Section 2 CPU

2.1 Register Configuration

The register set consists of sixteen 32-bit general registers, four 32-bit control registers, and four 32-bit system registers.

2.1.1 General Registers

Figure 2.1 shows the general registers.

The sixteen 32-bit general registers are numbered R0 to R15. General registers are used for data processing and address calculation. R0 is also used as an index register. Several instructions have R0 fixed as their only usable register. R15 is used as the hardware stack pointer (SP). Saving and restoring the status register (SR) and program counter (PC) in exception handling is accomplished by referencing the stack using R15.

2. R15 functions as a hardware stack pointer (SP) during exception processing.

Figure 2.1 General Registers

2.1.2 Control Registers

The control registers consist of four 32-bit registers: the status register (SR), the global base register (GBR), the vector base register (VBR), and the jump table base register (TBR).

The status register indicates instruction processing states.

The global base register functions as a base address for the GBR indirect addressing mode to transfer data to the registers of on-chip peripheral modules.

The vector base register functions as the base address of the exception handling vector area (including interrupts).

The jump table base register functions as the base address of the function table area.

Figure 2.2 Control Registers

(1) Status Register (SR)

(2) Global Base Register (GBR)

GBR is referenced as the base address in a GBR-referencing MOV instruction.

(3) Vector Base Register (VBR)

VBR is referenced as the branch destination base address in the event of an exception or an interrupt.

(4) Jump Table Base Register (TBR)

TBR is referenced as the start address of a function table located in memory in a JSR/N@@(disp8,TBR) table-referencing subroutine call instruction.

2.1.3 System Registers

The system registers consist of four 32-bit registers: the high and low multiply and accumulate registers (MACH and MACL), the procedure register (PR), and the program counter (PC). MACH and MACL store the results of multiply or multiply and accumulate operations. PR stores the return address from a subroutine procedure. PC points four bytes ahead of the current instruction and controls the flow of the processing.

Figure 2.3 System Registers

(1) Multiply and Accumulate Register High (MACH) and Multiply and Accumulate Register Low (MACL)

MACH and MACL are used as the addition value in a MAC instruction, and store the result of a MAC or MUL instruction.

(2) Procedure Register (PR)

PR stores the return address of a subroutine call using a BSR, BSRF, or JSR instruction, and is referenced by a subroutine return instruction (RTS).

(3) Program Counter (PC)

PC points four bytes ahead of the instruction being executed.

2.1.4 Register Banks

For the nineteen 32-bit registers comprising general registers R0 to R14, control register GBR, and system registers MACH, MACL, and PR, high-speed register saving and restoration can be carried out using a register bank. The register contents are automatically saved in the bank after the CPU accepts an interrupt that uses a register bank. Restoration from the bank is executed by issuing a RESBANK instruction in an interrupt processing routine.

This LSI has 15 banks. For details, see the SH-2A, SH2A-FPU Software Manual and section 6.8, Register Banks.

2.1.5 Initial Values of Registers

Table 2.1 lists the values of the registers after a reset.

2.2 Data Formats

2.2.1 Data Format in Registers

Register operands are always longwords (32 bits). If the size of memory operand is a byte (8 bits) or a word (16 bits), it is changed into a longword by expanding the sign-part when loaded into a register.

Figure 2.4 Data Format in Registers

2.2.2 Data Formats in Memory

Memory data formats are classified into bytes, words, and longwords. Memory can be accessed in 8-bit bytes, 16-bit words, or 32-bit longwords. A memory operand of fewer than 32 bits is stored in a register in sign-extended or zero-extended form.

A word operand should be accessed at a word boundary (an even address of multiple of two bytes: address 2n), and a longword operand at a longword boundary (an even address of multiple of four bytes: address 4n). Otherwise, an address error will occur. A byte operand can be accessed at any address.

Only big-endian byte order can be selected for the data format.

Data formats in memory are shown in figure 2.5.

Figure 2.5 Data Formats in Memory

2.2.3 Immediate Data Format

Byte (8-bit) immediate data is located in an instruction code. Immediate data accessed by the MOV, ADD, and CMP/EQ instructions is sign-extended and handled in registers as longword data. Immediate data accessed by the TST, AND, OR, and XOR instructions is zero-extended and handled as longword data. Consequently, AND instructions with immediate data always clear the upper 24 bits of the destination register.

20-bit immediate data is located in the code of a MOVI20 or MOVI20S 32-bit transfer instruction. The MOVI20 instruction stores immediate data in the destination register in sign-extended form. The MOVI20S instruction shifts immediate data by eight bits in the upper direction, and stores it in the destination register in sign-extended form.

Word or longword immediate data is not located in the instruction code, but rather is stored in a memory table. The memory table is accessed by an immediate data transfer instruction (MOV) using the PC relative addressing mode with displacement.

See examples given in section 2.3.1 (10), Immediate Data.

2.3 Instruction Features

2.3.1 RISC-Type Instruction Set

Instructions are RISC type. This section details their functions.

(1) 16-Bit Fixed-Length Instructions

Basic instructions have a fixed length of 16 bits, improving program code efficiency.

(2) 32-Bit Fixed-Length Instructions

The SH-2A additionally features 32-bit fixed-length instructions, improving performance and ease of use.

(3) One Instruction per State

Each basic instruction can be executed in one cycle using the pipeline system.

(4) Data Length

.DATA.W H'1234

Longword is the standard data length for all operations. Memory can be accessed in bytes, words, or longwords. Byte or word data in memory is sign-extended and handled as longword data. Immediate data is sign-extended for arithmetic operations or zero-extended for logic operations. It is also handled as longword data.

instruction.

Table 2.2 Sign Extension of Word Data

Note: @(disp, PC) accesses the immediate data.

(5) Load-Store Architecture

Basic operations are executed between registers. For operations that involve memory access, data is loaded to the registers and executed (load-store architecture). Instructions such as AND that manipulate bits, however, are executed directly in memory.

(6) Delayed Branch Instructions

With the exception of some instructions, unconditional branch instructions, etc., are executed as delayed branch instructions. With a delayed branch instruction, the branch is taken after execution of the instruction immediately following the delayed branch instruction. This reduces disturbance of the pipeline control when a branch is taken.

In a delayed branch, the actual branch operation occurs after execution of the slot instruction. However, instruction execution such as register updating excluding the actual branch operation, is performed in the order of delayed branch instruction \rightarrow delay slot instruction. For example, even though the contents of the register holding the branch destination address are changed in the delay slot, the branch destination address remains as the register contents prior to the change.

(7) Unconditional Branch Instructions with No Delay Slot

The SH-2A additionally features unconditional branch instructions in which a delay slot instruction is not executed. This eliminates unnecessary NOP instructions, and so reduces the code size.

(8) Multiply/Multiply-and-Accumulate Operations

-bit \times 16-bit \rightarrow 32-bit multiply operations are executed in one to two cycles. 16-bit \times 16-bit + -bit \rightarrow 64-bit multiply-and-accumulate operations are executed in two to three cycles. 32-bit \times -bit \rightarrow 64-bit multiply and 32 -bit \times 32-bit + 64-bit \rightarrow 64-bit multiply-and-accumulate operations are executed in two to four cycles.

(9) T Bit

The T bit in the status register (SR) changes according to the result of the comparison. Whether a conditional branch is taken or not taken depends upon the T bit condition (true/false). The number of instructions that change the T bit is kept to a minimum to improve the processing speed.

Table 2.4 T Bit

(10) Immediate Data

Byte immediate data is located in an instruction code. Word or longword immediate data is not located in instruction codes but in a memory table. The memory table is accessed by an immediate data transfer instruction (MOV) using the PC relative addressing mode with displacement.

With the SH-2A, 17- to 28-bit immediate data can be located in an instruction code. However, for 21- to 28-bit immediate data, an OR instruction must be executed after the data is transferred to a register.

Table 2.5 Immediate Data Accessing

Note: @(disp, PC) accesses the immediate data.

(11) Absolute Address

When data is accessed by an absolute address, the absolute address value should be placed in the memory table in advance. That value is transferred to the register by loading the immediate data during the execution of the instruction, and the data is accessed in register indirect addressing mode.

With the SH-2A, when data is referenced using an absolute address not exceeding 28 bits, it is also possible to transfer immediate data located in the instruction code to a register and to reference the data in register indirect addressing mode. However, when referencing data using an absolute address of 21 to 28 bits, an OR instruction must be used after the data is transferred to a register.

(12) 16-Bit/32-Bit Displacement

When data is accessed by 16-bit or 32-bit displacement, the displacement value should be placed in the memory table in advance. That value is transferred to the register by loading the immediate data during the execution of the instruction, and the data is accessed in the indexed indirect register addressing mode.

2.3.2 Addressing Modes

Addressing modes and effective address calculation are as follows:

Table 2.8 Addressing Modes and Effective Addresses

upper bits are sign-extended, and the lower bits are padded with zero.

2.3.3 Instruction Format

The instruction formats and the meaning of source and destination operands are described below. The meaning of the operand depends on the instruction code. The symbols used are as follows:

- xxxx: Instruction code
- mmm: Source register
- nnnn: Destination register
- iiii: Immediate data
- dddd: Displacement

Table 2.9 Instruction Formats

Note: $*$ In multiply-and-accumulate instructions, nnnn is the source register.

2.4 Instruction Set

2.4.1 Instruction Set by Classification

Table 2.10 lists the instructions according to their classification.

Table 2.10 Classification of Instructions

The table below shows the format of instruction codes, operation, and execution states. They are described by using this format according to their classification.

Notes: 1. Instruction execution cycles: The execution cycles shown in the table are minimums. In practice, the number of instruction execution states will be increased in cases such as the following:

- a. When there is a conflict between an instruction fetch and a data access
- b. When the destination register of a load instruction (memory \rightarrow register) is the same as the register used by the next instruction.
- 2. Depending on the operand size, displacement is scaled by \times 1, \times 2, or \times 4. For details, refer to the SH-2A, SH2A-FPU Software Manual.

2.4.2 Data Transfer Instructions

Table 2.11 Data Transfer Instructions

2.4.3 Arithmetic Operation Instructions

Table 2.12 Arithmetic Operation Instructions

2.4.4 Logic Operation Instructions

Table 2.13 Logic Operation Instructions

2.4.5 Shift Instructions

Table 2.14 Shift Instructions

2.4.6 Branch Instructions

Table 2.15 Branch Instructions

Note: * One cycle when the program does not branch.

2.4.7 System Control Instructions

Table 2.16 System Control Instructions

Notes: 1. Instruction execution cycles: The execution cycles shown in the table are minimums. In practice, the number of instruction execution states in cases such as the following:

a. When there is a conflict between an instruction fetch and a data access

b. When the destination register of a load instruction (memory \rightarrow register) is the same as the register used by the next instruction.

* In the event of bank overflow, the number of cycles is 19.

2.4.8 Floating Point Operation Instructions

Table 2.17 Floating Point Operation Instructions

2.4.9 FPU-Related CPU Instructions

Table 2.18 FPU-Related CPU Instructions

2.4.10 Bit Manipulation Instructions

Table 2.19 Bit Manipulation Instructions

2.5 Processing States

The CPU has four processing states: reset, exception handling, program execution, and powerdown. Figure 2.6 shows the transitions between the states.

Figure 2.6 Transitions between Processing States

(1) Reset State

In the reset state, the CPU is reset. There are two kinds of reset, power-on reset and manual reset.

(2) Exception Handling State

The exception handling state is a transient state that occurs when exception handling sources such as resets or interrupts alter the CPU's processing state flow.

For a reset, the initial values of the program counter (PC) (execution start address) and stack pointer (SP) are fetched from the exception handling vector table and stored; the CPU then branches to the execution start address and execution of the program begins.

For an interrupt, the stack pointer (SP) is accessed and the program counter (PC) and status register (SR) are saved to the stack area. The exception service routine start address is fetched from the exception handling vector table; the CPU then branches to that address and the program starts executing, thereby entering the program execution state.

(3) Program Execution State

In the program execution state, the CPU sequentially executes the program.

(4) Power-Down State

In the power-down state, the CPU stops operating to reduce power consumption. The SLEEP instruction places the CPU in sleep mode, software standby mode, or deep standby mode.

Section 3 Floating-Point Unit (FPU)

3.1 Features

The FPU has the following features.

- Conforms to IEEE754 standard
- 16 single-precision floating-point registers (can also be referenced as eight double-precision registers)
- Two rounding modes: Round to nearest and round to zero
- Denormalization modes: Flush to zero
- Five exception sources: Invalid operation, divide by zero, overflow, underflow, and inexact
- Comprehensive instructions: Single-precision, double-precision, and system control

3.2 Data Formats

3.2.1 Floating-Point Format

A floating-point number consists of the following three fields:

- $Sign(s)$
- Exponent (e)
- Fraction (f)

This LSI can handle single-precision and double-precision floating-point numbers, using the formats shown in figures 3.1 and 3.2.

Figure 3.1 Format of Single-Precision Floating-Point Number

63	\sim			

Figure 3.2 Format of Double-Precision Floating-Point Number

The exponent is expressed in biased form, as follows:

 $e = E + bias$

The range of unbiased exponent E is $E_{min} - 1$ to $E_{max} + 1$. The two values $E_{min} - 1$ and $E_{max} + 1$ are distinguished as follows. $E_{min} - 1$ indicates zero (both positive and negative sign) and a denormalized number, and $E_{\text{max}} + 1$ indicates positive or negative infinity or a non-number (NaN). Table 3.1 shows E_{min} and E_{max} values.

Floating-point number value v is determined as follows:

If $E = E_{max} + 1$ and $f \ne 0$, v is a non-number (NaN) irrespective of sign s If $E = E_{max} + 1$ and $f = 0$, $v = (-1)^s$ (infinity) [positive or negative infinity] If $E_{\text{min}} \le E \le E_{\text{max}}$, $v = (-1)^{s} 2^{E} (1.f)$ [normalized number] If $E = E_{min} - 1$ and $f \neq 0$, $v = (-1)^{s} 2^{Emin} (0.f)$ [denormalized number] If $E = E_{min} - 1$ and $f = 0$, $v = (-1)^{s}0$ [positive or negative zero]

Table 3.2 shows the ranges of the various numbers in hexadecimal notation.

Table 3.2 Floating-Point Ranges

3.2.2 Non-Numbers (NaN)

Figure 3.3 shows the bit pattern of a non-number (NaN). A value is NaN in the following case:

- Sign bit: Don't care
- Exponent field: All bits are 1
- Fraction field: At least one bit is 1

The NaN is a signaling NaN (sNaN) if the MSB of the fraction field is 1, and a quiet NaN (qNaN) if the MSB is 0.

An sNaN is input in an operation, except copy, FABS, and FNEG, that generates a floating-point value.

- When the EN.V bit in FPSCR is 0, the operation result (output) is a qNaN.
- When the EN.V bit in FPSCR is 1, an invalid operation exception will generate FPU exception processing. In this case, the contents of the operation destination register are unchanged.

If a qNaN is input in an operation that generates a floating-point value, and an sNaN has not been input in that operation, the output will always be a qNaN irrespective of the setting of the EN.V bit in FPSCR. An exception will not be generated in this case.

The qNAN values as operation results are as follows:

- Single-precision qNaN: H'7FBF FFFF
- Double-precision qNaN: H'7FF7 FFFF FFFF FFFF

See the individual instruction descriptions for details of floating-point operations when a nonnumber (NaN) is input.

3.2.3 Denormalized Numbers

For a denormalized number floating-point value, the exponent field is expressed as 0, and the fraction field as a non-zero value.

In the SH2A-FPU, the DN bit in the status register FPSCR is always set to 1, therefore a denormalized number (source operand or operation result) is always flushed to 0 in a floatingpoint operation that generates a value (an operation other than copy, FNEG, or FABS).

When the DN bit in FPSCR is 0, a denormalized number (source operand or operation result) is processed as it is. See the individual instruction descriptions for details of floating-point operations when a denormalized number is input.

3.3 Register Descriptions

3.3.1 Floating-Point Registers

Figure 3.4 shows the floating-point register configuration. There are sixteen 32-bit floating-point registers FPR0 to FPR15, referenced by specifying FR0 to FR15, DR0/2/4/6/8/10/12/14. The correspondence between FRPn and the reference name is determined by the PR and SZ bits in FPSCR. Refer figure 3.4.

- 1. Floating-point registers, FPRi (16 registers) FPR0 to FPR15
- 2. Single-precision floating-point registers, FRi (16 registers) FR0 to FR15 indicate FPR0 to FPR15
- 3. Double-precision floating-point registers or single-precision floating-point vector registers in pairs, DRi (8 registers)

A DR register comprises two FR registers.

 $DR0 = \{FR0, FR1\}$, $DR2 = \{FR2, FR3\}$, $DR4 = \{FR4, FR5\}$, $DR6 = \{FR6, FR7\}$, $DR8 = {FR8, FR9}, DR10 = {FR10, FR11}, DR12 = {FR12, FR13}, DR14 = {FR14, FR15}$

	Reference name		Register name
Transfer instruction case: Operation instruction case:	$FPSCR.SZ = 0$ $FPSCR.SZ = 1$ $FPSCR.PR = 0$ $FPSCR.PR = 1$		
	FR ₀ FR ₁	DR ₀	FPR ₀ FPR ₁
	FR ₂ FR ₃	DR ₂	FPR ₂ FPR ₃
	FR4 FR ₅	DR4	FPR4 FPR ₅
	FR ₆ FR ₇	DR ₆	FPR6 FPR7
	FR ₈ FR ₉	DR8	FPR8 FPR ₉
	FR10 FR11	DR10	FPR ₁₀ FPR11
	FR12 FR13	DR12	FPR ₁₂ FPR ₁₃
	FR14 FR15	DR14	FPR ₁₄ FPR ₁₅

Figure 3.4 Floating-Point Registers

3.3.2 Floating-Point Status/Control Register (FPSCR)

FPSCR is a 32-bit register that controls floating-point instructions, sets FPU exceptions, and selects the rounding mode.

Table 3.3 Bit Allocation for FPU Exception Handling

Note: No FPU error occurs in the SH2A-FPU.

3.3.3 Floating-Point Communication Register (FPUL)

Information is transferred between the FPU and CPU via FPUL. FPUL is a 32-bit system register that is accessed from the CPU side by means of LDS and STS instructions. For example, to convert the integer stored in general register R1 to a single-precision floating-point number, the processing flow is as follows:

 $R1 \rightarrow (LDS$ instruction) \rightarrow FPUL \rightarrow (single-precision FLOAT instruction) \rightarrow FR1

3.4 Rounding

In a floating-point instruction, rounding is performed when generating the final operation result from the intermediate result. Therefore, the result of combination instructions such as FMAC will differ from the result when using a basic instruction such as FADD, FSUB, or FMUL. Rounding is performed once in FMAC, but twice in FADD, FSUB, and FMUL.

Which of the two rounding methods is to be used is determined by the RM bits in FPSCR.

 $FPSCR.RM[1:0] = 00$: Round to Nearest $FPSCR.RM[1:0] = 01$: Round to Zero

(1) Round to Nearest

The operation result is rounded to the nearest expressible value. If there are two nearest expressible values, the one with an LSB of 0 is selected.

If the unrounded value is $2^{Emax} (2 - 2^{-P})$ or more, the result will be infinity with the same sign as the unrounded value. The values of Emax and P, respectively, are 127 and 24 for single-precision, and 1023 and 53 for double-precision.

(2) Round to Zero

The digits below the round bit of the unrounded value are discarded.

If the unrounded value is larger than the maximum expressible absolute value, the value will become the maximum expressible absolute value.

3.5 FPU Exceptions

3.5.1 FPU Exception Sources

FPU exceptions may occur on floating-point operation instruction and the exception sources are as follows:

- FPU error (E) : When FPSCR.DN = 0 and a denormalized number is input (No chance to occur in the SH2A-FPU)
- Invalid operation (V): In case of an invalid operation, such as NaN input
- Division by zero (Z): Division with a zero divisor
- Overflow (O): When the operation result overflows
- Underflow (U): When the operation result underflows
- Inexact exception (I): When overflow, underflow, or rounding occurs

The FPU exception cause field in FPSCR contains bits corresponding to all of above sources E, V, Z, O, U, and I, and the FPU exception flag and enable fields in FPSCR contain bits corresponding to sources V, Z, O, U, and I, but not E. Thus, FPU errors cannot be disabled.

When an FPU exception occurs, the corresponding bit in the FPU exception cause field is set to 1, and 1 is added to the corresponding bit in the FPU exception flag field. When an FPU exception does not occur, the corresponding bit in the FPU exception cause field is cleared to 0, but the corresponding bit in the FPU exception flag field remains unchanged.

3.5.2 FPU Exception Handling

FPU exception handling is initiated in the following cases:

- FPU error (E): FPSCR.DN = 0 and a denormalized number is input (No chance to occur in the SH2A-FPU)
- Invalid operation (V): FPSCR. Enable. $V = 1$ and invalid operation
- Division by zero (Z): FPSCR. Enable. $Z = 1$ and division with a zero divisor
- Overflow (O): FPSCR. Enable. $O = 1$ and instruction with possibility of operation result overflow
- Underflow (U): FPSCR. Enable. $U = 1$ and instruction with possibility of operation result underflow
- Inexact exception (I): FPSCR. Enable. $I = 1$ and instruction with possibility of inexact operation result

These possibilities of each exceptional handling on floating-point operation are shown in the individual instruction descriptions. All exception events that originate in the floating-point operation are assigned as the same FPU exceptional handling event. The meaning of an exception generated by floating-point operation is determined by software by reading from FPSCR and interpreting the information it contains. Also, the destination register is not changed when FPU exception handling operation occurs.

Except for the above, the FPU disables exception handling. In every processing, the bit corresponding to source V, Z, O, U, or I is set to 1, and a default value is generated as the operation result.

- Invalid operation (V): qNaN is generated as the result.
- Division by zero (Z) : Infinity with the same sign as the unrounded value is generated.
- Overflow (O):

When rounding mode $= RZ$, the maximum normalized number, with the same sign as the unrounded value, is generated.

When rounding mode = RN, infinity with the same sign as the unrounded value is generated.

- Underflow (U): Zero with the same sign as the unrounded value is generated.
	- Inexact exception (I): An inexact result is generated.
Section 4 Clock Pulse Generator (CPG)

This LSI has a clock pulse generator (CPG) that generates a CPU clock $(I\phi)$, a peripheral clock (Pφ), and a bus clock (Bφ). The CPG consists of a crystal oscillator, PLL circuits, and divider circuits.

4.1 Features

• Three clock operating modes

The mode is selected from among the three clock operating modes by the selection of the following three conditions: the frequency-divisor in use, whether the PLLs are on or off, and whether the internal crystal resonator or the input on the external clock-signal line is used.

- Three clocks generated independently A CPU clock $(I\phi)$ for the CPU and cache; a peripheral clock $(P\phi)$ for the on-chip peripheral modules; a bus clock ($B\phi = CKIO$) for the external bus interface.
- Frequency change function

CPU and peripheral clock frequencies can be changed independently using the PLL (phase locked loop) circuits and divider circuits within the CPG. Frequencies are changed by software using frequency control register (FRQCR) settings.

• Power-down mode control

The clock can be stopped by sleep mode, software standby mode, and deep standby mode. Specific modules can also be stopped using the module standby function. For details on clock control in the power-down modes, see section 27, Power-Down Modes.

Figure 4.1 shows a block diagram of the clock pulse generator.

Figure 4.1 Block Diagram of Clock Pulse Generator

The clock pulse generator blocks function as follows:

(1) PLL Circuit 1

PLL circuit 1 multiplies the input clock frequency from the CKIO pin by 1, 2, 3, 4, 6, or 8. The multiplication rate is set by the frequency control register. When this is done, the phase of the rising edge of the bus clock is controlled so that it will agree with the phase of the rising edge of the CKIO pin.

(2) PLL Circuit 2

PLL circuit 2 multiplies the input clock frequency from the crystal oscillator or EXTAL pin by 2 or 4. The multiplication rate is fixed according to the clock operating mode. The clock operating mode is specified by the MD_CLK1 and MD_CLK0 pins. For details on the clock operating mode, see table 4.2.

Note that the settings of these pins cannot be changed during operation. If changed, the operation of this LSI cannot be guaranteed.

(3) Crystal Oscillator

The crystal oscillator is an oscillation circuit in which a crystal resonator is connected to the XTAL pin or EXTAL pin. This can be used according to the clock operating mode.

(4) Divider

Divider generates a clock signal at the operating frequency used by the CPU or peripheral clock. The operating frequency can be 1, 1/2, 1/3, 1/4, 1/6, 1/8, or 1/12 times the output frequency of PLL circuit 1, as long as it stays at or above the clock frequency of the CKIO pin. The division ratio is set in the frequency control register (FRQCR).

(5) Clock Frequency Control Circuit

The clock frequency control circuit controls the clock frequency using the MD_CLK1 and MD_CLK0 pins and the frequency control register (FRQCR).

(6) Standby Control Circuit

The standby control circuit controls the states of the clock pulse generator and other modules during clock switching, or in sleep, software, and deep standby mode.

(7) Frequency Control Register (FRQCR)

 The frequency control register (FRQCR) has control bits assigned for the following functions: clock output/non-output from the CKIO pin during software standby mode, the frequency multiplication ratio of PLL circuit 1, and the frequency division ratio of the CPU clock and the peripheral clock (Pφ).

(8) Standby Control Register

The standby control register has bits for controlling the power-down modes. See section 27, Power-Down Modes, for more information.

4.2 Input/Output Pins

Table 4.1 lists the clock pulse generator pins and their functions.

Table 4.1 Pin Configuration and Functions of the Clock Pulse Generator

4.3 Clock Operating Modes

Table 4.2 shows the relationship between the combinations of the mode control pins (MD_CLK1 and MD CLK0) and the clock operating modes. Table 4.2 shows the usable frequency ranges in the clock operating modes.

Table 4.2 Clock Operating Modes

\bullet Mode 0

The frequency of the signal received from the EXTAL pin or crystal resonator is quadrupled by the PLL circuit 2 before it is supplied to the LSI as the clock signal. This enables to use the external clock of lower frequency. Either a crystal resonator with a frequency in the range from 10 to 15 MHz or an external signal in the same frequency range input on the EXTAL pin may be used. The frequency range of CKIO is from 40 to 60 MHz.

 \bullet Mode 2

The frequency of the signal received from the EXTAL pin or crystal resonator is doubled by the PLL circuit 2 before it is supplied to the LSI as the clock signal. This enables to use the external clock of lower frequency. An external signal with a frequency in the range from 10 to 30 MHz or a crystal resonator with 10 to 20 MHz may be used. The frequency range of CKIO is from 20 to 60 MHz.

Mode 3

In mode 3, the CKIO pin functions as an input pin and draws an external clock signal. The PLL circuit 1 shapes its waveform and the setting of the frequency control register multiplies its frequency before the clock enters the LSI. Frequency between 20 to 60 MHz can be input to the CKIO pin. For reduced current and hence power consumption, pull up the EXTAL pin and open the XTAL pin when the LSI is used in mode 3.

Table 4.3 Relationship between Clock Operating Mode and Frequency Range

Notes: 1. The ratio of clock frequencies, where the input clock frequency is assumed to be 1.

 2. In modes 0 and 2, the frequency of the clock input from the EXTAL pin or the frequency of the crystal resonator. In mode 3, the frequency of the clock input from the CKIO pin.

 3. Use an internal clock (Iφ) frequency of 120 MHz or lower for the regular specifications and 100 MHz or lower for the wide-range specifications. Use a CKIO pin or bus clock (B ϕ) frequency of 60 MHz or lower. P ϕ must be from 5 through 40 MHz.

Caution: Do not use this LSI for frequency settings other than those in table 4.3.

4.4 Register Descriptions

The clock pulse generator has the following registers.

Table 4.4 Register Configuration

4.4.1 Frequency Control Register (FRQCR)

FRQCR is a 16-bit readable/writable register used to specify whether a clock is output from the CKIO pin in software standby mode, the frequency multiplication ratio of PLL circuit 1, and the frequency division ratio of the CPU clock and peripheral clock $(P\phi)$. Only word access can be used on FRQCR.

FRQCR is initialized to H'1003 only by a power-on reset or in deep standby mode. FRQCR retains its previous value by a manual reset or in software standby mode. The previous value is also retained when an internal reset is triggered by an overflow of the WDT.

4.4.2 CKIO Control Register (CKIOCR)

CKIOCR is an 8-bit readable/writable register used to control output of the CKIO pin. When this LSI is started in clock operating mode 3, writing 1 to this register is invalid.

When this LSI is started in clock operating mode 3, CKIOCR is initialized to H'00 by a power-on reset caused by the $\overline{\rm RES}$ pin or in deep standby mode. When this LSI is started in clock operating mode 0 or 2, CKIOCR is initialized to H⁰¹ by a power-on reset caused by the \overline{RES} pin or in deep standby mode. This register is not initialized by an internal reset triggered by an overflow of the WDT, a manual reset, in sleep mode, or in software standby mode.

Note: * The initial value depends on the clock operating mode of the LSI.

4.5 Changing the Frequency

The frequency of the CPU clock $(I\phi)$ and peripheral clock $(P\phi)$ can be changed either by changing the multiplication rate of PLL circuit 1 or by changing the division rates of divider. All of these are controlled by software through the frequency control register (FRQCR). The methods are described below.

4.5.1 Changing the Multiplication Rate

A PLL settling time is required when the multiplication rate of PLL circuit 1 is changed. The onchip WDT counts the settling time.

- 1. In the initial state, the multiplication rate of PLL circuit 1 is 1 time.
- 2. Set a value that will become the specified oscillation settling time in the WDT and stop the WDT. The following must be set: $WTCSR.TME = 0$: WDT stops WTCSR.CKS[2:0]: Division ratio of WDT count clock WTCNT counter: Initial counter value
- 3. Set the desired value in the STC[2:0] bits. The division ratio can also be set in the IFC[2:0] and PFC[2:0] bits.
- 4. This LSI pauses temporarily and the WDT starts incrementing. The CPU and peripheral clocks both stop and the WDT is supplied with the clock. The clock will continue to be output at the CKIO pin. This state is the same as software standby mode. Whether or not registers are initialized depends on the module. For details, see section 27, Power-Down Modes.
- 5. Supply of the clock that has been set begins at WDT count overflow, and this LSI begins operating again. The WDT stops after it overflows.

4.5.2 Changing the Division Ratio

Counting by the WDT does not proceed if the frequency divisor is changed but the multiplier is not.

- 1. In the initial state, $IFC[2:0] = B'000$ and $PFC[2:0] = B'011$.
- 2. Set the desired value in the IFC[2:0] and PFC[2:0] bits. The values that can be set are limited by the clock operating mode and the multiplication rate of PLL circuit 1. Note that if the wrong value is set, this LSI will malfunction.
- 3. After the register bits (IFC[2:0] and PFC[2:0]) have been set, the clock is supplied of the new division ratio.
- Note: When executing the SLEEP instruction after the frequency has been changed, be sure to read the frequency control register (FRQCR) three times before executing the SLEEP instruction.

4.6 Notes on Board Design

4.6.1 Note on Inputting External Clock

Figure 4.2 is an example of connecting the external clock input. When putting the XTAL pin in open state, make sure the parasitic capacitance is less than or equal to 10 pF. To stably input the external clock with enough PLL stabilizing time at power on or releasing the standby, wait longer than the oscillation stabilizing time.

Figure 4.2 Example of Connecting External Clock

For details on input conditions of the external clock, see section 31.3.1, Clock Timing.

4.6.2 Note on Using Crystal Resonator

Place the crystal resonator and capacitors CL1 and CL2 as close to the XTAL and EXTAL pins as possible. In addition, to minimize induction and thus obtain oscillation at the correct frequency, the capacitors to be attached to the resonator must be grounded to the same ground. Do not bring wiring patterns close to these components.

Figure 4.3 Note on Using Crystal Resonator

4.6.3 Note on Resonator

Since various characteristics related to the resonator are closely linked to the user's board design, thorough evaluation is necessary on the user's part, using the resonator connection examples shown in this section as a guide. As the parameters for the oscillation circuit will depend on the floating capacitance of the resonator and the user board, the parameters should be determined in consultation with the resonator manufacturer. The design must ensure that a voltage exceeding the maximum rating is not applied to the resonator pin.

4.6.4 Note on Using a PLL Oscillation Circuit

In the PLLVcc and PLLVss connection pattern for the PLL, signal lines from the board power supply pins must be as short as possible and pattern width must be as wide as possible to reduce inductive interference.

In clock operating mode 3, the EXTAL pin is pulled up and the XTAL pin is left open.

Since the analog power supply pins of the PLL are sensitive to the noise, the system may malfunction due to inductive interference at the other power supply pins. To prevent such malfunction, the analog power supply pin Vcc and the digital power supply pins VccR and PVcc should not supply the same resources on the board if at all possible.

Figure 4.4 Note on Using PLL Oscillation Circuit

4.6.5 Note on Changing the Multiplication Rate

If the multiplication rate is changed by the frequency control register (FRQCR) during transfer by the DMAC, the DMAC stops its operation without waiting for the completion of the transfer. Thus, the DMA transfer is not guaranteed. Therefore, when changing the multiplication rate with the frequency control register (FRQCR), wait for the completion of the DMA transfer or stop the DMA transfer to change the setting of the frequency control register (FRQCR).

Section 5 Exception Handling

5.1 Overview

5.1.1 Types of Exception Handling and Priority

Exception handling is started by sources, such as resets, address errors, bus errors, register bank errors, interrupts, and instructions. Table 5.1 shows their priorities. When several exception handling sources occur at once, they are processed according to the priority shown.

Type	Exception Handling	Priority
Reset	Power-on reset	High
	Manual reset	
Address error	CPU address error	
Bus error	Bus error	
Instructions	FPU exception	
	Integer division exception (division by zero)	
	Integer division exception (overflow)	
Register bank error	Bank underflow	
	Bank overflow	
Interrupts	NMI	
	User break	
	H-UDI	
	IRQ	
	PINT	Low

Table 5.1 Types of Exception Handling and Priority Order

- Notes: 1. Delayed branch instructions: JMP, JSR, BRA, BSR, RTS, RTE, BF/S, BT/S, BSRF, BRAF.
	- 2. Instructions that rewrite the PC: JMP, JSR, BRA, BSR, RTS, RTE, BT, BF, TRAPA, BF/S, BT/S, BSRF, BRAF, JSR/N, RTV/N.
	- 3. 32-bit instructions: BAND.B, BANDNOT.B, BCLR.B, BLD.B, BLDNOT.B, BOR.B, BORNOT.B, BSET.B, BST.B, BXOR.B, MOV.B@disp12, MOV.W@disp12, MOV.L@disp12, MOVI20, MOVI20S, MOVU.B, MOVU.W.

5.1.2 Exception Handling Operations

The exception handling sources are detected and begin processing according to the timing shown in table 5.2.

When exception handling starts, the CPU operates as follows:

(1) Exception Handling Triggered by Reset

The initial values of the program counter (PC) and stack pointer (SP) are fetched from the exception handling vector table (PC and SP are respectively the H'00000000 and H'00000004 addresses for power-on resets and the H'00000008 and H'0000000C addresses for manual resets). See section 5.1.3, Exception Handling Vector Table, for more information. The vector base register (VBR) is then initialized to H'00000000, the interrupt mask level bits (I3 to I0) of the status register (SR) are initialized to H'F (B'1111), and the BO and CS bits are initialized. The BN bit in IBNR of the interrupt controller (INTC) is also initialized to 0. FPSCR is initialized to H'00040001 by a power-on reset. The program begins running from the PC address fetched from the exception handling vector table.

(2) Exception Handling Triggered by Address Errors, Bus Errors, Register Bank Errors, Interrupts, and Instructions

SR and PC are saved to the stack indicated by R15. In the case of interrupt exception handling other than the NMI or user break, with usage of the register banks enabled, general registers R0 to R14, control register GBR, system registers MACH, MACL, and PR, and the vector number of the interrupt exception handling to be executed are saved to the register banks. In the case of exception handling due to an address error, bus error, register bank error, NMI interrupt, user break interrupt, or instruction, saving to a register bank is not performed. When saving is performed to all register banks, automatic saving to the stack is performed instead of register bank saving. In this case, an interrupt controller setting must have been made so that register bank overflow exceptions are not accepted (the BOVE bit in IBNR of the INTC is 0). If a setting to accept register bank overflow exceptions has been made (the BOVE bit in IBNR of the INTC is 1), register bank overflow exception will be generated. In the case of interrupt exception handling, the interrupt priority level is written to the I3 to I0 bits in SR. In the case of exception handling due to an address error or instruction, the I3 to I0 bits are not affected. The start address is then fetched from the exception handling vector table and the program begins running from that address.

5.1.3 Exception Handling Vector Table

Before exception handling begins running, the exception handling vector table must be set in memory. The exception handling vector table stores the start addresses of exception service routines. (The reset exception handling table holds the initial values of PC and SP.)

All exception sources are given different vector numbers and vector table address offsets, from which the vector table addresses are calculated. During exception handling, the start addresses of the exception service routines are fetched from the exception handling vector table, which is indicated by this vector table address.

Table 5.3 shows the vector numbers and vector table address offsets. Table 5.4 shows how vector table addresses are calculated.

Table 5.3 Exception Handling Vector Table

Note: * The vector numbers and vector table address offsets for each external interrupt and onchip peripheral module interrupt are given in table 6.4 in section 6, Interrupt Controller (INTC).

Table 5.4 Calculating Exception Handling Vector Table Addresses

Notes: 1. Vector table address offset: See table 5.3.

2. Vector number: See table 5.3.

5.2 Resets

5.2.1 Input/Output Pins

Table 5.5 shows the configuration of pins relating to the resets.

Table 5.5 Pin Configuration

5.2.2 Types of Reset

A reset is the highest-priority exception handling source. There are two kinds of resets, power-on and manual. As shown in table 5.6, the CPU state is initialized by both a power-on reset and a manual reset. The FPU state is initialized by a power-on reset, but not by a manual reset. On-chip peripheral module registers except a few registers are initialized by a power-on reset, but not by a manual reset.

Table 5.6 Reset States

2. The BN bit in IBNR of the INTC is initialized.

5.2.3 Power-On Reset

(1) Power-On Reset by Means of RES **Pin**

When the RES pin is driven low, this LSI enters the power-on reset state. To reliably reset this LSI, the RES pin should be kept at the low level for the duration of the oscillation settling time at power-on or when in software standby mode (when the clock is halted), or at least 20-tcyc when the clock is running. In the power-on reset state, the internal state of the CPU and all the on-chip peripheral module registers are initialized. See appendix A, Pin States, for the status of individual pins during the power-on reset state.

In the power-on reset state, power-on reset exception handling starts when the RES pin is first driven low for a fixed period and then returned to high. The CPU operates as follows:

- 1. The initial value (execution start address) of the program counter (PC) is fetched from the exception handling vector table.
- 2. The initial value of the stack pointer (SP) is fetched from the exception handling vector table.
- 3. The vector base register (VBR) is cleared to H'00000000, the interrupt mask level bits (I3 to I0) of the status register (SR) are initialized to H'F (B'1111), and the BO and CS bits are initialized to 0. The BN bit in IBNR of the INTC is also initialized to 0. FPSCR is initialized to H'00040001.
- 4. The values fetched from the exception handling vector table are set in the PC and SP, and the program begins executing.

Be certain to always perform power-on reset processing when turning the system power on.

(2) Power-On Reset by Means of H-UDI Reset Assert Command

When the H-UDI reset assert command is set, this LSI enters the power-on reset state. Power-on reset by means of an H-UDI reset assert command is equivalent to power-on reset by means of the RES pin. Setting the H-UDI reset negate command cancels the power-on reset state. The time required between an H-UDI reset assert command and H-UDI reset negate command is the same as the time to keep the RES pin low to initiate a power-on reset. In the power-on reset state generated by an H-UDI reset assert command, setting the H-UDI reset negate command starts power-on reset exception handling. The CPU operates in the same way as when a power-on reset was caused by the RES pin.

(3) Power-On Reset Initiated by WDT

When a setting is made for a power-on reset to be generated in the WDT's watchdog timer mode, and WTCNT of the WDT overflows, this LSI enters the power-on reset state.

In this case, WRCSR of the WDT and FRQCR of the CPG are not initialized by the reset signal generated by the WDT.

If a reset caused by the RES pin or the H-UDI reset assert command occurs simultaneously with a reset caused by WDT overflow, the reset caused by the RES pin or the H-UDI reset assert command has priority, and the WOVF bit in WRCSR is cleared to 0. When power-on reset exception processing is started by the WDT, the CPU operates in the same way as when a poweron reset was caused by the RES pin.

5.2.4 Manual Reset

(1) Manual Reset by Means of MRES **Pin**

When the MRES pin is driven low, this LSI enters the manual reset state. To reset this LSI without fail, the MRES pin should be kept at the low level for at least 20-tcyc. In the manual reset state, the CPU's internal state is initialized, but all the on-chip peripheral module registers are not initialized. In the manual reset state, manual reset exception handling starts when the MRES pin is first driven low for a fixed period and then returned to high. The CPU operates as follows:

- 1. The initial value (execution start address) of the program counter (PC) is fetched from the exception handling vector table.
- 2. The initial value of the stack pointer (SP) is fetched from the exception handling vector table.
- 3. The vector base register (VBR) is cleared to H'00000000, the interrupt mask level bits (I3 to I0) of the status register (SR) are initialized to H'F (B'1111), and the BO and CS bits are initialized. The BN bit in IBNR of the INTC is also initialized to 0.
- 4. The values fetched from the exception handling vector table are set in the PC and SP, and the program begins executing.

(2) Manual Reset Initiated by WDT

When a setting is made for a manual reset to be generated in the WDT's watchdog timer mode, and WTCNT of the WDT overflows, this LSI enters the manual reset state.

When manual reset exception processing is started by the WDT, the CPU operates in the same way as when a manual reset was caused by the MRES pin.

(3) Notes at a Manual Reset

When a manual reset is generated, the bus cycle is retained. Thus, manual reset exception handling will be deferred until the CPU acquires the bus mastership. The CPU and the BN bit in IBNR of the INTC are initialized by a manual reset. The FPU and other modules are not initialized.

5.3 Address Errors

5.3.1 Address Error Sources

Address errors occur when instructions are fetched or data read or written, as shown in table 5.7.

Table 5.7 Bus Cycles and Address Errors

Notes: 1. For details on cache address array space, see section 8, Cache.

2. For details on peripheral module space, see section 9, Bus State Controller (BSC).

5.3.2 Address Error Exception Handling

When an address error occurs, address error exception handling starts after the bus cycle in which the address error occurred ends* and execution of the instruction being executed completes. The CPU operates as follows.

- 1. The exception service routine start address which corresponds to the address error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.
- Note: * In the case of address error related to data read/write. In the case of address error related to instruction fetch, if the bus cycle in which the address error occurred doesn't end until the entire three above-mentioned operations end, the CPU will start address error exception handling again until the bus cycle in which the address error occurred ends.

5.4 Bus Error

5.4.1 Bus Error Generation Source

In bus monitor, notification of bus error occurrence to the CPU can be set. The notification is generated when incorrect address access or bus timeout is detected. For details, see section 10, Bus Monitor.

5.4.2 Bus Error Exception Handling

When a bus error occurs, bus error exception handling starts after the bus cycle in which the bus error occurred ends and execution of the instruction being executed completes. The CPU operates as follows.

- 1. The exception service routine start address which corresponds to the bus error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

5.5 Register Bank Errors

5.5.1 Register Bank Error Sources

(1) Bank Overflow

In the state where saving has already been performed to all register bank areas, bank overflow occurs when acceptance of register bank overflow exception has been set by the interrupt controller (the BOVE bit in IBNR of the INTC is set to 1) and an interrupt that uses a register bank has occurred and been accepted by the CPU.

(2) Bank Underflow

Bank underflow occurs when an attempt is made to execute a RESBANK instruction while saving has not been performed to register banks.

5.5.2 Register Bank Error Exception Handling

When a register bank error occurs, register bank error exception handling starts. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the register bank error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction for a bank overflow, and the start address of the executed RESBANK instruction for a bank underflow.

To prevent multiple interrupts from occurring at a bank overflow, the interrupt priority level that caused the bank overflow is written to the interrupt mask level bits (I3 to I0) of the status register (SR).

4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

5.6 Interrupts

5.6.1 Interrupt Sources

Table 5.8 shows the sources that start up interrupt exception handling. These are divided into NMI, user breaks, H-UDI, IRQ, PINT, and on-chip peripheral modules.

Table 5.8 Interrupt Sources

Each interrupt source is allocated a different vector number and vector table offset. See table 6.4 in section 6, Interrupt Controller (INTC), for more information on vector numbers and vector table address offsets.

5.6.2 Interrupt Priority Level

The interrupt priority order is predetermined. When multiple interrupts occur simultaneously (overlap), the interrupt controller (INTC) determines their relative priorities and starts processing according to the results.

The priority order of interrupts is expressed as priority levels 0 to 16, with priority 0 the lowest and priority 16 the highest. The NMI interrupt has priority 16 and cannot be masked, so it is always accepted. The user break interrupt and H-UDI interrupt priority level is 15. Priority levels of IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts can be set freely using the interrupt priority registers 01, 02, and 05 to 16 (IPR01, IPR02, and IPR05 to IPR16) of the INTC as shown in table 5.9. The priority levels that can be set are 0 to 15. Level 16 cannot be set. See section 6.3.1, Interrupt Priority Registers 01, 02, 05 to 16 (IPR01, IPR02, IPR05 to IPR16), for details of IPR01, IPR02, and IPR05 to IPR16.

Table 5.9 Interrupt Priority Order

5.6.3 Interrupt Exception Handling

When an interrupt occurs, its priority level is ascertained by the interrupt controller (INTC). NMI is always accepted, but other interrupts are only accepted if they have a priority level higher than the priority level set in the interrupt mask level bits (I3 to I0) of the status register (SR).

When an interrupt is accepted, interrupt exception handling begins. In interrupt exception handling, the CPU fetches the exception service routine start address which corresponds to the accepted interrupt from the exception handling vector table, and saves SR and the program counter (PC) to the stack. In the case of interrupt exception handling other than the NMI or user break, with usage of the register banks enabled, general registers R0 to R14, control register GBR, system registers MACH, MACL, and PR, and the vector number of the interrupt exception handling to be executed are saved in the register banks. In the case of exception handling due to an address error, bus error, NMI interrupt, user break interrupt, or instruction, saving is not performed to the register banks. If saving has been performed to all register banks (0 to 14), automatic saving to the stack is performed instead of register bank saving. In this case, an interrupt controller setting must have been made so that register bank overflow exceptions are not accepted (the BOVE bit in IBNR of the INTC is 0). If the interrupt controller is set to accept register bank overflow exceptions (the BOVE bit in IBNR of INTC is set to 1), a register bank overflow exception will occur. Next, the priority level value of the accepted interrupt is written to the I3 to I0 bits in SR. For NMI, however, the priority level is 16, but the value set in the I3 to I0 bits is H[']F (level 15). Then, after jumping to the start address of the interrupt exception service routine fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch. See section 6.6, Operation, for further details of interrupt exception handling.

5.7 Exceptions Triggered by Instructions

5.7.1 Types of Exceptions Triggered by Instructions

Exception handling can be triggered by trap instructions, general illegal instructions, slot illegal instructions, integer division exceptions, and FPU exceptions, as shown in table 5.10.

5.7.2 Trap Instructions

When a TRAPA instruction is executed, trap instruction exception handling starts. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the vector number specified in the TRAPA instruction is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the TRAPA instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

5.7.3 Slot Illegal Instructions

An instruction placed immediately after a delayed branch instruction is said to be placed in a delay slot. When the instruction placed in the delay slot is undefined code, an instruction that rewrites the PC, a 32-bit instruction, an RESBANK instruction, a DIVS instruction, or a DIVU instruction, slot illegal exception handling starts when such kind of instruction is decoded. The CPU operates as follows:

- 1. The exception service routine start address is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the jump address of the delayed branch instruction immediately before the undefined code, the instruction that rewrites the PC, the 32-bit instruction, the RESBANK instruction, the DIVS instruction, or the DIVU instruction.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

5.7.4 General Illegal Instructions

When undefined code placed anywhere other than immediately after a delayed branch instruction (delay slot) is decoded, general illegal instruction exception handling starts. The CPU handles general illegal instructions in the same way as slot illegal instructions. Unlike processing of slot illegal instructions, however, the program counter value stored is the start address of the undefined code.

5.7.5 Integer Division Exceptions

When an integer division instruction performs division by zero or the result of integer division overflows, integer division instruction exception handling starts. The instructions that may become the source of division-by-zero exception are DIVU and DIVS. The only source instruction of overflow exception is DIVS, and overflow exception occurs only when the negative maximum value is divided by −1. The CPU operates as follows:

- 1. The exception service routine start address which corresponds to the integer division exception that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the integer division instruction at which the exception occurred.
- 4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

5.7.6 FPU Exceptions

An FPU exception handling is generated when the V, Z, O, U or I bit in the FPU exception enable field (Enable) of the floating point status/control register (FPSCR) is set. This indicates the occurrence of an invalid operation exception defined by the IEEE standard 754, a division-by-zero exception, overflow (in the case of an instruction for which this is possible), underflow (in the case of an instruction for which this is possible), or inexact exception (in the case of an instruction for which this is possible).

The floating-point operation instructions that may cause generation of an FPU exception handling are FADD, FSUB, FMUL, FDIV, FMAC, FCMP/EQ, FCMP/GT, FLOAT, FTRC, FCNVDS, FCNVSD, and FSQRT.

An FPU exception handling is generated only when the corresponding FPU exception enable bit (Enable) is set. When the FPU detects an exception source by a floating-point operation, FPU operation is halted and FPU exception handling generation is reported to the CPU. When exception handling is started, the CPU operations are as follows.

- 1. The start address of the exception service routine which corresponding to the FPU exception handling that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction.

4. After jumping to the address fetched from the exception handling vector table, program execution starts. The jump that occurs is not a delayed branch.

The FPU exception flag field (Flag) of FPSCR is always updated regardless of whether or not an FPU exception handling has been accepted, and remains set until explicitly cleared by the user through an instruction. The FPU exception source field (Cause) of FPSCR changes each time a floating-point operation instruction is executed.

When the V bit in the FPU exception enable field (Enable) of FPSCR is set and the QIS bit in FPSCR is also set, FPU exception handling is generated when qNAN or $\pm \infty$ is input to a floating point operation instruction source.

5.8 When Exception Sources Are Not Accepted

When an address error, bus error, FPU exception, register bank error (overflow), or interrupt is generated immediately after a delayed branch instruction, it is sometimes not accepted immediately but stored instead, as shown in table 5.11. When this happens, it will be accepted when an instruction that can accept the exception is decoded.

Table 5.11 Exception Source Generation Immediately after Delayed Branch Instruction

Note: * Delayed branch instructions: JMP, JSR, BRA, BSR, RTS, RTE, BF/S, BT/S, BSRF, BRAF.

5.9 Stack Status after Exception Handling Ends

The status of the stack after exception handling ends is as shown in table 5.12.

Table 5.12 Stack Status After Exception Handling Ends

5.10 Usage Notes

5.10.1 Value of Stack Pointer (SP)

The value of the stack pointer must always be a multiple of four. If it is not, an address error will occur when the stack is accessed during exception handling.

5.10.2 Value of Vector Base Register (VBR)

The value of the vector base register must always be a multiple of four. If it is not, an address error will occur when the stack is accessed during exception handling.

5.10.3 Address Errors Caused by Stacking of Address Error Exception Handling

When the stack pointer is not a multiple of four, an address error will occur during stacking of the exception handling (interrupts, etc.) and address error exception handling will start up as soon as the first exception handling is ended. Address errors will then also occur in the stacking for this address error exception handling. To ensure that address error exception handling does not go into an endless loop, no address errors are accepted at that point. This allows program control to be shifted to the address error exception service routine and enables error processing.

When an address error occurs during exception handling stacking, the stacking bus cycle (write) is executed. During the stacking of the status register (SR) and program counter (PC), the SP is decremented by 4 for both, so the value of SP will not be a multiple of four after the stacking either. The address value output during stacking is the SP value, so the address where the error occurred is itself output. This means the write data stacked will be undefined.

Section 6 Interrupt Controller (INTC)

The interrupt controller (INTC) ascertains the priority of interrupt sources and controls interrupt requests to the CPU. The INTC registers set the order of priority of each interrupt, allowing the user to process interrupt requests according to the user-set priority.

6.1 Features

• 16 levels of interrupt priority can be set

By setting the 14 interrupt priority registers, the priorities of the IRQ, PINT, and on-chip peripheral module interrupts can be set to one of 16 levels for each source.

• NMI noise canceller function

This controller provides an NMI input level bit that indicates the NMI pin state. The interrupt exception service routine can verify the pin state by reading this bit and use the information to implement a noise canceling function.

• Register banks

This LSI has register banks that enable register saving and restoration required in the interrupt processing to be performed at high speed.

Figure 6.1 shows a block diagram of the INTC.

Figure 6.1 Block Diagram of INTC

6.2 Input/Output Pins

Table 6.1 shows the pin configuration of the INTC.

Table 6.1 Pin Configuration

6.3 Register Descriptions

The INTC has the following registers. These registers are used to set the interrupt priorities and control detection of the external interrupt input signal.

Table 6.2 Register Configuration

Notes: 1. When the NMI pin is high, becomes H'8000; when low, becomes H'0000.

2. Only 0 can be written after reading 1, to clear the flag.

6.3.1 Interrupt Priority Registers 01, 02, 05 to 16 (IPR01, IPR02, IPR05 to IPR16)

IPR01, IPR02, and IPR05 to IPR16 are 16-bit readable/writable registers in which priority levels from 0 to 15 are set for IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts. Table 6.3 shows the correspondence between the interrupt request sources and the bits in IPR01, IPR02, and IPR05 to IPR16.

Table 6.3 Interrupt Request Sources and IPR01, IPR02, and IPR05 to IPR16

As shown in table 6.3, by setting the 4-bit groups (bits 15 to 12, bits 11 to 8, bits 7 to 4, and bits 3 to 0) with values from H'0 (0000) to H'F (1111), the priority of each corresponding interrupt is set. Setting of H'0 means priority level 0 (the lowest level) and H'F means priority level 15 (the highest level).

IPR01, IPR02, and IPR05 to IPR16 are initialized to H'0000 by a power-on reset or in deep standby mode.

6.3.2 Interrupt Control Register 0 (ICR0)

ICR0 is a 16-bit register that sets the input signal detection mode for the external interrupt input pin NMI, and indicates the input level at the NMI pin. ICR0 is initialized by a power-on reset or in deep standby mode.

Note: * 1 when the NMI pin is high, and 0 when the NMI pin is low.

6.3.3 Interrupt Control Register 1 (ICR1)

ICR1 is a 16-bit register that specifies the detection mode for external interrupt input pins IRQ7 to IRQ0 individually: low level, falling edge, rising edge, or both edges. ICR1 is initialized by a power-on reset or in deep standby mode.

[Legend]

 $n = 7$ to 0

6.3.4 Interrupt Control Register 2 (ICR2)

ICR2 is a 16-bit register that specifies the detection mode for external interrupt input pins PINT7 to PINT0 individually: low level or high level. ICR2 is initialized by a power-on reset or in deep standby mode.

 $n = 7$ to 0

6.3.5 IRQ Interrupt Request Register (IRQRR)

IRQRR is a 16-bit register that indicates interrupt requests from external input pins IRQ7 to IRQ0. If edge detection is set for the IRQ7 to IRQ0 interrupts, writing 0 to the IRQ7F to IRQ0F bits after reading IRQ7F to IRQ0F = 1 cancels the retained interrupts.

IRQRR is initialized by a power-on reset or in deep standby mode.

Note: * Only 0 can be written to clear the flag after 1 is read.

[Legend] $n = 7$ to 0

6.3.6 PINT Interrupt Enable Register (PINTER)

PINTER is a 16-bit register that enables interrupt request inputs to external interrupt input pins PINT7 to PINT0. PINTER is initialized by a power-on reset or in deep standby mode.

[Legend]

 $n = 7$ to 0

6.3.7 PINT Interrupt Request Register (PIRR)

PIRR is a 16-bit register that indicates interrupt requests from external input pins PINT7 to PINT0. PIRR is initialized by a power-on reset or in deep standby mode.

[Legend]

 $n = 7$ to 0

6.3.8 Bank Control Register (IBCR)

IBCR is a 16-bit register that enables or disables use of register banks for each interrupt priority level. IBCR is initialized to H'0000 by a power-on reset or in deep standby mode.

6.3.9 Bank Number Register (IBNR)

IBNR is a 16-bit register that enables or disables use of register banks and register bank overflow exception. IBNR also indicates the bank number to which saving is performed next through the bits BN3 to BN0.

IBNR is initialized to H'0000 by a power-on reset or in deep standby mode.

Note: * Bits BN[3:0] are initialized at a manual reset.

6.3.10 DMA Transfer Request Enable Register 0 (DREQER0)

DMA transfer request enable register 0 (DREQER0) is an 8-bit readable/writable register that enables/disables the IIC3 DMA transfer requests, and enables/disables CPU interrupt requests.

DMA transfer request enable register 0 is initialized by a power-on reset or in deep standby mode.

6.3.11 DMA Transfer Request Enable Register 1 (DREQER1)

DMA transfer request enable register 1 (DREQER1) is an 8-bit readable/writable register that enables/disables the SCIF (channels 0 to 3) DMA transfer requests, and enables/disables CPU interrupt requests.

DMA transfer request enable register 1 is initialized by a power-on reset or in deep standby mode.

6.3.12 DMA Transfer Request Enable Register 2 (DREQER2)

DMA transfer request enable register 2 (DREQER2) is an 8-bit readable/writable register that enables/disables the SCIF (channels 4 to 7) DMA transfer requests, and enables/disables CPU interrupt requests.

DMA transfer request enable register 2 is initialized by a power-on reset or in deep standby mode.

6.3.13 DMA Transfer Request Enable Register 3 (DREQER3)

DMA transfer request enable register 3 (DREQER3) is an 8-bit readable/writable register that enables/disables the ADC, MTU2 (channels 0 to 4), and RCAN-ET (channels 0 and 1) DMA transfer requests, and enables/disables CPU interrupt requests.

DMA transfer request enable register 3 is initialized by a power-on reset or in deep standby mode.

6.4 Interrupt Sources

There are six types of interrupt sources: NMI, user break, H-UDI, IRQ, PINT, and on-chip peripheral modules. Each interrupt has a priority level (0 to 16), with 0 the lowest and 16 the highest. When set to level 0, that interrupt is masked at all times.

6.4.1 NMI Interrupt

The NMI interrupt has a priority level of 16 and is accepted at all times. NMI interrupt requests are edge-detected, and the NMI edge select bit (NMIE) in interrupt control register 0 (ICR0) selects whether the rising edge or falling edge is detected.

Though the priority level of the NMI interrupt is 16, the NMI interrupt exception handling sets the interrupt mask level bits (I3 to I0) in the status register (SR) to level 15.

6.4.2 User Break Interrupt

A user break interrupt which occurs when a break condition set in the user break controller (UBC) matches has a priority level of 15. The user break exception handling sets the I3 to I0 bits in SR to level 15. For user break interrupts, see section 7, User Break Controller (UBC).

6.4.3 H-UDI Interrupt

The user debugging interface (H-UDI) interrupt has a priority level of 15, and occurs at serial input of an H-UDI interrupt instruction. H-UDI interrupt requests are edge-detected and retained until they are accepted. The H-UDI exception handling sets the I3 to I0 bits in SR to level 15. For H-UDI interrupts, see section 28, User Debugging Interface (H-UDI).

6.4.4 IRQ Interrupts

IRQ interrupts are input from pins IRQ7 to IRQ0. As regard to the setting method of pins IRQ7 to IRQ0, see section 25, Pin Function Controller (PFC). For the IRQ interrupts, low-level, fallingedge, rising-edge, or both-edge detection can be selected individually for each pin by the IRQ sense select bits (IRQ71S to IRQ01S and IRQ70S to IRQ00S) in interrupt control register 1 (ICR1). The priority level can be set individually in a range from 0 to 15 for each pin by interrupt priority registers 01 and 02 (IPR01 and IPR02).

When using low-level sensing for IRQ interrupts, an interrupt request signal is sent to the INTC while the IRQ7 to IRQ0 pins are low. An interrupt request signal is stopped being sent to the INTC when the IRQ7 to IRQ0 pins are driven high. The status of the interrupt requests can be checked by reading the IRQ interrupt request bits (IRQ7F to IRQ0F) in the IRQ interrupt request register (IRQRR).

When using edge-sensing for IRQ interrupts, an interrupt request is detected due to change of the IRQ7 to IRQ0 pin states, and an interrupt request signal is sent to the INTC. The result of IRQ interrupt request detection is retained until that interrupt request is accepted. Whether IRQ interrupt requests have been detected or not can be checked by reading the IRQ7F to IRQ0F bits in IRQRR. Writing 0 to these bits after reading them as 1 clears the result of IRQ interrupt request detection.

The IRQ interrupt exception handling sets the I3 to I0 bits in SR to the priority level of the accepted IRQ interrupt.

When restoring from the service routine of IRQ interrupt exception handling, execute the RTE instruction after an interrupt request has been cleared in the IRQ interrupt request register (IRQRR).

6.4.5 PINT Interrupts

PINT interrupts are input from pins PINT7 to PINT0. As regard to the setting method of pins PINT7 to PINT0, see section 25, Pin Function Controller (PFC). Input of the interrupt requests is enabled by the PINT enable bits (PINT7E to PINT0E) in the PINT interrupt enable register (PINTER). For the PINT7 to PINT0 interrupts, low-level or high-level detection can be selected individually for each pin by the PINT sense select bits (PINT7S to PINT0S) in interrupt control register 2 (ICR2). A single priority level in a range from 0 to 15 can be set for all PINT7 to PINT0 interrupts by bits 15 to 12 in interrupt priority register 05 (IPR05).

When using low-level sensing for the PINT7 to PINT0 interrupts, an interrupt request signal is sent to the INTC while the PINT7 to PINT0 pins are low. An interrupt request signal is stopped being sent to the INTC when the PINT7 to PINT0 pins are driven high. The status of the interrupt requests can be checked by reading the PINT interrupt request bits (PINT7R to PINT0R) in the PINT interrupt request register (PIRR). The above description also applies to when using highlevel sensing, except for the polarity being reversed. The PINT interrupt exception handling sets the I3 to I0 bits in SR to the priority level of the PINT interrupt.

When restoring from the service routine of PINT interrupt exception handling, execute the RTE instruction after an interrupt request has been cleared in the PINT interrupt request register (PIRR).

6.4.6 On-Chip Peripheral Module Interrupts

On-chip peripheral module interrupts are generated by the following on-chip peripheral modules:

- A/D converter (ADC)
- CD-ROM decoder (ROM-DEC)
- Multi-function timer pulse unit 2 (MTU2)
- Realtime clock (RTC)
- Watchdog timer (WDT)
- I²C bus interface 3 (IIC3)
- Direct memory access controller (DMAC)
- Serial communication interface with FIFO (SCIF)
- Controller area network (RCAN-ET)
- IEBusTM controller (IEB)
- Serial sound interface (SSI)
- 8-bit timer (TMR)

As every source is assigned a different interrupt vector, the source does not need to be identified in the exception service routine. A priority level in a range from 0 to 15 can be set for each module by interrupt priority registers 05 to 16 (IPR05 to IPR16). The on-chip peripheral module interrupt exception handling sets the I3 to I0 bits in SR to the priority level of the accepted on-chip peripheral module interrupt.

6.5 Interrupt Exception Handling Vector Table and Priority

Table 6.4 lists interrupt sources and their vector numbers, vector table address offsets, and interrupt priorities.

Each interrupt source is allocated a different vector number and vector table address offset. Vector table addresses are calculated from the vector numbers and vector table address offsets. In interrupt exception handling, the exception service routine start address is fetched from the vector table indicated by the vector table address. For details of calculation of the vector table address, see table 5.4, Calculating Exception Handling Vector Table Addresses, in section 5, Exception Handling.

The priorities of IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts can be set freely between 0 and 15 for each pin or module by setting interrupt priority registers 01, 02, and 05 to 16 (IPR01, IPR02, and IPR05 to IPR16). However, if two or more interrupts specified by the same IPR among IPR05 to IPR16 occur, the priorities are defined as shown in the IPR setting unit internal priority of table 6.4, and the priorities cannot be changed. A power-on reset assigns priority level 0 to IRQ interrupts, PINT interrupts, and on-chip peripheral module interrupts. If the same priority level is assigned to two or more interrupt sources and interrupts from those sources occur simultaneously, they are processed by the default priorities indicated in table 6.4.

Table 6.4 Interrupt Exception Handling Vectors and Priorities

6.6 Operation

6.6.1 Interrupt Operation Sequence

The sequence of interrupt operations is described below. Figure 6.2 shows the operation flow.

- 1. The interrupt request sources send interrupt request signals to the interrupt controller.
- 2. The interrupt controller selects the highest-priority interrupt from the interrupt requests sent, following the priority levels set in interrupt priority registers 01, 02, and 05 to 16 (IPR01, IPR02, and IPR05 to IPR16). Lower priority interrupts are ignored*. If two of these interrupts have the same priority level or if multiple interrupts occur within a single IPR, the interrupt with the highest priority is selected, according to the default priority and IPR setting unit internal priority shown in table 6.4.
- 3. The priority level of the interrupt selected by the interrupt controller is compared with the interrupt level mask bits (I3 to I0) in the status register (SR) of the CPU. If the interrupt request priority level is equal to or less than the level set in bits I3 to I0, the interrupt request is ignored. If the interrupt request priority level is higher than the level in bits I3 to I0, the interrupt controller accepts the interrupt and sends an interrupt request signal to the CPU.
- 4. The CPU detects the interrupt request sent from the interrupt controller when the CPU decodes the instruction to be executed. Instead of executing the decoded instruction, the CPU starts interrupt exception handling (figure 6.4).
- 5. The start address of the interrupt exception service routine is fetched from the exception handling vector table corresponding to the accepted interrupt.
- 6. The status register (SR) is saved onto the stack, and the priority level of the accepted interrupt is copied to bits I3 to I0 in SR.
- 7. The program counter (PC) is saved onto the stack.
- 8. The CPU jumps to the fetched start address of the interrupt exception service routine and starts executing the program. The jump that occurs is not a delayed branch.
- Notes: The interrupt source flag should be cleared in the interrupt handler. After clearing the interrupt source flag, "time from occurrence of interrupt request until interrupt controller identifies priority, compares it with mask bits in SR, and sends interrupt request signal to CPU" shown in table 6.5 is required before the interrupt source sent to the CPU is actually cancelled. To ensure that an interrupt request that should have been cleared is not inadvertently accepted again, read the interrupt source flag after it has been cleared, and then execute an RTE instruction.
	- Interrupt requests that are designated as edge-sensing are held pending until the interrupt requests are accepted. IRQ interrupts, however, can be cancelled by accessing the IRQ interrupt request register (IRQRR). For details, see section 6.4.4, IRQ Interrupts.

Interrupts held pending due to edge-sensing are cleared by a power-on reset or in deep standby mode.

Figure 6.2 Interrupt Operation Flow

6.6.2 Stack after Interrupt Exception Handling

Figure 6.3 Stack after Interrupt Exception Handling

6.7 Interrupt Response Time

Table 6.5 lists the interrupt response time, which is the time from the occurrence of an interrupt request until the interrupt exception handling starts and fetching of the first instruction in the interrupt exception service routine begins. The interrupt processing operations differ in the cases when banking is disabled, when banking is enabled without register bank overflow, and when banking is enabled with register bank overflow. Figures 6.4 and 6.5 show examples of pipeline operation when banking is disabled. Figures 6.6 and 6.7 show examples of pipeline operation when banking is enabled without register bank overflow. Figures 6.8 and 6.9 show examples of pipeline operation when banking is enabled with register bank overflow.

Table 6.5 Interrupt Response Time

Number of States

Notes: m1 to m4 are the number of states needed for the following memory accesses.

- m1: Vector address read (longword read)
- m2: SR save (longword write)
- m3: PC save (longword write)
- m4: Banked registers (R0 to R14, GBR, MACH, MACL, and PR) are restored from the stack.
- 1. In the case of $m1 = m2 = m3 = m4 = 1$ lcyc.
- 2. In the case of $I\phi$: $B\phi$: $P\phi = 120$:60:30 [MHz].

Figure 6.4 Example of Pipeline Operation when IRQ Interrupt is Accepted (No Register Banking)

Figure 6.6 Example of Pipeline Operation when IRQ Interrupt is Accepted (Register Banking without Register Bank Overflow)

Figure 6.7 Example of Pipeline Operation when Interrupt is Accepted during RESBANK Instruction Execution (Register Banking without Register Bank Overflow)

Figure 6.8 Example of Pipeline Operation when IRQ Interrupt is Accepted (Register Banking with Register Bank Overflow)

Figure 6.9 Example of Pipeline Operation when Interrupt is Accepted during RESBANK Instruction Execution (Register Banking with Register Bank Overflow)

6.8 Register Banks

This LSI has fifteen register banks used to perform register saving and restoration required in the interrupt processing at high speed. Figure 6.10 shows the register bank configuration.

Figure 6.10 Overview of Register Bank Configuration

6.8.1 Register Banks and Bank Control Registers

(1) Banked Register

The contents of the general registers (R0 to R14), global base register (GBR), multiply and accumulate registers (MACH and MACL), and procedure register (PR), and the vector table address offset (VTO) are banked.

(2) Input/Output of Banks

This LSI has fifteen register banks, bank 0 to bank 14. Register banks are stacked in first-in lastout (FILO) sequence. Saving takes place in order, beginning from bank 0, and restoration takes place in the reverse order, beginning from the last bank saved to.

6.8.2 Bank Save and Restore Operations

(1) Saving to Bank

Figure 6.11 shows register bank save operations. The following operations are performed when an interrupt for which usage of register banks is allowed is accepted by the CPU:

- (a) Assume that the bank number bit value in the bank number register (IBNR), BN, is i before the interrupt is generated.
- (b) The contents of registers R0 to R14, GBR, MACH, MACL, and PR, and the interrupt vector table address offset (VTO) of the accepted interrupt are saved in the bank indicated by BN, bank i.
- (c) The BN value is incremented by 1.

Figure 6.11 Bank Save Operations

Figure 6.12 shows the timing for saving to a register bank. Saving to a register bank takes place between the start of interrupt exception handling and the start of fetching the first instruction in the exception service routine.

Figure 6.12 Bank Save Timing

(2) Restoration from Bank

The RESBANK (restore from register bank) instruction is used to restore data saved in a register bank. After restoring data from the register banks with the RESBANK instruction at the end of the interrupt service routine, execute the RTE instruction to return from exception handling.

6.8.3 Save and Restore Operations after Saving to All Banks

If an interrupt occurs and usage of the register banks is enabled for the interrupt accepted by the CPU in a state where saving has been performed to all register banks, automatic saving to the stack is performed instead of register bank saving if the BOVE bit in the bank number register (IBNR) is cleared to 0. If the BOVE bit in IBNR is set to 1, register bank overflow exception occurs and data is not saved to the stack.

Save and restore operations when using the stack are as follows:

(1) Saving to Stack

- 1. The status register (SR) and program counter (PC) are saved to the stack during interrupt exception handling.
- 2. The contents of the banked registers (R0 to R14, GBR, MACH, MACL, and PR) are saved to the stack. The registers are saved to the stack in the order of MACL, MACH, GBR, PR, R14, R13, …, R1, and R0.
- 3. The register bank overflow bit (BO) in SR is set to 1.
- 4. The bank number bit (BN) value in the bank number register (IBNR) remains set to the maximum value of 15

(2) Restoration from Stack

When the RESBANK (restore from register bank) instruction is executed with the register bank overflow bit (BO) in SR set to 1, the CPU operates as follows:

- 1. The contents of the banked registers (R0 to R14, GBR, MACH, MACL, and PR) are restored from the stack. The registers are restored from the stack in the order of R0, R1, …, R13, R14, PR, GBR, MACH, and MACL.
- 2. The bank number bit (BN) value in the bank number register (IBNR) remains set to the maximum value of 15

6.8.4 Register Bank Exception

There are two register bank exceptions (register bank errors): register bank overflow and register bank underflow.

(1) Register Bank Overflow

This exception occurs if, after data has been saved to all of the register banks, an interrupt for which register bank use is allowed is accepted by the CPU, and the BOVE bit in the bank number register (IBNR) is set to 1. In this case, the bank number bit (BN) value in the bank number register (IBNR) remains set to the bank count of 15 and saving is not performed to the register bank.

(2) Register Bank Underflow

This exception occurs if the RESBANK (restore from register bank) instruction is executed when no data has been saved to the register banks. In this case, the values of R0 to R14, GBR, MACH, MACL, and PR do not change. In addition, the bank number bit (BN) value in the bank number register (IBNR) remains set to 0.

6.8.5 Register Bank Error Exception Handling

When a register bank error occurs, register bank error exception handling starts. When this happens, the CPU operates as follows:

- 1. The exception service routine start address which corresponds to the register bank error that occurred is fetched from the exception handling vector table.
- 2. The status register (SR) is saved to the stack.
- 3. The program counter (PC) is saved to the stack. The PC value saved is the start address of the instruction to be executed after the last executed instruction for a register bank overflow, and the start address of the executed RESBANK instruction for a register bank underflow. To prevent multiple interrupts from occurring at a register bank overflow, the interrupt priority level that caused the register bank overflow is written to the interrupt mask level bits (I3 to I0) of the status register (SR).
- 4. Program execution starts from the exception service routine start address.

6.9 Data Transfer with Interrupt Request Signals

Interrupt request signals can be used to activate the DMAC and transfer data.

Interrupt sources that are specified to activate the DMAC are masked by setting the DMA transfer enable bit in DREQER0 to DREQER3 to 1 without being input to the INTC.

6.9.1 Handling Interrupt Request Signals as Sources for CPU Interrupt but not DMAC Activation

- 1. Clear the corresponding DMAC transfer request enable bit in DREQER0 to DREQER3 to 0.
- 2. When an interrupt occurs, the interrupt request will be sent to the CPU.
- 3. The CPU clears the interrupt source and performs the necessary processing in the interrupt handling routine.

6.9.2 Handling Interrupt Request Signals as Sources for DMAC Activation but not CPU Interrupt

- 1. Select* the signals as DMAC activating sources by setting the corresponding DMAC transfer request enable bit in DREQER0 to DREQER3 to 1. This masks the CPU interrupt source regardless of the interrupt priority register settings.
- 2. When an interrupt occurs, the activation source will be sent to the DMAC.
- 3. The DMAC clears the activation source during the transfer.
- Note: * As for the method to select the DMAC request sources, see section 11, Direct Memory Access Controller (DMAC).

6.10 Usage Note

6.10.1 Timing to Clear an Interrupt Source

The interrupt source flags should be cleared in the interrupt handler. After clearing the interrupt source flag, "time from occurrence of interrupt request until interrupt controller identifies priority, compares it with mask bits in SR, and sends interrupt request signal to CPU" shown in table 6.5 is required before the interrupt source sent to the CPU is actually cancelled. To ensure that an interrupt request that should have been cleared is not inadvertently accepted again, read the interrupt source flag after it has been cleared, and then execute an RTE instruction.

Section 7 User Break Controller (UBC)

The user break controller (UBC) provides functions that simplify program debugging. These functions make it easy to design an effective self-monitoring debugger, enabling the chip to debug programs without using an in-circuit emulator. Instruction fetch or data read/write of CPU, data size, data contents, address value, and stop timing in the case of instruction fetch are break conditions that can be set in the UBC. Since this LSI uses a Harvard architecture, instruction fetch on the CPU bus (C bus) is performed by issuing bus cycles on the instruction fetch bus (F bus), and data access on the C bus is performed by issuing bus cycles on the memory access bus (M bus). The UBC monitors the C bus and internal bus (I bus).

7.1 Features

1. The following break comparison conditions can be set.

Number of break channels: two channels (channels 0 and 1)

User break can be requested as the independent condition on channels 0 and 1.

• Address

Comparison of the 32-bit address is maskable in 1-bit units.

One of the three address buses (F address bus (FAB), M address bus (MAB), and I address bus (IAB)) can be selected.

• Data

Comparison of the 32-bit data is maskable in 1-bit units.

One of the two data buses (M data bus (MDB) and I data bus (IDB)) can be selected.

• Bus cycle

Instruction fetch (only when C bus is selected) or data access

- Read/write
- Operand size
	- Byte, word, and longword
- 2. In an instruction fetch cycle, it can be selected whether the start of user break interrupt exception processing is set before or after an instruction is executed.
- 3. When a break condition is satisfied, a trigger signal is output from the UBCTRG pin.

Figure 7.1 shows a block diagram of the UBC.

7.2 Input/Output Pin

Table 7.1 shows the pin configuration of the UBC.

Table 7.1 Pin Configuration

7.3 Register Descriptions

The UBC has the following registers.

Table 7.2 Register Configuration

7.3.1 Break Address Register (BAR)

BAR is a 32-bit readable/writable register. BAR specifies the address used as a break condition in each channel. The control bits CD[1:0] in the break bus cycle register (BBR) select one of the three address buses for a break condition. BAR is initialized to H'00000000 by a power-on reset or in deep standby, but retains its previous value by a manual reset or in software standby mode or sleep mode.

Note: When setting the instruction fetch cycle as a break condition, clear the LSB in BAR to 0.

7.3.2 Break Address Mask Register (BAMR)

BAMR is a 32-bit readable/writable register. BAMR specifies bits masked in the break address bits specified by BAR. BAMR is initialized to H'00000000 by a power-on reset or in deep standby, but retains its previous value by a manual reset or in software standby mode or sleep mode.

7.3.3 Break Data Register (BDR)

BDR is a 32-bit readable/writable register. The control bits $CD[1:0]$ in the break bus cycle register (BBR) select one of the two data buses for a break condition. BDR is initialized to H'00000000 by a power-on reset or in deep standby, but retains its previous value by a manual reset or in software standby mode or sleep mode.

Notes: 1. Set the operand size when specifying a value on a data bus as the break condition.

 2. When the byte size is selected as a break condition, the same byte data must be set in bits 31 to 24, 23 to 16, 15 to 8, and 7 to 0 in BDR as the break data. Similarly, when the word size is selected, the same word data must be set in bits 31 to 16 and 15 to 0.

7.3.4 Break Data Mask Register (BDMR)

BDMR is a 32-bit readable/writable register. BDMR specifies bits masked in the break data bits specified by BDR. BDMR is initialized to H'00000000 by a power-on reset or in deep standby, but retains its previous value by a manual reset or in software standby mode or sleep mode.

Notes: 1. Set the operand size when specifying a value on a data bus as the break condition.

 2. When the byte size is selected as a break condition, the same byte data must be set in bits 31 to 24, 23 to 16, 15 to 8, and 7 to 0 in BDMR as the break mask data. Similarly, when the word size is selected, the same word data must be set in bits 31 to 16 and 15 to 0.

7.3.5 Break Bus Cycle Register (BBR)

BBR is a 16-bit readable/writable register, which specifies (1) disabling or enabling of user break interrupts, (2) including or excluding of the data bus value, (3) bus master of the I bus, (4) C bus cycle or I bus cycle, (5) instruction fetch or data access, (6) read or write, and (7) operand size as the break conditions. BBR is initialized to H'0000 by a power-on reset and in deep standby, but retains its previous value by a manual reset or in software standby mode or sleep mode.

x: Don't care

7.3.6 Break Control Register (BRCR)

BRCR sets the following conditions:

- 1. Specifies whether a start of user break interrupt exception processing by instruction fetch cycle is set before or after instruction execution.
- 2. Specifies the pulse width of the UBCTRG output when a break condition is satisfied.

BRCR is a 32-bit readable/writable register that has break condition match flags and bits for setting other break conditions. For the condition match flags of bits 15 to 12, writing 1 is invalid (previous values are retained) and writing 0 is only possible. To clear the flag, write 0 to the flag bit to be cleared and 1 to all other flag bits. BRCR is initialized to H'00000000 by a power-on reset and in deep standby, but retains its previous value by a manual reset or in software standby mode or sleep mode.

7.4 Operation

7.4.1 Flow of the User Break Operation

The flow from setting of break conditions to user break exception handling is described below:

- 1. The break address is set in the break address register (BAR). The masked address bits are set in the break address mask register (BAMR). The break data is set in the break data register (BDR). The masked data bits are set in the break data mask register (BDMR). The bus break conditions are set in the break bus cycle register (BBR). Three control bit groups of BBR (C bus cycle/I bus cycle select, instruction fetch/data access select, and read/write select) are each set. No user break will be generated if even one of these groups is set to 00. The relevant break control conditions are set in the bits of the break control register (BRCR). Make sure to set all registers related to breaks before setting BBR, and branch after reading from the last written register. The newly written register values become valid from the instruction at the branch destination.
- 2. In the case where the break conditions are satisfied and the user break interrupt request is enabled, the UBC sends a user break request to the INTC, sets the C bus condition match flag (SCMFC) or I bus condition match flag (SCMFD) for the appropriate channel, and outputs a pulse to the UBCTRG pin with the width set by the CKS[1:0] bits. Setting the UBID bit in BBR to 1 enables external monitoring of the trigger output without requesting user break interrupts.
- 3. On receiving a user break interrupt request signal, the INTC determines its priority. Since the user break interrupt has a priority level of 15, it is accepted when the priority level set in the interrupt mask level bits (I3 to I0) of the status register (SR) is 14 or lower. If the I3 to I0 bits are set to a priority level of 15, the user break interrupt is not accepted, but the conditions are checked, and condition match flags are set if the conditions match. For details on ascertaining the priority, see section 6, Interrupt Controller (INTC).
- 4. Condition match flags (SCMFC and SCMFD) can be used to check which condition has been satisfied. Clear the condition match flags during the user break interrupt exception processing routine. The interrupt occurs again if this operation is not performed.
- 5. There is a chance that the break set in channel 0 and the break set in channel 1 occur around the same time. In this case, there will be only one break request to the INTC, but these two break channel match flags may both be set.
- 6. When selecting the I bus as the break condition, note as follows:
	- Whether or not the access the CPU issued on the C bus is issued on the I bus depends on the setting of the cache. As regard to the I bus operation that depends on cache conditions, see table 8.8 in section 8, Cache.
- ⎯ When a break condition is specified for the I bus, only the data access cycle is monitored. The instruction fetch cycle (including cache update cycle) is not monitored.
- ⎯ If a break condition is specified for the I bus, even when the condition matches in an I bus cycle resulting from an instruction executed by the CPU, at which instruction the break is to be accepted cannot be clearly defined.

7.4.2 Break on Instruction Fetch Cycle

- 1. When C bus/instruction fetch/read/word or longword is set in the break bus cycle register (BBR), the break condition is the FAB bus instruction fetch cycle. Whether a start of user break interrupt exception processing is set before or after the execution of the instruction can be selected with the PCB0 or PCB1 bit in the break control register (BRCR) for the appropriate channel. If an instruction fetch cycle is set as a break condition, clear LSB in the break address register (BAR) to 0. A break cannot be generated as long as this bit is set to 1.
- 2. A break for instruction fetch which is set as a break before instruction execution occurs when it is confirmed that the instruction has been fetched and will be executed. This means a break does not occur for instructions fetched by overrun (instructions fetched at a branch or during an interrupt transition, but not to be executed). When this kind of break is set for the delay slot of a delayed branch instruction, the user break interrupt request is not received until the execution of the first instruction at the branch destination.
	- Note: If a branch does not occur at a delayed branch instruction, the subsequent instruction is not recognized as a delay slot.
- 3. When setting a break condition for break after instruction execution, the instruction set with the break condition is executed and then the break is generated prior to execution of the next instruction. As with pre-execution breaks, a break does not occur with overrun fetch instructions. When this kind of break is set for a delayed branch instruction and its delay slot, the user break interrupt request is not received until the first instruction at the branch destination.
- 4. When an instruction fetch cycle is set, the break data register (BDR) is ignored. Therefore, break data cannot be set for the break of the instruction fetch cycle.
- 5. If the I bus is set for a break of an instruction fetch cycle, the setting is invalidated.

7.4.3 Break on Data Access Cycle

- 1. If the C bus is specified as a break condition for data access break, condition comparison is performed for the logical addresses (and data) accessed by the executed instructions, and a break occurs if the condition is satisfied. If the I bus is specified as a break condition, condition comparison is performed for the addresses (and data) of the data access cycles on the bus specified by the I bus select bits, and a break occurs if the condition is satisfied. For details on the CPU bus cycles issued on the I bus, see 6 in section 7.4.1, Flow of the User Break Operation.
- 2. The relationship between the data access cycle address and the comparison condition for each operand size is listed in table 7.3.

This means that when address H'00001003 is set in the break address register (BAR), for example, the bus cycle in which the break condition is satisfied is as follows (where other conditions are met).

Longword access at H'00001000

Word access at H'00001002

Byte access at H'00001003

3. When the data value is included in the break conditions:

When the data value is included in the break conditions, either longword, word, or byte is specified as the operand size in the break bus cycle register (BBR). When data values are included in break conditions, a break is generated when the address conditions and data conditions both match. To specify byte data for this case, set the same data in the four bytes at bits 31 to 24, 23 to 16, 15 to 8, and 7 to 0 of the break data register (BDR) and break data mask register (BDMR). To specify word data for this case, set the same data in the two words at bits 31 to 16 and 15 to 0.

- 4. Access by a PREF instruction is handled as read access in longword units without access data. Therefore, if including the value of the data bus when a PREF instruction is specified as a break condition, a break will not occur.
- 5. If the data access cycle is selected, the instruction at which the break will occur cannot be determined.

7.4.4 Value of Saved Program Counter

When a user break interrupt request is received, the address of the instruction from where execution is to be resumed is saved to the stack, and the exception handling state is entered. If the C bus (FAB)/instruction fetch cycle is specified as a break condition, the instruction at which the break should occur can be uniquely determined. If the C bus/data access cycle or I bus/data access cycle is specified as a break condition, the instruction at which the break should occur cannot be uniquely determined.

1. When C bus (FAB)/instruction fetch (before instruction execution) is specified as a break condition:

The address of the instruction that matched the break condition is saved to the stack. The instruction that matched the condition is not executed, and the break occurs before it. However, when a delay slot instruction matches the condition, the instruction is executed, and the branch destination address is saved to the stack.

2. When C bus (FAB)/instruction fetch (after instruction execution) is specified as a break condition:

The address of the instruction following the instruction that matched the break condition is saved to the stack. The instruction that matches the condition is executed, and the break occurs before the next instruction is executed. However, when a delayed branch instruction or delay slot matches the condition, the instruction is executed, and the branch destination address is saved to the stack.

3. When C bus/data access cycle or I bus/data access cycle is specified as a break condition: The address after executing several instructions of the instruction that matched the break condition is saved to the stack.

7.4.5 Usage Examples

(1) Break Condition Specified for C Bus Instruction Fetch Cycle

(Example 1-1)

• Register specifications

```
BAR_0 = H'00000404, BAMR_0 = H'00000000, BBR_0 = H'0054, BAR_1 = H'00008010, 
BAMR_1 = H'00000006, BBR_1 = H'0054, BDR_1 = H'00000000, BDMR_1 = H'00000000, 
BRCR = H'00000020 
<Channel 0> 
Address: H'00000404, Address mask: H'00000000 
Bus cycle: C bus/instruction fetch (after instruction execution)/read (operand size is not 
           included in the condition) 
<Channel 1> 
Address: H'00008010, Address mask: H'00000006 
Data: H'00000000, Data mask: H'00000000 
Bus cycle: C bus/instruction fetch (before instruction execution)/read (operand size is not
```
included in the condition)

A user break occurs after an instruction of address H'00000404 is executed or before instructions of addresses H'00008010 to H'00008016 are executed.

(Example 1-2)

• Register specifications

BAR_0 = H'00027128, BAMR_0 = H'00000000, BBR_0 = H'005A, BAR_1= H'00031415, BAMR_1 = H'00000000, BBR_1 = H'0054, BDR_1 = H'00000000, BDMR_1 = H'00000000, BRCR = H'00000000

<Channel 0>

Address: H'00027128, Address mask: H'00000000

Bus cycle: C bus/instruction fetch (before instruction execution)/write/word

<Channel 1>

Address: H'00031415, Address mask: H'00000000

Data: H'00000000, Data mask: H'00000000

Bus cycle: C bus/instruction fetch (before instruction execution)/read (operand size is not included in the condition)

On channel 0, a user break does not occur since instruction fetch is not a write cycle. On channel 1, a user break does not occur since instruction fetch is performed for an even address. (Example 1-3)

Register specifications

BAR_0 = H'00008404, BAMR_0 = H'00000FFF, BBR_0 = H'0054, BAR_1= H'00008010, BAMR_1 = H'00000006, BBR_1 = H'0054, BDR_1 = H'00000000, BDMR_1 = H'00000000, $BRCR = H'00000020$

<Channel 0>

Address: H'00008404, Address mask: H'00000FFF

Bus cycle: C bus/instruction fetch (after instruction execution)/read (operand size is not included in the condition)

<Channel 1>

Address: H'00008010, Address mask: H'00000006

Data: H'00000000, Data mask: H'00000000

Bus cycle: C bus/instruction fetch (before instruction execution)/read (operand size is not included in the condition)

A user break occurs after an instruction with addresses H'00008000 to H'00008FFE is executed or before an instruction with addresses H'00008010 to H'00008016 are executed.

(2) Break Condition Specified for C Bus Data Access Cycle

(Example 2-1)

Register specifications

BAR_0 = H'00123456, BAMR_0 = H'00000000, BBR_0 = H'0064, BAR_1= H'000ABCDE, BAMR_1 = H'000000FF, BBR_1 = H'106A, BDR_1 = H'A512A512, BDMR_1 = H'00000000, BRCR = H'00000000 <Channel 0> Address: H'00123456, Address mask: H'00000000 Bus cycle: C bus/data access/read (operand size is not included in the condition) <Channel 1> Address: H'000ABCDE, Address mask: H'000000FF Data: H'0000A512, Data mask: H'00000000 Bus cycle: C bus/data access/write/word On channel 0, a user break occurs with longword read from address H'00123456, word read from address H'00123456, or byte read from address H'00123456. On channel 1, a user break occurs when word H'A512 is written in addresses H'000ABC00 to H'000ABCFE.

(3) Break Condition Specified for I Bus Data Access Cycle

(Example 3-1)

• Register specifications

BAR_0 = H'00314156, BAMR_0 = H'00000000, BBR_0 = H'0094, BAR_1= H'00055555, BAMR_1 = H'00000000, BBR_1 = H'11A9, BDR_1 = H'78787878, BDMR_1 = H'0F0F0F0F, BRCR = H'00000000

<Channel 0>

Address: H'00314156, Address mask: H'00000000

Bus cycle: I bus/instruction fetch/read (operand size is not included in the condition)

<Channel 1>

Address: H'00055555, Address mask: H'00000000

Data: H'00000078, Data mask: H'0000000F

Bus cycle: I bus/data access/write/byte

On channel 0, the setting of I bus/instruction fetch is ignored.

On channel 1, a user break occurs when the CPU writes byte data H'7x in address H'00055555 on the I bus.

7.5 Usage Notes

- 1. The CPU can read from or write to the UBC registers via the I bus. Accordingly, during the period from executing an instruction to rewrite the UBC register till the new value is actually rewritten, the desired break may not occur. In order to know the timing when the UBC register is changed, read from the last written register. Instructions after then are valid for the newly written register value.
- 2. The UBC cannot monitor access to the C bus and I bus cycles in the same channel.
- 3. When a user break interrupt request and another exception source occur at the same instruction, which has higher priority is determined according to the priority levels defined in table 5.1 in section 5, Exception Handling. If an exception source with higher priority occurs, the user break interrupt request is not received.
- 4. Note the following when a break occurs in a delay slot. If a pre-execution break is set at a delay slot instruction, the user break interrupt request is not received immediately before execution of the branch destination.
- 5. User breaks are disabled during UBC module standby mode. Do not read from or write to the UBC registers during UBC module standby mode; the values are not guaranteed.
- 6. Do not set an address within an interrupt exception handling routine whose interrupt priority level is at least 15 (including user break interrupts) as a break address.
- 7. Do not set break after instruction execution for the SLEEP instruction or for the delayed branch instruction where the SLEEP instruction is placed at its delay slot.
- 8. When setting a break for a 32-bit instruction, set the address where the upper 16 bits are placed. If the address of the lower 16 bits is set and a break before instruction execution is set as a break condition, the break is handled as a break after instruction execution.
- 9. Do not set a break after instruction execution for the DIVU or DIVS instruction. If a break after instruction execution is set for the DIVU or DIVS instruction and an exception or interrupt occurs during execution of the DIVU or DIVS instruction, a break after instruction execution occurs even though execution of the DIVU or DIVS instruction is halted.
- 10. Do not set a pre-execution break for the instruction that comes after the DIVU or DIVS instruction. If a pre-execution break is set for the instruction that comes after the DIVU or DIVS instruction and an exception or interrupt occurs during execution of the DIVU or DIVS instruction, a pre-execution break occurs even though execution of the DIVU or DIVS instruction is halted.
- 11. Do not set a pre-execution break and a break after instruction execution simultaneously in one address. For example, if a pre-execution break for channel 0 and a break after instruction execution for channel 1 are set simultaneously for one address, a break generated prior to instruction execution for channel 0 can set a condition-match flag after the instruction execution for channel 1.
Section 8 Cache

8.1 Features

- Capacity Instruction cache: 8 Kbytes Operand cache: 8 Kbytes
- Structure: Instructions/data separated, 4-way set associative
- Cache lock function (only for operand cache): Way 2 and way 3 are lockable
- Line size: 16 bytes
- Number of entries: 128 entries/way
- Write system: Write-back/write-through selectable
- Replacement method: Least-recently-used (LRU) algorithm

8.1.1 Cache Structure

The cache separates data and instructions and uses a 4-way set associative system. It is composed of four ways (banks), each of which is divided into an address section and a data section.

Each of the address and data sections is divided into 128 entries. The data section of the entry is called a line. Each line consists of 16 bytes (4 bytes \times 4). The data capacity per way is 2 Kbytes $(16 \text{ bytes} \times 128 \text{ entries})$, with a total of 8 Kbytes in the cache as a whole (4 ways) . Figure 8.1 shows the operand cache structure. The instruction cache structure is the same as the operand cache structure except for not having the U bit.

Figure 8.1 Operand Cache Structure

(1) Address Array

The V bit indicates whether the entry data is valid. When the V bit is 1, data is valid; when 0, data is not valid.

The U bit (only for operand cache) indicates whether the entry has been written to in write-back mode. When the U bit is 1, the entry has been written to; when 0, it has not.

The tag address holds the physical address used in the external memory access. It is composed of 21 bits (address bits 31 to 11) used for comparison during cache searches. In this LSI, as values of addresses in the cache valid space are from H'00000000 to H'1FFFFFFF (see section 9, Bus State Controller (BSC)), the upper three bits of the tag address are cleared to 0.

The V and U bits are initialized to 0 by a power-on reset and in deep standby mode but not initialized by a manual reset or in software standby mode.

The tag address is not initialized by a power-on reset or manual reset or in software standby mode. The tag address becomes undefined after deep standby.

(2) Data Array

Holds a 16-byte instruction or data. Entries are registered in the cache in line units (16 bytes).

The data array is not initialized by a power-on reset or manual reset or in software standby mode. The data array becomes undefined after deep standby.

(3) LRU

With the 4-way set associative system, up to four instructions or data with the same entry address can be registered in the cache. When an entry is registered, LRU shows which of the four ways it is recorded in. There are six LRU bits, controlled by hardware. A least-recently-used (LRU) algorithm is used to select the way that has been least recently accessed.

Six LRU bits indicate the way to be replaced in case of a cache miss. The relationship between LRU and way replacement is shown in table 8.1 when the cache lock function (only for operand cache) is not used (concerning the case where the cache lock function is used, see section 8.2.2, Cache Control Register 2 (CCR2)). If a bit pattern other than those listed in table 8.1 is set in the LRU bits by software, the cache will not function correctly. When modifying the LRU bits by software, set one of the patterns listed in table 8.1.

The LRU bits are initialized to B'000000 by a power-on reset and in deep standby but not initialized by a manual reset or in software standby mode.

LRU (Bits 5 to 0)	Way to be Replaced
000000, 000100, 010100, 100000, 110000, 110100	З
000001, 000011, 001011, 100001, 101001, 101011	
000110, 000111, 001111, 010110, 011110, 011111	
111000, 111001, 111011, 111100, 111110, 111111	O

Table 8.1 LRU and Way Replacement (Cache Lock Function Not Used)

8.2 Register Descriptions

The cache has the following registers.

Table 8.2 Register Configuration

8.2.1 Cache Control Register 1 (CCR1)

The instruction cache is enabled or disabled using the ICE bit. The ICF bit controls disabling of all instruction cache entries. The operand cache is enabled or disabled using the OCE bit. The OCF bit controls disabling of all operand cache entries. The WT bit selects either write-through mode or write-back mode for operand cache.

Programs that change the contents of CCR1 should be placed in an address space that is not cached, and an address space that is cached should be accessed after reading the contents of CCR₁

CCR1 is initialized to H'00000000 by a power-on reset and in deep standby but not initialized by a manual reset or in software standby mode.

SH7261 Group

Initial

Section 8 Cache

8.2.2 Cache Control Register 2 (CCR2)

CCR2 is used to enable or disable the cache locking function for operand cache and is valid in cache locking mode only. In cache locking mode, the lock enable bit (the LE bit) in CCR2 is set to 1. In non-cache-locking mode, the cache locking function is invalid.

When a cache miss occurs in cache locking mode by executing the prefetch instruction (PREF @Rn), the line of data pointed to by Rn is loaded into the cache according to bits 9 and 8 (the W3LOAD and W3LOCK bits) and bits 1 and 0 (the W2LOAD and W2LOCK bits) in CCR2. The relationship between the setting of each bit and a way, to be replaced when the prefetch instruction is executed, are listed in table 8.3. On the other hand, when the prefetch instruction is executed and a cache hit occurs, new data is not fetched and the entry which is already enabled is held. For example, when the prefetch instruction is executed with W3LOAD = 1 and W3LOCK = 1 specified in cache locking mode while one-line data already exists in way 0 which is specified by Rn, a cache hit occurs and data is not fetched to way 3.

In the cache access other than the prefetch instruction in cache locking mode, ways to be replaced by bits W3LOCK and W2LOCK are restricted. The relationship between the setting of each bit in CCR2 and ways to be replaced are listed in table 8.4.

Programs that change the contents of CCR2 should be placed in an address space that is not cached, and an address space that is cached should be accessed after reading the contents of CCR₂.

CCR2 is initialized to H'00000000 by a power-on reset and in deep standby but not initialized by a manual reset or in software standby mode.

Note: $*$ The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

SH7261 Group

Note: * The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

LE	W3LOAD*	W3LOCK	W2LOAD*	W2LOCK	Way to be Replaced
0	x	x	x	x	Decided by LRU (table 8.1)
	x		x		Decided by LRU (table 8.1)
	x	Ω			Decided by LRU (table 8.5)
			x		Decided by LRU (table 8.6)
					Decided by LRU (table 8.7)
		x			Way 2
				х	Way 3

Table 8.3 Way to be Replaced when a Cache Miss Occurs in PREF Instruction

[Legend]

x: Don't care

Note: * The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

[Legend]

x: Don't care

Note: * The W3LOAD and W2LOAD bits should not be set to 1 at the same time.

Table 8.5 LRU and Way Replacement (when W2LOCK=1 and W3LOCK=0)

Table 8.6 LRU and Way Replacement (when W2LOCK=0 and W3LOCK=1)

Table 8.7 LRU and Way Replacement (when W2LOCK=1 and W3LOCK=1)

8.3 Operation

Operations for the operand cache are described here. Operations for the instruction cache are similar to those for the operand cache except for the address array not having the U bit, and there being no prefetch operation or write operation, or a write-back buffer.

8.3.1 Searching Cache

If the operand cache is enabled (OCE bit in CCR1 is 1), whenever data in a cache-enabled area is accessed, the cache will be searched to see if the desired data is in the cache. Figure 8.2 illustrates the method by which the cache is searched.

Entries are selected using bits 10 to 4 of the address used to access memory and the tag address of that entry is read. At this time, the upper three bits of the tag address are always cleared to 0. Bits 31 to 11 of the address used to access memory are compared with the read tag address. The address comparison uses all four ways. When the comparison shows a match and the selected entry is valid $(V = 1)$, a cache hit occurs. When the comparison does not show a match or the selected entry is not valid $(V = 0)$, a cache miss occurs. Figure 8.2 shows a hit on way 1.

Figure 8.2 Cache Search Scheme

8.3.2 Read Access

(1) Read Hit

In a read access, data is transferred from the cache to the CPU. LRU is updated so that the hit way is the latest.

(2) Read Miss

An external bus cycle starts and the entry is updated. The way replaced follows table 8.4. Entries are updated in 16-byte units. When the desired data that caused the miss is loaded from external memory to the cache, the data is transferred to the CPU in parallel with being loaded to the cache. When it is loaded in the cache, the V bit is set to 1, and LRU is updated so that the replaced way becomes the latest. In operand cache, the U bit is additionally cleared to 0. When the U bit of the entry to be replaced by updating the entry in write-back mode is 1, the cache update cycle starts after the entry is transferred to the write-back buffer. After the cache completes its update cycle, the write-back buffer writes the entry back to the memory. The write-back unit is 16 bytes.

8.3.3 Prefetch Operation (Only for Operand Cache)

(1) Prefetch Hit

LRU is updated so that the hit way becomes the latest. The contents in other caches are not modified. No data is transferred to the CPU.

(2) Prefetch Miss

No data is transferred to the CPU. The way to be replaced follows table 8.3. Other operations are the same in case of read miss.

8.3.4 Write Operation (Only for Operand Cache)

(1) Write Hit

In a write access in write-back mode, the data is written to the cache and no external memory write cycle is issued. The U bit of the entry written is set to 1 and LRU is updated so that the hit way becomes the latest.

In write-through mode, the data is written to the cache and an external memory write cycle is issued. The U bit of the written entry is not updated and LRU is updated so that the replaced way becomes the latest.

(2) Write Miss

In write-back mode, an external bus cycle starts when a write miss occurs, and the entry is updated. The way to be replaced follows table 8.4. When the U bit of the entry to be replaced is 1, the cache update cycle starts after the entry is transferred to the write-back buffer. Data is written to the cache, the U bit is set to 1, and the V bit is set to 1. LRU is updated so that the replaced way becomes the latest. After the cache completes its update cycle, the write-back buffer writes the entry back to the memory. The write-back unit is 16 bytes.

In write-through mode, no write to cache occurs in a write miss; the write is only to the external memory.

8.3.5 Write-Back Buffer (Only for Operand Cache)

When the U bit of the entry to be replaced in the write-back mode is 1, it must be written back to the external memory. To increase performance, the entry to be replaced is first transferred to the write-back buffer and fetching of new entries to the cache takes priority over writing back to the external memory. After the cache completes to fetch the new entry, the write-back buffer writes the entry back to external memory. During the write-back cycles, the cache can be accessed. The write-back buffer can hold one line of cache data (16 bytes) and its physical address. Figure 8.3 shows the configuration of the write-back buffer.

Figure 8.3 Write-Back Buffer Configuration

Operations in sections 8.3.2 to 8.3.5 are compiled in table 8.8

Table 8.8 Cache Operations

[Legend]

x: Don't care

- Note: Cache renewal cycle: 16-byte read access, write-back cycle in write-back buffer: 16-byte write access
	- * Neither LRU renewed. LRU is renewed in all other cases.

8.3.6 Coherency of Cache and External Memory

Use software to ensure coherency between the cache and the external memory. When memory shared by this LSI and another device is mapped in the address space to be cached, operate the memory-mapped cache to invalidate and write back as required.

8.4 Memory-Mapped Cache

To allow software management of the cache, cache contents can be read and written by means of MOV instructions. The instruction cache address array is mapped onto addresses H'F0000000 to H'F07FFFFF, and the data array onto addresses H'F1000000 to H'F17FFFFF. The operand cache address array is mapped onto addresses H'F0800000 to H'F0FFFFFF, and the data array onto addresses H'F1800000 to H'F1FFFFFF. Only longword can be used as the access size for the address array and data array, and instruction fetches cannot be performed.

8.4.1 Address Array

To access an address array, the 32-bit address field (for read/write accesses) and 32-bit data field (for write accesses) must be specified.

In the address field, specify the entry address selecting the entry, the W bit for selecting the way, and the A bit for specifying the existence of associative operation. In the W bit, B'00 is way 0, B'01 is way 1, B'10 is way 2, and B'11 is way 3. Since the access size of the address array is fixed at longword, specify B'00 for bits 1 and 0 of the address.

The tag address, LRU bits, U bit (only for operand cache), and V bit are specified as data. Always specify 0 for the upper three bits (bits 31 to 29) of the tag address.

For the address and data formats, see figure 8.4.

The following three operations are possible for the address array.

(1) Address Array Read

The tag address, LRU bits, U bit (only for operand cache), and V bit are read from the entry address specified by the address and the entry corresponding to the way. For the read operation, associative operation is not performed regardless of whether the associative bit (A bit) specified by the address is 1 or 0.

(2) Address-Array Write (Non-Associative Operation)

When the associative bit (A bit) in the address field is cleared to 0, write the tag address, LRU bits, U bit (only for operand cache), and V bit, specified by the data field, to the entry address specified by the address and the entry corresponding to the way. When writing to a cache line for which the U bit = 1 and the V bit = 1 in the operand cache address array, write the contents of the cache line back to memory, then write the tag address, LRU bits, U bit, and V bit specified by the data field. When 0 is written to the V bit, 0 must also be written to the U bit of that entry.

(3) Address-Array Write (Associative Operation)

When writing with the associative bit (A bit) of the address field set to 1, the addresses in the four ways for the entry specified by the address field are compared with the tag address that is specified by the data field. Write the U bit (only for operand cache) and the V bit specified by the data field to the entry of the way that has a hit. However, the tag address and LRU bits remain unchanged. When there is no way that has a hit, nothing is written and there is no operation.

This function is used to invalidate a specific entry in the cache. When the U bit of the entry that has had a hit is 1 in the operand cache, writing back should be performed. However, when 0 is written to the V bit, 0 must also be written to the U bit of that entry.

8.4.2 Data Array

To access a data array, the 32-bit address field (for read/write accesses) and 32-bit data field (for write accesses) must be specified. The address field specifies information for selecting the entry to be accessed; the data field specifies the longword data to be written to the data array.

Specify the entry address for selecting the entry, the L bit indicating the longword position within the (16-byte) line, and the W bit for selecting the way. In the L bit, B'00 is longword 0, B'01 is longword 1, B'10 is longword 2, and B'11 is longword 3. In the W bit, B'00 is way 0, B'01 is way 1, B'10 is way 2, and B'11 is way 3. Since the access size of the data array is fixed at longword, specify B'00 for bits 1 and 0 of the address.

For the address and data formats, see figure 8.4.

The following two operations are possible for the data array. Information in the address array is not modified by this operation.

(1) Data Array Read

The data specified by the L bit in the address is read from the entry address specified by the address and the entry corresponding to the way.

(2) Data Array Write

The longword data specified by the data is written to the position specified by the L bit in the address from the entry address specified by the address and the entry corresponding to the way.

Figure 8.4 Specifying Address and Data for Memory-Mapped Cache Access

8.4.3 Usage Examples

(1) Invalidating Specific Entries

Specific cache entries can be invalidated by writing 0 to the entry's V bit in the memory mapping cache access. When the A bit is 1, the tag address specified by the write data is compared to the tag address within the cache selected by the entry address, and data is written to the bits V and U specified by the write data when a match is found. If no match is found, there is no operation. When the V bit of an entry in the address array is set to 0, the entry is written back if the entry's U bit is 1.

An example when a write data is specified in R0 and an address is specified in R1 is shown below.

```
; R0=H'0110 0010; tag address(28-11)=B'0 0001 0001 0000 0000 0, U=0, V=0 
; R1=H'F080 0088; operand cache address array access, entry=B'000 1000, A=1 
; 
 MOV.L R0,@R1
```
(2) Reading the Data of a Specific Entry

The data section of a specific cache entry can be read by the memory mapping cache access. The longword indicated in the data field of the data array in figure 8.4 is read into the register.

An example when an address is specified in R0 and data is read in R1 is shown below.

```
; R0=H'F100 004C; instruction cache data array access, entry=B'000 0100, 
; Way=0, longword address=3 
; 
 MOV.L @R0,R1
```
8.4.4 Notes

- 1. Programs that access memory-mapped cache of the operand cache should be placed in a cachedisabled space. Programs that access memory-mapped cache of the instruction cache should be placed in a cache-disabled space, and in each of the beginning and the end of that, two or more read accesses to on-chip peripheral modules or external address space (cache-disabled address) should be executed.
- 2. Rewriting the address array contents so that two or more ways are hit simultaneously is prohibited. Operation is not guaranteed if the address array contents are changed so that two or more ways are hit simultaneously.
- 3. Memory-mapped cache can be accessed only by the CPU and not by the DMAC. Registers can be accessed by the CPU and the DMAC.

Section 9 Bus State Controller (BSC)

The bus state controller (BSC) outputs control signals for various types of memory that is connected to the external address space and external devices. This enables the LSI to connect directly with SRAM, SDRAM, and other memory storage devices, and external devices.

9.1 Features

- 1. External address space
- A maximum of 64 Mbytes for the SDRAM and each for areas CS0 to CS6 (256 Mbytes for $CS₆$
- Ability to select the data bus width (8, 16, or 32 bits) independently for each address space
- 2. Normal space interface
- Supports an interface for direct connection to SRAM
- Cycle wait function: Maximum of 31 wait states (maximum of seven wait states for page access cycles)
- Wait control
	- Ability to select the assert/negate timing for chip select signals
	- Ability to select the assert/negate timing for the read strobe and write strobe signals
	- ⎯ Ability to select the data output start/end timing
	- ⎯ Ability to select the delay for chip select signals
- Write access modes: One-write strobe and byte-write strobe modes
- Page access mode: Support for page read and page write (64-bit, 128-bit, and 256-bit page units)
- 3. SDRAM interface
- Ability to set SDRAM in up to two areas
- Refresh functions
	- Auto-refresh (on-chip programmable refresh counter)
	- Self-refresh
- Ability to select the access timing (support for low column latency, column latency, and low active interval settings)
- Initialization sequencer function, power-down function, deep-power-down function, and mode register setting function implemented on-chip

Figure 9.1 shows a block diagram of the BSC.

Figure 9.1 Block Diagram of BSC

9.2 Input/Output Pins

Table 9.1 shows the pin configuration of the BSC.

Table 9.1 Pin Configuration

9.3 Area Overview

9.3.1 Address Map

In the architecture, this LSI has a 32-bit address space, which is divided into cache-enabled, cache-disabled, and on-chip spaces (on-chip RAM, on-chip peripheral modules, and reserved areas) according to the upper bits of the address.

External address spaces CS5 to CS0 are cache-enabled when internal address $A29 = 0$ and cachedisabled when A29 = 1. The CS6 space is always cache-disabled.

The kind of memory to be connected and the data bus width are specified independently for each partial space. The address map for the external address space is listed below.

Note: * For the on-chip RAM space, access the addresses shown in section 26, On-Chip RAM. For the on-chip peripheral module space, access the addresses shown in section 30, List of Registers. Do not access addresses which are not described in these sections. Otherwise, correct operation cannot be guaranteed.

9.3.2 Data Bus Width and Pin Function Setting for Individual Areas

In this LSI the data bus width of area 0 can be set to 8, 16, or 32 bits through external pins during a power-on reset. The data bus widths of areas 1 to 6 can be modified through register settings during program execution. Note that the selectable data bus widths may be limited depending on the connected memory type.

After a power-on reset, the LSI starts execution of the program stored in the external memory allocated in area 0.

For details on pin function settings, see section 25, Pin Function Controller (PFC).

9.4 Register Descriptions

The BSC has the following registers.

All registers are initialized by a power-on reset or in deep standby mode.

Do not access spaces other than area 0 until settings are completed for the connected memory interface.

Table 9.4 Register Configuration

Note: * Depends on the setting of the MD pin.

9.4.1 CSn Control Register (CSnCNT) (n = 0 to 6)

CSnCNT selects the width of the external bus and controls the operation of the CSC interface.

Notes: 1. The initial value of the BSIZE bits in CS0 differs depending on the settings of pins MD1 and MD0.

2. The initial value of the EXENB bit in CS0 is 1.

To disable the operation for each channel, forcibly write out data tentatively stored in internal write buffer. The procedure is as follows:

- 1. Execute read access to the channel whose operation is to be disabled.
- 2. Then, write 0 to the EXENB bit (operation disabled).

9.4.2 CSn Recovery Cycle Setting Register (CSnREC) (n = 0 to 6)

CSnREC specifies the number of data recovery cycles to be inserted after read or write accesses.

- Notes: 1. When accessing SDRAM, there is no danger of data collision on the bus due to timing. Consequently, there is no data recovery cycle setting for SDRAM. (The value is fixed at 0 cycles.)
	- 2. Writing to the CSn recovery cycle setting register (CSnREC) must be done while CSC for the corresponding channel is disabled ($EXENB = 0$). Only channel 0 (CS0) can be enabled by setting EXENB = 1. To enable channel 0, stop the DMAC and set EXENB to 1 between the reset release and data write access to CS0.

9.4.3 SDRAMCm Control Register (SDCmCNT) (m = 0, 1)

To disable the operation for each channel, forcibly write out data tentatively stored in internal write buffer. The procedure is as follows:

- 1. Execute read access to the channel whose operation is to be disabled.
- 2. Then, write 0 to the EXENB bit (operation disabled).

9.4.4 CSn Mode Register (CSMODn) (n = 0 to 6)

CSMODn selects the mode for page read access and the bit boundary for page access, enables page read/write access and external wait, and selects the mode for write access.

Writing to the CSn mode register (CSMODn) must be done while CSC for the corresponding channel is disabled (EXENB = 0). Only channel 0 (CS0) can be enabled by setting EXENB = 1. To enable channel 0, stop the DMAC and set EXENB to 1 between the reset release and data write access to CS0.

9.4.5 CSn Wait Control Register 1 (CS1WCNTn) (n = 0 to 6)

CS1WCNTn specifies the number of wait states inserted into the read/write cycle or page read/page write cycle.

- Notes: 1. Make sure the page read and page write cycle wait select (CSPRWAIT and CSPWWAIT) settings are within the range defined by the read and write cycle wait select (CSRWAIT and CSWWAIT) settings. Select each wait cycle number according the system configuration incorporated.
	- 2. Writing to the CSn wait control register 1 (CS1WCNTn) must be done while CSC for the corresponding channel is disabled ($EXENB = 0$). Only channel 0 (CS0) can be enabled by setting EXENB = 1. To enable channel 0, stop the DMAC and set EXENB to 1 between the reset release and data write access to CS0.
9.4.6 CSn Wait Control Register 2 (CS2WCNTn) (n = 0 to 6)

CS2WCNTn specifies the number of wait states and the number of delay cycles.

9.4.7 SDRAM Refresh Control Register 0 (SDRFCNT0)

Bit Bit Name Initial Value R/W Description $31 \text{ to } 1$ $\qquad \qquad$ All 0 R Reserved These bits are always read as 0. The write value should always be 0. 0 DSFEN 0 R/W SDRAM Common Self-Refresh Operation Enable This bit controls self-refresh operation for all channels simultaneously. Setting DSFEN to 1 performs autorefresh cycle operation, immediately after which selfrefresh operation begins. Clearing DSFEN to 0 ends self-refresh operation, and auto-refresh operation resumes immediately afterward. The value written to this bit is reflected when self-refresh operation starts, if DSFEN was set to 1, or when auto-refresh operation starts following the end of self-refresh operation, if DSFEN was cleared to 0. 0: Self-refresh disabled 1: Self-refresh enabled

SDRFCNT0 controls self-refresh operation.

9.4.8 SDRAM Refresh Control Register 1 (SDRFCNT1)

SDRFCNT1 controls auto-refresh operation.

Note: Auto-refresh requests are not accepted while multiple read or write accesses are in progress, or during a transfer using DMAC, so the auto-refresh interval may become enlarged in some cases. Set the DRFC bits to an auto-refresh request interval value that satisfies the auto-refresh interval specification of the SDRAM being used. Furthermore, make sure to set the auto-refresh request interval to a duration longer than the auto-refresh cycle.

Auto-Refresh Request Interval and DRFC Set Value:

SDRAMC includes a 12-bit refresh counter that generates auto-refresh requests at fixed intervals. The following equation is used to calculate the set value for the DRFC bits from the auto-refresh request interval.

 $DRFC = (Auto-refresh request interval / Bus clock cycle) - 1$

Auto-refresh requests are not accepted while SDRAM is being accessed; they must wait until the access completes. However, the counter value is updated regardless or whether or not the request was accepted. Note that if two or more auto-refresh requests are generated while SDRAM is being accessed, the second and subsequent requests are ignored.

9.4.9 SDRAM Initialization Register 0 (SDIR0)

SDIR0 specifies the SDRAM initialization sequence timing.

Note: Make settings that satisfy the specifications of the connected SDRAM before starting the initialization sequence.

9.4.10 SDRAM Initialization Register 1 (SDIR1)

SDIR1 controls activation of the SDRAM initialization sequence.

9.4.11 SDRAM Power-Down Control Register (SDPWDCNT)

SDPWDCNT controls transition to and recovery from power-down mode.

9.4.12 SDRAM Deep-Power-Down Control Register (SDDPWDCNT)

Bit: 31 30 29 28 27 26 25 24 23 22 21 20 19 18 17 16 Initial value: R/W: Bit: 15 14 13 12 11 10 9 8 7 6 5 4 3 2 1 0 Initial value: R/W: 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 R R R R R R R R R R R R R R R R 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 R R R R R R R R R R R R R R R R/W — — — — — — — — — \equiv — — — — \overline{a} — — \overline{a} — — — — — — — — \overline{a} — — — $-$ DDPD

SDDPWDCNT controls transition to and recovery from deep-power-down mode.

9.4.13 SDRAMm Address Register (SDmADR) (m = 0, 1)

SDmADR specifies the data bus width and the channel size of SDRAM.

9.4.14 SDRAMm Timing Register (SDmTR) (m = 0, 1)

SDmTR specifies the timing for read and write accesses to SDRAM.

[Legend] x: Don't care

9.4.15 SDRAMm Mode Register (SDmMOD) (m = 0, 1)

SDmMOD specifies the values to be written to the SDRAM mode register or extended mode register. Writing to this register causes a mode register set command or extended mode register set command to be issued automatically to SDRAM.

Notes: The following points should be kept in mind regarding SDRAMm mode register settings.

- 1. Make sure to set a burst length of 1 for SDRAM. Operation cannot be guaranteed with settings other than burst length 1.
- 2. The SDRAM column latency must match the setting of the SDRAM controller column latency setting bits (DCL) in SDRAMC. Operation cannot be guaranteed if the latency settings do not agree.
- 3. Check to make sure the status bits (DSRFST, DPWDST, DDPDST, and DMRSST) in the SDRAM status register (SDSTR) are all cleared to 0.

9.4.16 SDRAM Status Register (SDSTR)

SDSTR consists of the status flags that indicate the status of operation during self-refresh, initialization sequences, power-down mode, deep-power-down mode, and mode register setting.

"Transition to or recovery from in progress" refers to the interval from the point at which the bits listed in table 9.5 are written to until the corresponding commands are issued.

Note: Execution of a self-refresh, a transition to or recovery from power-down or deep-powerdown mode, an initialization sequence, or mode register setting may only be performed when all status bits are cleared to 0. Do not rewrite the registers (bits) listed below when any of the status bits (DSRFST, DINIST, DPWDST, DDPDST, DMRSST) is set to 1.

9.4.17 SDRAM Clock Stop Control Signal Setting Register (SDCKSCNT)

SDCKSCNT enables or disables the clock stop control signal (internal signal in the chip) and specifies the number of assert cycles.

9.4.18 AC Characteristics Switching Register (ACSWR)

When writing to the external address space or making SDRAM settings in power-on reset exception handling or cancellation of deep standby mode, be sure to set bits ACOSW[3:0] in ACSWR to B'0011 beforehand.

ACSWR is initialized to H'00000000 by a power-on reset and entry to deep standby mode, but is not initialized by a manual reset, entry to sleep mode, or entry to software standby mode.

9.5 Operation

9.5.1 CSC Interface

(1) Normal Access

Normal read/write operation is used for all bus access when page read/write access is disabled $(PRENB = 0, PWENB = 0)$. Even when page read/write access is enabled $(PRENB = 1, PWENB)$ = 1), normal read/write operation is employed in cases where page access cannot be used. Figure 9.2 shows the basic operation of the external bus control signals in read operation, and figure 9.3 shows the basic operation of these signals in write operation.

Figure 9.2 Basic Bus Timing (Read Operation)

Figure 9.3 Basic Bus Timing (Write Operation)

1. Ts (Internal Bus Access Start)

This is a bus access request cycle initiated by the internal bus master and with the external bus as the target. CSn is always high during this cycle. In the next cycle A27 to A0 and the write data change.

2. Tw1 to Twn (Read Cycle Wait, Write Cycle Wait)

These are the cycles between internal bus access start and the wait end cycle. A duration of from 0 to 31 clocks may be selected. During this interval the CSn, RD, and WR control signals are asserted (low level) in accordance with the wait settings. The assert timing can be controlled using the CS assert wait, RD assert wait, WR assert wait, and write data output wait bits in CSn control registers 1 and 2. The number of wait cycles can be set to from 0 to 7 clocks, with the count starting from the cycle following internal bus access start (Ts). The number of clocks selected must be no greater than the number of read/write cycle wait cycles.

3. Tend (Wait End Cycle)

This is the final cycle in a series of read cycle wait or write cycle wait cycles. The RD or WR signal is negated (high level) in the next cycle.

4. Tn1 to Tnm (CS Delay Cycle)

These are the cycles between the wait end cycle and when CSn is negated (high level). The negation timing can be controlled using write data output delay cycles. The number of cycles is counted beginning from the wait end cycle. In write access or if the number of CS delay cycles during a read is other than 0 or 1, the succeeding bus access can start from the cycle following the CS delay cycle end. If the number of CS delay cycles is 0 or 1 in read access, the succeeding bus access can start after the end of the read data sample cycle (see below).

5. Trd (Read Data Sample Cycle)

This is the sample cycle for read data.

(2) Page Access

Page read and write operation is employed for bus accesses for which page access can be used if page write access enable ($PWENB = 1$) and page read access enable ($PRENB = 1$) have been selected. Page access is used in the following cases.

- 1. CPU burst access (cache replacement)
- 2. When longword (32-bit) access to an 8-bit or 16-bit external data bus has been performed
- 3. When word (16-bit) access to an 8-bit external data bus has been performed

Table 9.6 shows the way addresses are modified in cases 1 above.

Figure 9.4 shows the basic operation of the external bus control signals in page read operation, and figure 9.5 shows the basic operation of these signals in write operation. Note that if the number of data bits accessed in a single burst is greater than the single page access bit boundary setting of the PBCNT bits in the mode register, a single burst access will trigger multiple page accesses. Regardless of whether the bust mode is increment or wraparound, page access stops temporarily (the CSn signal is negated) at the point when the address exceeds the page boundary, and page access operation starts again. If the number of data bits accessed in a single burst is smaller than the page boundary bit count, a single page access is sufficient to complete the burst transfer.

Figure 9.4 Basic Bus Timing (Page Read Operation)

Figure 9.5 Basic Bus Timing (Page Write Operation)

1. Ts (Internal Bus Access Start)

This is a bus access request cycle initiated by the internal bus master and with the external bus as the target. CSn is always high during this cycle. In the next cycle A27 to A0 and the write data change.

2. Tw1 to Twn (Read Cycle Wait, Write Cycle Wait)

For the first page access, the wait operation from internal bus access start to the wait end cycle is the same as in normal access.

3. Tend (First Wait End Cycle)

This is the final cycle in the first series of read cycle wait or write cycle wait cycles. In write access, the second and subsequent page accesses start from the next cycle, unless a write data output delay cycle has been specified (with a value other than 0). The RD or WR signal is negated (high level) in the next cycle if the RD assert wait or WD assert wait setting is other than 0. If the RD assert wait or WD assert wait setting is 0, the RD or WR signal continues to be asserted (low level). The CSn signal is not negated and continues to be asserted (low level). In page read access, the succeeding bus access starts without waiting for the read data sample cycle (Trd).

4. Tdw1 to Tdwn (Write Data Output Delay Cycle)

In write access write data output delay cycles are inserted between the wait end cycle and the following page access if the write data output delay wait setting is other than 0. Assertion of the address and output data is extended for the duration of this interval. Also, the WR signal is negated (high level).

5. Tpw1 to Tpwn (Page Read Cycle Wait, Page Write Cycle Wait)

In page access the page read cycle wait and page write cycle wait settings are used in place of the read cycle wait and write cycle wait settings for the second and subsequent bus cycles. The WR assert wait setting works the same as during the first bus cycle. The RD assert wait setting operates differently depending on the page read access mode (PRMOD) setting value.

 $PRMOD = 0$: RD assert wait setting operates identically to first bus cycle.

 $PRMOD = 1$: RD assert wait setting is invalid. Operation is the same as an RD assert wait setting of 0.

- 6. Tend/Tdw1 to Tdwn (Wait End Cycle/Write Data Output Delay Cycle) These operate the same as during the first access (3 and 4 above).
- 7. Tn1 to Tnm (CS Delay Cycle)

These are the cycles between the final wait end cycle and when CSn is negated (high level). The number of CS delay cycles is counted beginning from the wait end cycle.

8. Trd (Final Read Data Sample Cycle)

This is the final sample cycle for read data.

(3) External Wait Function

The external wait signal (\overline{WATT}) can be used to extend the wait cycle duration beyond the value specified by the cycle wait (CSRWAIT, CSWWAIT) or page access cycle wait (CSPRWAIT, CSPWWAIT) settings in the CSn wait control register (CSWCNTn). If external wait enable $(EWENB = 1)$ has been selected, wait cycles are inserted for as long as the WAIT signal remains low level. The WAIT signal is disabled if external wait disable (EWENB = 0) has been selected.

Note that the wait cycles specified by the settings of the CSn wait control register (CSWCNTn) are inserted regardless of the state of the WAIT signal.

(a) Normal Read/Write Operation

The \overline{WAIT} signal is sampled all the time and its result is reflected two cycles later. Thus, when the WAIT signal is low two cycles before the end of the wait cycles, external cycles are inserted. After the WAIT signal has gone high, the wait cycles end two cycles later.

(b) Page Access Operation

The initial data read/write operation is the same as a normal read/write operation. That is, when the WAIT signal is low two cycles before the end of the wait cycles (Tend), external wait cycles are inserted. After the WAIT signal has gone high, the wait cycles end (Tend) two cycles later.

In the second and subsequent read accesses, the page wait cycle is extended if the WAIT signal is low two cycles before the end of the page access wait cycle (Tend), and the page wait cycles end two cycles after the \overline{WAIT} signal has gone high.

Figure 9.6 shows an example of external wait timing for page read access using longword (32-bit) access to a 16-bit channel.

Figure 9.6 External Wait Timing Example (Page Read Access to 16-Bit Channel)

(4) Access Type and Data Alignment

(a) 32-Bit Bus Channel

If a 32-bit bus is selected by the external bus width select bits in the CSn control register, A27 to A2 are enabled as address signals for longword units and A1 and A0 are disabled (fixed low level). Table 9.7 shows the data alignment corresponding to byte addresses for different data sizes.

Pins $\overline{WR3}$ to $\overline{WR0}$ are enabled when byte strobe mode (WRMOD = 0) is selected. Pins $\overline{BC3}$ to BC0 are not used.

Only the $\overline{WR3}$ pin is enabled when one-write strobe mode (WRMOD = 1) is selected. A low-level signal is output from the $\overline{WR3}$ pin during write access, regardless of the data size. At this time pins $\overline{WR2}$ to $\overline{WR0}$ are disabled (fixed high level). The valid byte positions are indicated by pins $\overline{BC3}$ to $\overline{BC0}$.

Table 9.7 Data Alignment (32-Bit Bus Channel)

Note: The valid bits in the data bus for each data size are indicated by circles (O). Crosses (x) indicate bus data bits that are undefined.

(b) 16-Bit Bus Channel

If a 16-bit bus is selected by the external bus width select bits in the CSn control register, A27 to A1 are enabled as address signals for word units and A0 is disabled (fixed low level). Table 9.8 shows the data alignment corresponding to byte addresses for different data sizes.

Pins $\overline{WR1}$ and $\overline{WR0}$ are enabled when byte strobe mode (WRMOD = 0) is selected. Pins $\overline{WR3}$ and $\overline{WR2}$ are disabled. Pins $\overline{BC3}$ to $\overline{BC0}$ are not used.

Only the $\overline{WR1}$ pin is enabled when one-write strobe mode (WRMOD = 1) is selected. A low-level signal is output from the WR1 pin during write access, regardless of the data size. At this time the $\overline{WR0}$ pin is disabled (fixed high level). The valid byte positions are indicated by pins $\overline{BC1}$ and $\overline{BC0}$

Table 9.8 Data Alignment (16-Bit Bus Channel)

Note: The valid bits in the data bus for each data size are indicated by circles (O). Crosses (×) indicate bus data bits that are undefined.

Asterisks (*) indicate write/byte control bits that are disabled (fixed high level).

(c) 8-Bit Bus Channel

If an 8-bit bus is selected by the external bus width select bits in the CSn control register, A27 to A0 are enabled as address signals for byte units. Table 9.9 shows the data alignment corresponding to byte addresses for different data sizes.

With an 8-bit bus channel only the \overline{WRO} pin is enabled, regardless of the strobe mode setting. A low-level signal is output to $\overline{WR0}$ during write access. $\overline{BC0}$ constantly outputs low level. Pins $\overline{WR3}$ to $\overline{WR1}$ and pins $\overline{BC3}$ to $\overline{BC1}$ are not used.

Table 9.9 Data Alignment (8-Bit Bus Channel)

Note: The valid bits in the data bus for each data size are indicated by circles (O). Crosses (x) indicate bus data bits that are undefined. Asterisks (*) indicate write/byte control bits that are disabled (fixed high level).

9.5.2 SDRAM Interface

A description is provided here of the SDRAM controller (SDRAMC) operation enable and SDRAM bus width settings as well as operations involving SDRAM (read, write, auto-refresh, self-refresh, initialization sequence, and mode register settings).

(1) SDRAM Access Enable/Disable and SDRAM Bus Width Settings

Enabling and disabling SDRAM access is performed by making settings in the individual SDRAMCm control registers to enable or prohibit SDRAMC operation. SDRAM bus width settings are also performed by means of the SDRAMCm control registers.

Even if the SDRAMC control register is set to disable SDRAMC operation, refresh operation will still take place if self-refresh or auto-refresh operation is set as enabled.

(2) SDRAM Commands

SDRAMC controls the SDRAM by issuing commands each bus cycle. These commands are defined by combinations of RAS, CAS, WE, CKE, CS, etc.

Table 9.10 lists the commands issued by SDRAMC.

[Legend]

H: High level, L: Low level, V: Valid, X: Don't care

(3) SDRAMC Register Setting Conditions

Rewriting of SDRAMC registers should only be performed when all of the conditions listed in table 9.11 are satisfied.

Notes: 1. After writing 0 to EXENB, check to confirm that the EXENB bit has been cleared to 0. 2. Do not fail to confirm that all status bits in the SDRAM status register (SDSTR) have been cleared to 0 before rewriting this bit.

(4) Self-Refresh

Transition to and from self-refresh mode is controlled by means of settings to SDRAM refresh control register 0 (SDRFCNT0). Transition to and from self-refresh mode takes place simultaneously for all channels.

An auto-refresh cycle operation takes place immediately before transition to self-refresh mode. While in self-refresh mode the CKE signal is low level. Immediately after recovery from selfrefresh mode an auto-refresh cycle is triggered.

Figure 9.7 shows the timing of transition to self-refresh mode, and figure 9.8 shows the timing of recovery from self-refresh mode.

Figure 9.7 Example of Timing of Transition to Self-Refresh Mode (DREFW Bit Set Value: 0010)

Figure 9.8 Example of Timing of Recovery from Self-Refresh Mode (DREFW Bit Set Value: 0010)

(5) Auto-Refresh

An auto-refresh cycle starts when the auto-refresh operation enable bit (DRFEN) in SDRAM refresh control register 1 (SDRFCNT1) is set to 1. After that refresh requests are issued at fixed intervals, activating auto-refresh cycles. However, the activation of auto-refresh cycles may sometimes be delayed because refresh requests are not accepted during read or write accesses.

A refresh request is issued immediately if the auto-refresh operation enable bit (DRFEN) in SDRAM refresh control register 1 (SDRFCNT1) is set to 1 while auto-refresh is enabled.

The refresh counter is halted in self-refresh or deep-power-down mode. After recovery from selfrefresh or deep-power-down mode an auto-refresh cycle is activated, after which the counter value is reset and the counter begins operating again

Make auto-refresh settings in SDRAM refresh control register 1 (SDRFCNT1). Note that refresh cycles affect all SDRAM channels. Figure 9.9 shows an auto-refresh cycle timing example.

Figure 9.9 Auto-Refresh Cycle Timing Example (DREFW Bit Set Value: 0010)

(6) Initialization Sequencer

SDRAMC is provided with a sequencer for issuing the commands for SDRAM initialization. The initialization sequence should always be initiated a single time only following a reset (all channels) and following recovery from deep-power-down mode (individual channels). In such cases operation cannot be guaranteed if the initialization sequence is not performed, or if it is performed more than once.

The SDRAM initialization sequence issues the precharge-all-banks command followed by $n (n = 1)$ to 15) auto-refresh commands, in that order. Make timing settings for the initialization sequencer to SDRAM initialization register 0 (SDIR0). Initialization sequences are initiated using SDRAM initialization register 1 (SDIR1).

Note that an initialization sequence for all channels is initiated using the DINIRQ bit.

Figure 9.10 shows a timing example for the initialization sequence. Setting DARFC to specify two or more times causes multiple initialization auto-refresh cycles to be performed.

Figure 9.10 Initialization Sequence Timing Example (DPC Bit Set Value: 001, DARFI Bit Set Value: 0001, DARFC Bit Set Value: 001)

(7) Power-Down Mode

SDRAMC supports an SDRAM power-down mode. In power-down mode the SDCKE signal from SDRAMC goes low level. While in power-down mode auto-refresh operations are performed at the interval specified by the auto-refresh request interval setting (DRFC) bits in SDRAM refresh control register 1 (SDRFCNT1). The SDCKE signal only goes high when an auto-refresh command is issued.

Transition to and recovery from power-down mode are performed using the SDRAM power-down control register (SDPWDCNT).

Setting the DPWD bit to 1 causes SDRAMC to transition to power-down mode. Clearing the DPWD bit to 0 causes SDRAMC to recover from power-down mode.

The SDCKE signal from SDRAMC goes high level when recovery from power-down mode occurs.

Figure 9.11 SDRAMC Power-Down Mode

Figure 9.12 Auto-Refresh Operation in SDRAMC Power-Down Mode
(8) Deep-Power-Down Mode

SDRAMC supports an SDRAM deep-power-down mode. In deep-power-down mode SDRAMC issues a deep-power-down command and drives the SDCKE signal low level.

Transition to and recovery from deep-power-down mode are performed using the SDRAM deeppower-down control register (SDDPDCNT).

Setting the DDPD bit to 1 causes SDRAMC to put all channels into deep-power-down mode. Clearing the DDPD bit to 0 causes SDRAMC to recover from deep-power-down mode.

During recovery from deep-power-down mode, SDRAMC issues a deep-power-down exit command and drives the SDCKE signal high level.

Following recovery from deep-power-down exit, wait for the duration designated for the SDRAM being used and then execute an initialization sequence.

Figure 9.13 SDRAMC Deep-Power-Down Mode

(9) Read/Write Access

The following two types of read/write access are supported.

- Multiple read/multiple write
- Single read/single write

Multiple read/multiple write occurs in the following cases.

- 1. CPU burst access (cache replace)
- 2. Access with longword (32-bit) to the SDRAM data bus having 8-bit or 16-bit width
- 3. Access with word (16-bit) to the SDRAM data bus having 8-bit width
- 4. Multiple data transfer in DMA pipeline transfer

The access timing can be set independently for each channel using the SDRAMI timing register (SDITR). Access timing examples are described below.

(a) Multiple Read/Multiple Write Access

Figure 9.14 shows a timing example for multiple read of 4 units of data, and figure 9.15 for multiple write of 4 units of data.

The number of DMA transfers performed will vary depending on factors such as the number of transfers and the transfer data size per operand and the SDRAM bus width. Read commands or write commands may or may not be issued consecutively in response to an access request from the bus master. When read commands or write commands are not issued consecutively, a deselect command is issued between them.

Furthermore, deactivation and activation are performed automatically when the SDRAM row address changes during a DMA transfer operation.

Figure 9.16 shows a timing example for multiple read of 4 units of data, and figure 9.17 for multiple write of 4 units of data, when read/write commands are not issued consecutively. Figure 9.18 shows a timing example for multiple write with a row address change.

The access timing is modified by means of settings in the SDRAMm timing register (SDmTR).

Figure 9.14 Multiple Read Timing Example (Multiple Read of 4 Data Units, Shortest Timing Settings) Consecutive Read Commands Issued

Figure 9.15 Multiple Write Timing Example (Multiple Write of 4 Data Units, Shortest Timing Settings) Consecutive Write Commands Issued

Figure 9.16 Multiple Read Timing Example (Multiple Read of 4 Data Units, Shortest Timing Settings) Non-Consecutive Read Commands Issued

Figure 9.17 Multiple Write Timing Example (Multiple Write of 4 Data Units, Shortest Timing Settings) Non-Consecutive Write Commands Issued

Figure 9.18 Multiple Write Timing Example (Multiple Write of 4 Data Units, Shortest Timing Settings) Access Spanning Rows

(b) Single Read/Single Write Access

Figure 9.19 shows a timing example for single read operation and figure 9.20 for single write operation. The access timing is modified by means of settings in the SDRAMm timing register (SDmTR).

Figure 9.19 Single Read Timing Example (Shortest Timing Settings)

Figure 9.20 Single Write Timing Example (Shortest Timing Settings)

(c) Byte Access Control by DQM

Figures 9.21 and 9.22 show timing examples for byte accesses to the SDRAM with a 16-bit bus width. In the SDRAM access, the DQM signal is asserted when data is masked.

Figure 9.21 Byte Write Timing to SDRAM with 16-Bit Bus Width (Example)

Figure 9.22 Byte Read Timing from SDRAM with 16-Bit Bus Width (Example)

(10) Mode Register Setting

Writing to the SDRAMm mode register (SDmMOD) causes mode register set commands and extended mode register set commands to be issued to the various channels. Settings to the SDRAMm mode register (SDmMOD) should be made individually for each channel.

Figure 9.23 shows the operation timing for mode register setting.

Figure 9.23 Operation Timing for Mode Register Setting

(11) Clock Stop Control Signal

SDRAMC outputs a clock stop control signal (CLKSTOP). CLKSTOP can be enabled or disabled using the DCKSEN bit in the SDRAM clock stop control signal setting register (SDCKSCNT).

The CLKSTOP signal remains low level when the clock stop control signal is disabled.

When clock stop control signal is enabled, the CLKSTOP and CKIO signals operate in conjunction with transition to and recovery from deep-power-down mode.

During a transition to deep-power-down mode, the CLKSTOP signal goes high after the deeppower-down entry command is issued. During a recovery from deep-power-down mode, the CLKSTOP signal goes low and a deep-power-down exit command is issued when the clearing of the DDPD bit to 0 is accepted by SDRAMC and the CKIO starts operation.

DCKSC, the period between the change of CLKSTOP along with CKIO and the issuance of deep power-down entry or exit command, can be set by the SDRAM clock stop control signal setting register.

Figures 9.24 and 9.25 show the operation timing of the clock stop control signal.

(12) SDRAMC Setting Examples

The SDRAMC setting procedure, timing register setting examples, and the procedure for transitioning to and recovering from self-refresh mode, power-down mode, and deep-power-down mode are described below.

(a) SDRAMC Setting Procedure

Figure 9.26 shows the SDRAMC setting procedure.

Note that the specifications of the power-up sequence, etc., may vary depending on the SDRAM used. Study the SDRAM specifications carefully before making system settings.

Figure 9.26 SDRAMC Setting Procedure

(b) Procedure for Transition to and Recovery from Self-Refresh Mode

Figure 9.27 shows the procedure for transitioning to and recovering from self-refresh mode.

Figure 9.27 Procedure for Transition to and Recovery from Self-Refresh Mode

- Note: Before transitioning to or recovering from self-refresh mode it is necessary to halt SDRAM access to the affected channels. Consequently, it is not possible to transition to or recover from self-refresh mode while programs or DMA operations that access SDRAM are in progress. Pay attention to the following points when writing programs.
	- Before transitioning to self-refresh mode, halt any DMA channel transfers that access the SDRAM area of the affected channels.
	- Make sure that programs run while transitioning to self-refresh mode, while in selfrefresh mode, or while recovering from self-refresh mode do not access operands or fetch (or pre-fetch) instructions stored in the SDRAM area.

(c) Procedure for Transition to and Recovery from Deep-Power-Down Mode

Figure 9.28 shows the procedure for transitioning to deep-power-down mode.

Figure 9.28 Procedure for Transition to Deep-Power-Down Mode

Figure 9.29 shows the procedure for recovering from deep-power-down mode.

Figure 9.29 Procedure for Recovery from Deep-Power-Down Mode

Note: Before transitioning to or recovering from deep-power-down mode it is necessary to halt SDRAM access to the affected channels. Consequently, it is not possible to transition to or recover from deep-power-down mode while programs or DMA operations that access SDRAM are in progress. Pay attention to the following points when writing programs.

- Before transitioning to deep-power-down mode, halt any DMA channel transfers that access the SDRAM area of the affected channels.
- Make sure that programs run while transitioning to deep-power-down mode, while in deep-power-down mode, or while recovering from deep-power-down mode do not access operands or fetch (or pre-fetch) instructions stored in the SDRAM area.

(d) Timing Register Set Values and Access Timing

The correspondence between the SDRAMm timing register (SDmTR) set values and the read and write access timing is described below.

• Multiple Read Timing Setting Examples

Figures 9.30 to 9.32 show the correspondence between the timing of multiple read operations involving 4 data units and the set values of the SDRAMm timing register (SDmTR). Table 9.12 shows the SDRAMm timing register (SDmTR) set values for each figure.

Table 9.12 SDITR Set Value Correspondence Table (Multiple Read Timing)

Figure	DRAS	DRCD	DPCG	DCL
Figure 9.30	010	00	001	010
Figure 9.31	000		001	010
Figure 9.32	000		001	011

Figure 9.30 Multiple Read Timing Example 1

Figure 9.31 Multiple Read Timing Example 2

Figure 9.32 Multiple Read Timing Example 3

• Multiple Write Timing Setting Examples

Figures 9.33 to 9.35 show the correspondence between the timing of multiple write operations involving 4 data units and the set values of the SDRAMm timing register (SDmTR). Table 9.13 shows the SDRAMm timing register (SDmTR) set values for each figure.

Table 9.13 SDITR Set Value Correspondence Table (Multiple Write Timing)

Figure 9.33 Multiple Write Timing Example 1

Figure 9.35 Multiple Write Timing Example 3

• Single Read Timing Setting Examples

Figures 9.36 to 9.38 show the correspondence between the timing of single read operations and the set values of the SDRAMm timing register (SDmTR). Table 9.14 shows the SDRAMm timing register (SDmTR) set values for each figure.

Table 9.14 SDITR Set Value Correspondence Table (Single Read Timing)

Figure 9.36 Single Read Timing Example 1

Figure 9.37 Single Read Timing Example 2

Figure 9.38 Single Read Timing Example 3

• Single Write Timing Setting Examples

Figures 9.39 to 9.41 show the correspondence between the timing of single write operations and the set values of the SDRAMm timing register (SDmTR). Table 9.15 shows the SDRAMm timing register (SDmTR) set values for each figure.

Table 9.15 SDITR Set Value Correspondence Table (Single Write Timing)

Figure 9.39 Single Write Timing Example 1

Figure 9.40 Single Write Timing Example 2

Figure 9.41 Single Write Timing Example 3

(13) External Address/SDRAM Address Signal Multiplex

(a) Address Multiplex

Either of addresses used for accessing external device or SDRAM is output through external address pins. The SDRAM address is shifted internally by changing the settings of DDBW and DSZ in SDmADR and BSIZE in SDCmCNT. The bank address is output on A16 and A15 and the address is output on A14 to A2.

Table 9.16 External Address/SDRAM Address Pins

(b) Address Register Setting Value and Supported SDRAM Configuration

Tables 9.17 to 9.19 are the SDRAM configurations that to support for 8-, 16-, or 32-bit bus width. These tables are featured to ease the understanding of the relationships between the SDRAM to support and address multiplex.

Addresses addr27 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM. The table below shows how the settings of DSZ and DDBW determine which signals are output on the SDRAM-access pins.

SDRAM

Table 9.17 Case for 8-Bit External Data Bus Width $(BSIZE^* = (1, 0))$

Notes: 1. The legend BSIZE represents the BSIZE bit in the SDRAMCm control register.

- 3. The legend DDBW represents the DDBW bit in the SDRAMm address register.
- 4. The legends BA1, BA0, and MA12 to MA0 represent the SDRAM bank address and SDRAM address respectively.
- 5. Addresses addr25 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM.
- 6. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

Table 9.18 Case for 16-Bit External Data Bus Width $(BSIZE^{*1} = (0, 0))$ (1)

Notes: 1. The legend BSIZE represents the BSIZE bit in the SDRAMCm control register.

- 3. The legend DDBW represents the DDBW bit in the SDRAMm address register.
- 4. The legends BA1, BA0, and MA12 to MA0 represent the SDRAM bank address and SDRAM address respectively.
- 5. Addresses addr24 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM.
- 6. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

SDRAM

Table 9.18 Case for 16-Bit External Data Bus Width $(BSIZE^{*1} = (0, 0))$ (2)

Notes: 1. The legend BSIZE represents the BSIZE bit in the SDRAMCm control register.

- 3. The legend DDBW represents the DDBW bit in the SDRAMm address register.
- 4. The legends BA1, BA0, and MA12 to MA0 represent the SDRAM bank address and SDRAM address respectively.
- 5. Addresses addr26 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM.
- 6. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

Notes: 1. The legend BSIZE represents the BSIZE bit in the SDRAMCm control register.

- 2. The legend DSZ represents the DSZ bit in the SDRAMm address register.
- 3. The legend DDBW represents the DDBW bit in the SDRAMm address register.
- 4. The legends BA1, BA0, and MA12 to MA0 represent the SDRAM bank address and SDRAM address respectively.
- 5. Addresses addr24 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM.
- 6. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

Table 9.19 Case for 32-Bit External Data Bus Width $(BSIZE^{\ast 1} = (0, 1))$ (2)

Notes: 1. The legend BSIZE represents the BSIZE bit in the SDRAMCm control register.

- 3. The legend DDBW represents the DDBW bit in the SDRAMm address register.
- 4. The legends BA1, BA0, and MA12 to MA0 represent the SDRAM bank address and SDRAM address respectively.
- 5. Addresses addr25 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM.
- 6. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

Notes: 1. The legend BSIZE represents the BSIZE bit in the SDRAMCm control register.

- 2. The legend DSZ represents the DSZ bit in the SDRAMm address register.
- 3. The legend DDBW represents the DDBW bit in the SDRAMm address register.
- 4. The legends BA1, BA0, and MA12 to MA0 represent the SDRAM bank address and SDRAM address respectively.
- 5. Addresses addr27 to addr0 are the logical addresses used by the CPU and DMAC in access to the SDRAM.
- 6. When the RD, WR or PRA command is issued, this carries the pre-charge option signal.

(c) Example of SDRAM Connection

Figures 9.42 and 9.43 show examples of the connection of SDRAM with this LSI.

Figure 9.42 Example of Connecting a 32-Bit Data-Width SDRAM

Figure 9.43 Example of Connecting a 16-Bit Data-Width SDRAM

9.6 Usage Note

9.6.1 Note on Power-on Reset Exception Handling and Deep Standby Mode Cancellation

When writing to the external address space or making SDRAM settings in power-on reset exception handling or cancellation of deep standby mode, be sure to set bits ACOSW[3:0] in ACSWR to B'0011 beforehand.

9.6.2 Write Buffer

In write access to normal or SDRAM space, the write data are stored once in the internal write buffer of the BSC, and only after that does actual writing to the device (external device) connected in the normal or SDRAM space proceed. Since writing from the write buffer to the external device is performed automatically, no processing by software is necessary.

However, care must be taken on the following point. Write access from the CPU or DMAC appears complete at the point where the data are stored in the above write buffer. That is, at the point where the write access from the CPU or DMA controller has been completed, writing to the external device might not have been completed. To confirm the completion of writing to the external device, dummy read the normal or SDRAM space. Completion of the dummy-read operation guarantees the completion of writing to the external device in response to previous write access. The target address for the dummy read operation does not have to be in the same device as the target for write access. Furthermore, it does not have to be in the same space.

9.6.3 Note on Transition to Software Standby Mode or Deep Standby Mode

When a transition to software standby mode or deep standby mode is made after write access to the normal or SDRAM space, there is a possibility that data remains in the internal write buffer of the BSC. To confirm that no data remain in the write buffer, execute a dummy read of the external device in the same way as described above.

Section 10 Bus Monitor

The bus monitor is a module that monitors bus errors on each bus. When an illegal address access or a bus timeout is detected, a bus error interrupt is generated and an access canceling signal is output for the bus timeout. (The bus timeout function is used for debugging.)

Figure 10.1 shows a block diagram of the bus monitor.

Figure 10.1 Block Diagram of Bus Monitor

10.1 Register Descriptions

The bus monitor has the following registers.

All registers are initialized by a power-on reset or in deep standby mode.

Table 10.1 Register Configuration

10.1.1 Bus Monitor Enable Register (SYCBEEN)

SYCBEEN clears the bus monitor status register and controls the detection function.

Note: When a bus access is performed with the detection function disabled (TOEN = 0), the bus may freeze.

10.1.2 Bus Monitor Status Register 1 (SYCBESTS1)

SYCBESTS1 indicates the status of slave buses (peripheral bus (1)/peripheral bus (3)) regarding whether a timeout occurred, whether an illegal address access was made, or which bus master accessed the slave bus. Table 10.2 shows the correspondence between the bus spaces and the slave buses.

Table 10.2 Bus Space and Slave Bus

Notes: 1. This means bus spaces in the slave bus space other than those for the external bus and peripheral buses (1), (2), and (3).

2. An illegal address access error does not occur.

10.1.3 Bus Monitor Status Register 2 (SYCBESTS2)

SYCBESTS2 indicates the status of slave buses (external bus/peripheral bus (2)/others) regarding whether a timeout occurred, whether an illegal address access was made, or which bus master accessed the slave bus.

Section 10 Bus Monitor

10.1.4 Bus Error Control Register (SYCBESW)

SYCBESW controls the notification of various types of bus errors to the CPU.

10.2 Bus Monitor Function

The bus monitor function detects two types of bus error: illegal address access and bus timeout. Bus error detection is performed in one bus access.

Even when data is transferred in multiple bus accesses such as burst transfer, a bus error can be detected in one bus access.

10.2.1 Operation when a Bus Error is Detected

When a bus error is detected, the status is saved in the bus monitor status register 1 (SYCBESTS1) and bus monitor status register 2 (SYCBESTS2) and the CPU is notified of the bus error is notified to the CPU.

(1) Saving Status in Bus Monitor Status Register or Bus Monitor Status Register 2

When a bus error occurs, the status at the time (what type of error occurred and which bus was being accessed by which bus master) is saved in the bus monitor status register 1 (SYCBESTS1) or bus monitor status register 2 (SYCBESTS2).

Even if another bus error occurs after this, the value in the bus monitor status register (SYCBESTS) or bus monitor status register 2 (SYCBESTS2) is not updated. When multiple bus errors occur at the same time, multiple status bits may be set.

The bus monitor status register 1 (SYCBESTS1) or bus monitor status register 2 (SYCBESTS2) can be cleared by writing 1 to the status clear bit (STSCLR) in the bus monitor enable register (SYCBEEN) from the bus master. After being cleared, the status of a bus error, if generated, is saved in the bus monitor status register 1 (SYCBESTS1) or bus monitor status register 2 (SYCBESTS2) again.

When a clear operation and a bus error happen at the same time, the clear operation has priority and the bus error is ignored.

(2) Error Notification to the CPU

The CPU is notified of a bus error through the OR condition of the timeout bits (PTO/CTO/ETO) and illegal address access bits (PER/CER/EER/OER/SHER) in the bus monitor status register 1 (SYCBESTS1) and bus monitor status register 2 (SYCBESTS2). The CPU is notified of a bus error interrupt according to the setting of the bus error control register (SYCBESW).

When the bus monitor status register 1 (SYCBESTS1) and bus monitor status register 2 (SYCBESTS2) are cleared by the CPU, the bus error interrupt signal is also negated.

(3) Termination of Bus Access

When a bus error is detected, the bus access is terminated. For details, see section 10.2.4, Combinations of Masters and Bus Errors.

For the detailed operations when each type of error is detected, see section 10.2.2, Illegal Address Access Detection Function and section 10.2.3, Bus Timeout Detection Function.

10.2.2 Illegal Address Access Detection Function

The illegal address access detection function detects attempted accesses to illegal addresses.

(1) Conditions of Illegal Address Access Error Generation

Illegal address access errors occur when the following illegal addresses are accessed.

- External spaces for which the operation enable bit (EXENB) in the control register of the BSC is not set to "operation enabled"
- Other address areas that are not mapped to any slave bus
- Address areas that are mapped to the slave buses but do not correspond to slave devices

Tables 10.3 to 10.5 show the address areas to which slave devices are not mapped within the spaces for peripheral buses (1) , (2) , and (3) .

Table 10.3 Address Areas without Slave Devices in the Space for Peripheral Bus (1)

Table 10.4 Address Areas without Slave Devices in the Space for Peripheral Bus (2)

Table 10.5 Address Areas without Slave Devices in the Space for Peripheral Bus (3)

10.2.3 Bus Timeout Detection Function

The bus timeout detection function detects bus accesses whose cycles are extended to 768 cycles or more.

(1) Conditions of Bus Timeout Error Generation

Bus timeout errors occur in the following cases. This function should be used when debugging software.

- A bus access is not completed on peripheral bus (1)
- A bus access is not completed on peripheral bus (3)
- The WAIT signal remains asserted during an external bus access

(2) Operation When a Bus Timeout Error is Generated

The operation when a bus timeout error occurs is explained below.

- 1. The timeout counter starts counting from the next cycle after the start of a bus access.
- 2. If the bus access is not completed in 768 cycles, a bus timeout occurs and an access canceling signal is asserted for 256 cycles.

Bus signals such as address, data, BC, read/write, and burst are held.

 The timeout error is recorded in the bus monitor status register 1 (SYCBESTS1) or bus monitor status register 2 (SYCBESTS2).

A bus error interrupt is generated and sent to the CPU.

- 3. The bus access is terminated.
- 4. The CPU processes the bus error.

Locked buses are all released.

(3) Bus Timeout Operation in Consecutive Accesses

For transfers where multiple bus accesses are made (such as burst transfer), the next bus access might not be terminated when a bus timeout occurs. In this case, a bus timeout may occur continuously.

Even if a bus timeout occurs continuously, the timeout process of terminating a bus access is performed in the same way as the first time. However, the status is saved in the bus monitor status register 1 (SYSCESTS1) or bus monitor status register 2 (SYCBESTS2) only the first time.

10.2.4 Combinations of Masters and Bus Errors

The types of detectable bus error depend on the master and access mode.

(1) CPU Transfer Modes and Types of Bus Error Generated

Table 10.6 shows the types of bus error that may be generated by accesses from the CPU.

[Legend]

O: A bus error is generated.

 \equiv : A bus error is not generated.

Notes: 1. To enable bus error detection, the bus monitor enable register (SYCBEEN) should be set.

 2. To notify the CPU of a bus error, the 00CPEN bit in the bus error control register (SYCBESW) should be set to 1.

 3. The number of bus errors detected is the same as the number of accesses that resulted in an error.

(2) DMAC Transfer Modes and Operations of Each Bus

Table 10.7 shows the DMAC transfer modes and the types of bus error that may be generated by accesses from the DMAC.

[Legend]

O: A bus error is generated.

 \equiv : A bus error is not generated.

Note: * To enable bus error detection, the bus monitor enable register (SYCBEEN) should be set.

10.3 Usage Note

10.3.1 Operation when the CPU is Not Notified of a Bus Error

Table 10.8 describes the operations when bus error notification to the CPU is disabled with the bus error detection enabled (by the setting of the bus monitor enable register (SYCBEEN)).

Section 11 Direct Memory Access Controller (DMAC)

The DMA controller (hereafter DMAC) is a module that handles high-speed data transfer without CPU intervention in response to requests from software, on-chip peripheral I/O modules, or external pins (external modules). The DMAC itself does not distinguish between requests from on-chip peripheral I/O or external pins (external modules). The DMA supports data transfer between memory units, memory and I/O modules, and I/O modules.

11.1 Features

- Channel number: Up to eight channels (with four channels capable of external requests)
- Transfer requests: Requests from 38 sources including software trigger, on-chip peripheral I/O, and external pins (external modules)
- Maximum transfer bytes: 64 Mbytes
- Address space: 4 Gbytes
- Transfer data sizes:
	- \sim Single data transfer: 8, 16, 32, 64, and 128 bits
	- $-$ Single operand transfer: 1, 2, 4, 8, 16, 32, 64, and 128 data
	- ⎯ Non-stop transfer: Up to the byte count "0"
- Transfer mode:
	- Cycle-stealing transfer (dual-address transfer)
	- Pipelined transfer (dual-address transfer)
- Maximum transfer speed:
	- ⎯ Cycle-stealing transfer: Minimum of three clock cycles per unit data transfer
	- ⎯ Pipelined transfer: Minimum of one clock cycle per unit data transfer
- Transfer conditions:
	- ⎯ Unit operand transfer: a single sequence of single operand data transfer in response to a DMA request
	- ⎯ Sequential operand transfer: single operand transfers are repeated until the byte count reaches "0"
	- ⎯ Non-stop transfer: data is continuously transferred until the byte count reaches "0" in response to a single DMA request
- Channel priority:

Channel 0 > channel $1 > \rightarrow$ > channel 6 > channel 7 (this priority order is fixed)

- Interrupt request
	- ⎯ Two types of interrupt requests (generated when the byte count reaches "0")
		- Interrupt request signals for each channel
		- Interrupt request signal common to all channels
- Reload function (source address, destination address, byte count) settable
- Rotate function settable
- DMAC stop/restart/suspend function settable

Notes: Terminologies in this section are as follows:

- 1. Single data transfer: Transfer in one read cycle and one write cycle by the DMAC (in the case of dual address transfer)
- 2. Single operand transfer: Continuous data transfer by the DMAC on one channel (amount of data to be transferred is set in a register)
- 3. One DMA transfer: Transferring a number of data, from the start address to the end address set in the byte count register
- 4. Channel number: $n = 0$ to 7
- 5. Request source number: $k = 1$ to 37, m = 0 to 37
- 6. BIU: Bus Interface Unit (peripheral module). One of the following four kinds according to the source or destination of transfer.
	- BIU_E: External space (normal space and SDRAM space)
	- BIU_P: Peripheral bus (1) (see figure 1.1)
	- BIU SH: Peripheral bus (2) (see figure 1.1), on-chip RAM space
	- BIU_C: Peripheral bus (3) (see figure 1.1)

Figure 11.1 is a block diagram of the DMAC

Figure 11.1 DMAC Block Diagram

11.2 Input/Output Pins

Table 11.1 Pin Configuration

11.3 Register Descriptions

The DMAC has the following registers. All registers are initialized by a power-on reset or in deep standby mode.

Table 11.2 Register Configuration

11.3.1 DMA Current Source Address Register (DMCSADR)

DMCSADR is a register used to specify the start address of the transfer source.

The value in this register is transferred to the working source-address register at the start of DMA transfer.

The default behavior is for the contents of the working source-address register to be returned on completion of single operand transfer. However, the contents of the working source address register are not returned in two cases: when the rotate setting $(SAMOD = 011)$ is made for the source address and when the source-address reload function is enabled. In the latter case, the contents of the DMA reload source address register (DMRSADRn) are returned to this register on completion of DMA transfer.

This register must be set before transfer is initiated, regardless of whether the reload function is enabled or disabled.

Notes: 1. Set this register so that DMA transfer is performed within the correctly aligned address boundaries for the transfer sizes listed below.

- When the transfer size is set to 16 bits $(SZSEL = "001")$: $(b0) = "0"$.
- When the transfer size is set to 32 bits $(SZSEL = "010")$: $(b1, b0) = (0, 0)$.
- 2. Only write to this register when single operand transfer is not in process on the corresponding channel (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is "0") and DMA transfer is disabled (DMST in the DMA activation control register (DMSCNT) or DEN in DMA control register B for the channel (DMCNTBn) is set to "0"). Operation is not guaranteed if this register is written to when both conditions are not satisfied.

11.3.2 DMA Current Destination Address Register (DMCDADR)

DMCDADR is a register used to specify the start address of the transfer destination.

The value in this register is transferred to the working destination-address register at the start of DMA transfer.

The default behavior is for the contents of the working destination-address register to be returned on completion of each single operand transfer. However, the contents of the working destinationaddress register are not returned in two cases: when the rotate setting $(SAMOD = 011)$ is made for the destination address and when the destination-address reload function is enabled. In the latter case, the contents of the DMA reload destination address register (DMRDADRn) are returned to this register on completion of DMA transfer.

This register must be set before transfer is initiated, regardless of whether the reload function is enabled or disabled.

Notes: 1. Set this register so that DMA transfer is performed within the correctly aligned address boundaries for the transfer sizes listed below.

- When the transfer size is set to 16 bits $(SZSEL = "001")$: $(b0) = "0"$.
- When the transfer size is set to 32 bits (SZSEL = $"010"$): (b1, b0) = (0, 0).
- 2. Only write to this register when single operand transfer is not in process on the corresponding channel (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is "0") and DMA transfer is disabled (DMST in the DMA activation control register (DMSCNT) or DEN in DMA control register B for the channel (DMCNTBn) is set to "0"). Operation is not guaranteed if this register is written to when both conditions are not satisfied.

11.3.3 DMA Current Byte Count Register (DMCBCT)

DMCBCT is a register used to specify the number of bytes to be transferred by DMA.

The value in this register is transferred to the working byte-count register at the start of DMA transfer, and is then decremented by the number of bytes transferred on each unit data transfer. Decrementation is thus by the following values.

- When the transfer size is set to 8 bits ($SZSEL = "000"$): -1
- When the transfer size is set to 16 bits (SZSEL = "001"): -2
- When the transfer size is set to 32 bits (SZSEL = $"010"$): -4

When the value in the working byte count register reaches H'000 0000, DMA transfer ends (an end due to byte count "0"). The corresponding bit of the DMA transfer end detection register (DMEDET) is set to 1.

If the byte count reload function is disabled, the contents of the working byte count register are returned to this register at the moment the channel for DMA transfer switches or DMA transfer ends. If the byte count reload function is enabled, the contents of the DMA reload byte counter register (DMRBCTn) are returned to this register.

This register must be set before transfer is initiated, regardless of whether the reload function is enabled or disabled.

Notes: 1. Note that a setting of H'000 0000 leads to transfer of the maximum number of bytes, i.e. 64 Mbytes.

- 2. Set this register so that DMA transfer is performed within the correctly aligned address boundaries for the transfer sizes listed below.
	- When the transfer size is set to 16 bits (SZSEL = $"001"$): (b0) = $"0"$.
	- When the transfer size is set to 32 bits $(SZSEL = "010")$: $(b1, b0) = (0, 0)$.
- 3. Only write to this register when single operand transfer is not in process on the corresponding channel (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is "0") and DMA transfer is disabled (DMST in the DMA activation control register (DMSCNT) or DEN in DMA control register B for the channel (DMCNTBn) is set to "0"). Operation is not guaranteed if this register is written to when both conditions are not satisfied.

11.3.4 DMA Reload Source Address Register (DMRSADR)

DMRSADR is used to set an address for reloading to the DMA current source address register (DMCSADRn).

To enable reloading, set the DMA source address reload function enable bit (SRLOD) in DMA control register A (DMCNTAn) for the channel to "1". In this case, set both the DMA current source address register (DMCSADRn) and DMA reload source address register (DMRSADRn).

boundaries for the transfer sizes listed below.

- When the transfer size is set to 16 bits (SZSEL = "001"): (b0) = "0".
- When the transfer size is set to 32 bits $(SZSEL = "010")$: $(b1, b0) = (0, 0)$.

11.3.5 DMA Reload Destination Address Register (DMRDADR)

DMRDADR is a register used to set an address for reloading to the DMA current destination address register (DMCDADRn).

To enable reloading, set the DMA destination address reload function enable bit (DRLOD) in DMA control register A (DMCNTAn) to 1. In this case, set both the DMA current destination address register (DMCDADRn) and DMA reload destination address register (DMRDADRn).

• When the transfer size is set to 16 bits $(SZSEL = "001")$: $(b0) = "0"$.

• When the transfer size is set to 32 bits $(SZSEL = "010")$: $(b1, b0) = (0, 0)$.

11.3.6 DMA Reload Byte Count Register (DMRBCT)

DMRBCT is a register used to set the number for reloading to the DMA current byte count register (DMCBCTn).

To enable reloading, set the DMA byte count reload function enable bit (BRLOD) in the DMA control register A (DMCNTAn) to 1. In this case, set both the DMA current byte count register (DMCBTn) and DMA reload byte count address register (DMRBCTn).

Note: Set this register so that DMA transfer is performed within the correctly aligned address boundaries for the transfer sizes listed below.

- When the transfer size is set to 16 bits $(SZSEL = "001")$: $(b0) = "0"$.
- When the transfer size is set to 32 bits $(SZSEL = "010")$: $(b1, b0) = (0, 0)$.

11.3.7 DMA Mode Register (DMMOD)

DMMOD controls the amount of data, data size selection, address direction, and various types of signal outputs.

Note: Only write to this register when the corresponding channel is not engaged in single operand transfer (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is "0") and DMA transfer is disabled (DMST in the DMA activation control register (DMSCNT) or DEN in DMA control register B for the channel (DMCNTBn) is set to "0"). Operation is not guaranteed if this register is written to when both conditions are not satisfied. When SACT and DACT are set to 1, output of a low DACT signal from the cycle following a DMAC read or write cycle is enabled.

Table 11.3 shows the DMA source/destination address registers. For details on the rotation address "indexing" mode, see section 11.11, Rotate Function. Note that when performing pipelined transfer to or from external devices and modules that support burst access, make sure to set the direction bits to select address incrementation ("001") or rotation ("011").

Table 11.3 Increment/Decrement for DMA Source/Destination Address Registers

11.3.8 DMA Control Register A (DMCNTA)

DMCNTA handles the selections of the transfer mode and the condition of transfer, control of reload functions, and selection of DMA sources.

Note: Only write to bits of this register other than the reload function enable bits (BRLOD, SRLOD, and DRLOD) when a transfer operation is not in process on the corresponding channel (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is "0") and DMA transfer is disabled (DMST in the DMA activation control register (DMSCNT) or DEN in DMA control register B for the channel (DMCNTBn) is set to "0"). Operation is not guaranteed if this register is written to when both conditions are not satisfied.

Table 11.4 Relationships between DMA Request Sources and Input Sense Mode

STRG Bit Settings

[Legend]

 \times : Setting prohibited $\sqrt{ }$: Can be set

√: Can be set

11.3.9 DMA Control Register B (DMCNTB)

DMCNTB enables or disables DMA transfer, clears the DMA transfer enable bit, and also clears the internal state. In addition, this register can check the status of a DMA request.

Note: When the software trigger is selected as the DMA request source, the DMA request bit (DREQ) can be set to "1" regardless of the settings of the DMA transfer enable bit (DEN) and DMAC module activation bit (DMST) and whether or not a transfer operation is currently in progress. However, even if the software trigger is selected as the DMA request source, only clear the DMA request bit (DREQ) to "0" or write to the DMAC internal state clearing bit (DSCLR) when a transfer operation is not in process on the corresponding channel (the corresponding DASTS bit in the DMA arbitration status register (DMASTS) is "0") and DMA transfer has been disabled (DMST in the DMA activation control register (DMSCNT) or DEN in the DMA control register B (DMCNTBn) is set to "0"). Operation is not guaranteed if this register is written to when these conditions are not satisfied.

11.3.10 DMA Activation Control Register (DMSCNT)

DMSCNT controls the operation of the DMAC.

11.3.11 DMA Interrupt Control Register (DMICNT)

DMICNT controls DMA interrupts for the respective channels.

Note: Bits 31 to 24 correspond to channels 0 to 7, respectively (31: channel 0, 30: channel 1. …24: channel 7).

11.3.12 DMA Common Interrupt Control Register (DMICNTA)

DMICNTA determines which channels contribute to the output of a common interrupt request signal.

Note: Bits 31 to 24 correspond to channel 0 to 7, respectively (31: channel 0, 30: channel 1, …, 24: channel 7).

11.3.13 DMA Interrupt Status Register (DMISTS)

DMISTS consists of the DMA interrupt request status bits.

Notes: 1. This register is read-only.

 2. Bits 31 to 24 correspond to channels 0 to 7, respectively (31: channel 0, 30: channel 1, …, 24: channel 7).

11.3.14 DMA Transfer End Detection Register (DMEDET)

DMEDET verifies the status of DMA transfer end detection for each channel. Writing 0 to the DEDET bit is invalid and 1 written to the bit is not retained.

Note: Bits 31 to 24 correspond to channels 0 to 7, respectively (31: channel 0, 30: channel 1, …, 24: channel 7).

11.3.15 DMA Arbitration Status Register (DMASTS)

DMASTS verifies the status of DMA transfer on each channel. Writing 0 to the DASTS bit is invalid and 1 written to the bit is not retained.

Note: Bits 31 to 24 correspond to channels 0 to 7, respectively (31: channel 0, 30: channel 1, …, 24: channel 7)

11.4 Operation

11.4.1 DMA Transfer Mode

There are two DMA transfer modes — cycle-stealing mode and pipelined mode. These modes are selectable through the setting of the DMA transfer mode select bits (MDSEL) in DMA Control Register A (DMCNTAn).

Figure 11.2 gives examples of how bus mastership alternates between the DMAC and CPU in various DMA transfer modes.

(1) Cycle-stealing Transfer Mode

Cycle-stealing transfer mode is selected when the DMA transfer mode select bits are set to "00".

In cycle-stealing transfer mode, the DMAC leaves at least one cycle between the read and write access cycles of each single data transfer. During this interval, the CPU can access the same target BIU as the source or destination of its own operations. For details on the BIU, see section 11.1, Features.

(2) Pipelined Transfer Mode

Pipelined transfer mode is selected when the DMA transfer mode select bits are set to "01".

In pipelined transfer mode, DMAC activates the bus for read or write access, or both, on consecutive cycles. Therefore, the CPU cannot access the target BIU as a source or destination during single operand transfer.

Pipelined transfer through a single BIU is not possible either.

Figure 11.2 Examples of the Alternation of Bus Mastership between the DMAC and CPU in Various DMA Transfer Modes

11.4.2 DMA Transfer Condition

There are three methods of DMA transfer — the unit transfer operation, sequential operand transfer, and non-stop transfer. These are selectable through the setting of the DMA transfer condition selection bits (DSEL) in DMA Control Register A (DMCNTAn). Each of the conditions is explained below. Table 11.5 and figure 11.3 are a list and chart of the DMA transfer conditions.

(1) Unit Operand Transfer

Setting the DMA transfer condition selection bits (DSEL) to 00 selects this mode. A single DMA request initiates continuous transfer of the number of bytes selected by the OPSEL bits in the DMA mode register. If the byte counter does not reach 0 in single operand transfer, the DMA transfer is completed by repeating unit transfer operations until the byte counter does reach 0.

In the case that the DMA transfer condition is the unit operand transfer and the input sense mode of DMA request is the level sense, there is the mask period of the DMA request in the channel arbitration period after one operand transfer end (please refer to section 11.7.3, Sense Mode for DMA Requests for details). Therefore, in the channel arbitration period after one operand transfer end, in the case that there is no DMA request of the higher-priority channel than the transferring channel and there is the DMA request of the lower-priority channel than the transferring channel, the DMA transfer of the low-priority channel starts. To execute the DMA transfer of the highpriority channel in succession, please set the DMA transfer condition to the sequential operand transfer or the non-stop transfer.

(2) Sequential Operand Transfer

Setting the DMA transfer condition selection bits (DSEL) to 01 selects this mode. A single DMA request initiates transfer in units of the number of bytes selected by the OPSEL bits in the DMA mode register (i.e., unit transfer operations) until the DMA transfer is complete (i.e., until the byte counter reaches zero). Channel arbitration is performed on completion of each unit transfer operation. Transfer on the channel for the sequential operand transfer automatically resumes unless there is a DMA request from a higher-priority channel.

In the case that the DMA transfer condition is the sequential operand transfer, even if the input sense mode of DMA request is the level sense, there is no mask period before the byte count becomes 0. Therefore, the DMA transfer of the low-priority channel than the transferring channel cannot start.

(3) Non-Stop Transfer

Setting the DMA transfer condition selection bits (DSEL) to 11 selects this mode. A single DMA request initiates DMA transfer that continues until the transfer is complete (i.e., until the byte counter reaches zero). There are no gaps for channel arbitration, so even DMA requests from highpriority channels will not be accepted.

Table 11.5 List of DMA Transfer Conditions

Figure 11.3 DMA Transfer Conditions

Relations between the mode and conditions of DMA transfer are shown in table 11.6.

Table 11.6 Relations between the mode and conditions of DMA transfer.

Note: * The restriction means that non-stop transfer to the external SDRAM in pipelined transfer mode cannot be set up.

11.4.3 DMA Activation

(1) Initial Settings of the DMAC

Initial settings must be made in each of the relevant registers before the DMA transfer enable bit is set (DEN = "1"). These settings cannot be changed once transfer has started.

An example of DMAC registers that require initial settings is given below.

- 1. DMA mode register (DMMODn)
- 2. DMA control register A (DMCNTAn)
- 3. DMA control register B (DMCNTBn)
- 4. DMA current source address register (DMCSADRn)
- 5. DMA reload source address register (DMRSADRn) when the reload function is used
- 6. DMA current destination address register (DMCDADRn)
- 7. DMA reload destination address register (DMRDADRn) when the reload function is used
- 8. DMA current byte count register (DMCBCTn)
- 9. DMA reload byte count register $(DMRBCTn)$ when the reload function is used
- 10. DMA interrupt control register (DMICNT) when an interrupt is used
- 11. DMA common interrupt control register (DMICNTA) when an interrupt is used
- 12. DMA transfer enable bit (DEN)
- 13. DMA activation control register (DMSCNT)

(2) DMA Activation

DMA transfer for a channel is enabled by setting the DMA transfer enable bit (DEN) in DMA control register B for the channel and the DMAC module activation bit (DMST) in the DMAC activation register (DMSCNT) to "1".

When multiple DMA transfer requests are present, there is no complex mechanism for the determination of channel priority. The DMA request that corresponds to the highest priority channel is simply accepted and DMA transfer on that channel starts.

Whether a DMA request on a given channel is or is not present can be verified by testing the value of the DMA request bit (DREQ) in DMA control register B (DMCNTBn) for that channel.

When a DMA request is accepted and DMA transfer starts, the DMA arbitration status bit (DASTS) for the corresponding channel in the DMA arbitration status register (DMASTS) is set to "1".

11.5 Completion of DMA Transfer and Interrupts

11.5.1 Completion of DMA Transfer

When the value H'0000 0000 is transferred from the working byte count register to the DMA current byte count register (DMCBCTn) (all data has been transferred), the DMA transfer end condition is fulfilled and one DMA transfer is complete.

The operations following detection of the DMA transfer end condition are as follows.

• DMA transfer end condition

The DMA transfer end condition detection bit (DEDET) for the corresponding channel in the DMA transfer end detection register (DMEDET) is set to "1".

• Interrupt request generation

An interrupt request is generated for the interrupt controller according to the settings of the DMA interrupt control register (DMICNT) and the DMA common interrupt control register (DMICNTA).

• Output of DMA end signal

The DMA end signal (DTENDm) is output according the setting of the DMA end signal output control bit (DTCM) in the DMA mode register (DMMODn) for the channel.

• Clearing the DMA transfer enable bit (DEN)

If the DMA transfer enable clear bit (ECLR) in DMA control register B (DMCNTBn) is set to "1", the DEN bit in the DMA control register B (DMCNTBn) is cleared to "0", suspending any subsequent DMA transfer for the channel.

If the DMA transfer enable clear bit (ECLR) is clear ("0"), the DEN bit is not cleared.

• Reloading the source address register

If the DMA source address reload function enable bit (SRLOD) in the DMA control register A (DMCNTAn) is set to "1", the DMA current source address register (DMCSADRn) is reloaded with the value in the DMA reload source address register (DMRSADRn).

• Reloading the destination address register

If the DMA destination address reload function enable bit (DRLOD) in DMA control register A (DMCNTAn) is set to "1", the DMA current destination address register (DMCDADRn) is reloaded with the value in the DMA reload destination address register (DMRDADRn).

• Reloading the byte count register

If the DMA byte count reload function enable bit (BRLOD) in the DMA control register A (DMCNTAn) is set to "1", the DMA current byte count register (DMCBCTn) is reloaded with the value in the DMA reload byte count register (DMRBCTn).

Note: If reloading is not to be executed, set $ECLR = "1"$ to ensure that the DEN bit is cleared.

11.5.2 DMA Interrupt Requests

The DMAC generates two types of interrupt request signal for the interrupt controller. One consists of the interrupt request signals for the individual channels (DMINT_N) and the other is the common interrupt request signal in which the interrupt request signals from all channels are pooled to produce a common interrupt request signal (DMINTA_N).

Figure 11.4 is a block diagram showing how the per-channel and common interrupt requests are generated.

When a DMA transfer ends and the DMA interrupt control bit (DINTM) for the corresponding channel in the DMA interrupt control register (DMICNT) is set to "1", interrupt requests for the corresponding channel are generated.

Only those channels for which the DMA common interrupt request signal control bit (DINTA) in the DMA common interrupt control register (DMICNTA) is set to "1" contribute to the output of common interrupt request.

Once generated, an interrupt request is cleared to "0" by writing a "1" to the corresponding DMA transfer end condition detection bit (DEDET).

11.5.3 DMA End Signal Output

The form in which the DMA end signal (DTENDm) is output differs with the setting of the DMA end signal output control bit (DTCM) in the DMA mode register (DMMODn) for the corresponding channel.

- When DTCM is set to "00", output of the DTEND signal is not valid so the signal remains fixed at the "H" level when and after the DMA transfer ends.
- When DTCM is set to "01", the DTEND signal becomes active (low) one cycle after the start of the read cycle immediately before the end of DMA transfer (the read cycle for the last data transfer).
- When DTCM is set to "10", the DTEND signal becomes active for one cycle after the write cycle immediately before the end of DMA transfer (the write cycle for the last data transfer).
- When DTCM is set to "11", the DTEND signal becomes active for one clock cycle at the same time as the DMA transfer end interrupt is generated.

Output of the DTEND signal is not valid in the case of DMA requests from external peripheral circuits, so the signal remains fixed to "H" regardless of the setting of this bit.

Charts of the timing of DMA end signal output are given in figure 11.5.

Note: The BSC is provided with a write buffer. Writing data to this buffer while writing to the external devices stops bus access in the chip. Because of this, in DMA transfer to or from external devices, the DTEND signal become disabled ("H") before the end of external bus access. In this case the DTEND signal is not synchronized with the external bus access.

11.6 Suspending, Restarting, and Stopping of DMA Transfer

11.6.1 Suspending and Restarting DMA Transfer

Transfer on all channels of the DMAC can be suspended by clearing the DMST bit in the DMA activation control register (DMSCNT) to "0". Transfer on a specific channel can also be suspended by clearing the DMA transfer enable bit (DEN) in DMA control register B (DMCNTBn) for that channel.

If the DMST bit or the corresponding DEN bit is cleared to "0" while single operand transfer or sequential operand transfer is in progress, transfer is suspended on completion of the current single operand transfer regardless of the transfer mode (whether transfer is in cycle-stealing or pipelined mode).

When transfer in the non-stop transfer condition is in progress, DMA transfer is not suspended and continues to completion (until the byte counter reaches "0") even if the DMST bit or corresponding DEN bit is cleared to "0".

To restart DMA transfer on a channel for which transfer has been suspended, set (to "1") whichever of DMST and the corresponding DEN bit has been cleared.

11.6.2 Stopping DMA Transfer on Any Channel

To stop transfer on any channel, suspend transfer on that channel and then initialize the interior state of the DMAC for that channel by setting the DMAC internal state clear bit (DSCLR) in the corresponding DMA control register B (DMCNTBn). In this case, only the transfer state of the DMAC internal circuits is initialized; the registers retain their values.

11.7 DMA Requests

11.7.1 Sources of DMA Requests

The 38 sources of DMA requests include the software trigger and various DMA request signal inputs.

The DMA request source for each channel is specified by the DMA request source select bits (DTCG) in the corresponding DMA control register A (DMCNTAn).

11.7.2 Synchronous Circuits for DMA Request Signals

For each channel of the DMAC, a synchronous circuit is incorporated to manage DMA requests, which are asynchronously input. As a result, a blank period of a few clock cycles appears between activation of the DMA request and actual reflection of the request in the DMA request bits (DREQ) of DMA control register B (DMCNTBn). Figure 11.6 shows an example of timing between the input of a DMA request and the DMA request bit.

Figure 11.6 Example of Timing between DMA Request Input and DMA Request Bit

11.7.3 Sense Mode for DMA Requests

When pins DREQ0 to DREQ3 (DCTG = $"000001"$ to $"000100"$) are specified by the DMA request source selection bits (DTCG), either level sense or edge sense might be required. Make the appropriate setting ("01" or "11" for level sense and "00" or "10" for edge sense) in the input sense selection bits (STRG) of DMA control register A (DMCNTAn).

When the software trigger (DCTG = "000000") is selected as a DMA request source, set these bits to "00" to select the rising-edge sense. When IIC3, SCIF, SSI, RCAN-ET, MTU2, or ADC (DCTG = "000101" to "100101") is selected, set the bits to "10" to select the falling-edge sense. Table 11.4 shows the relationships between the DMA request sources and input sense mode.

Below are further details on level- and edge-sense operation.

(1) Level Sense

When a level sense is specified ($STRG = "01"$ or $T1"$), one level of the DMA request signal indicates the DMA request. Since DMA requests detected in this way are not retained in the DMAC, maintain the requesting level until acceptance of the DMA request has been confirmed.

Figure 11.7 is an example of DMA request reception processing when a level sense has been selected.

Figure 11.7 Example of DMA Request Reception Processing for a Level Sense

When a level sense has been selected, DMA request bit for the channel is masked over the period from the start of the last write access of single operand transfer until four clock pulses (system

clock) after the end of the single operand transfer. This provides a margin in which continued requests for DMA transfer on the same channel are rejected.

Figure 11.8 shows the period over which DMA request bit is masked when a level sense has been selected.

Therefore, for a channel on which level sense has been selected, even when the DMA request signal level is maintained (requesting further DMA transfer) well after the DMA request has been accepted and handled, DMA requests on other channels, if they exist, are accepted. This is because the DMA request on the channel on which level sense has been selected is not considered to exist during the DMA request bit masking period.

In the case of sequential operand transfer, masking is only applied from the end of operand transfer, i.e. when the byte count is 0. The DMA request is not masked while the byte count is non-zero, so channel arbitration is executed without masking of the DMA request during the actual unit transfer operation.

In the case of non-stop transfer, masking is only enabled from the end of the transfer operation, i.e. when the byte count is 0.

If the DMA transfer is not done sequentially, the DMA request must be canceled within three cycles after the end of single operand transfer.

(2) Edge sense

When an edge sense is specified (STRG = $"00"$ or $"10"$), the rising or falling edge of the DMA request signal indicates a DMA request.

When the selected edge is detected, the DMA request bit (DREQ) in the DMA control register B (DMCNTBn) is set to "1". After that, the value in the DMA request bit (DREQ) is retained regardless of shifts in the level of the DMA request signal. After the DMA request has been accepted and the DAM acknowledge signal output, the DMA request bit (DREQ) is automatically cleared to "0".

Since DMA requests are internally retained for a channel in edge sense mode, further occurrences of the selected edge of the DMA request signal are ignored since the DMA request bit (DREQ) has already been set back to "1".

Figure 11.9 is an example of DMA request reception processing when an edge sense is selected.

Figure 11.9 Example of DMA Request Reception Processing when an Edge Sense is Selected

11.8 Determining DMA Channel Priority

11.8.1 Channel Priority Order

Channel priority is allocated in descending order from channel 0; that is priority follows the below relation, where P indicates priority.

 $P_{\text{channel}} > P_{\text{channel}} > P_{\text{channel}} \dots P_{\text{channel}} > P_{\text{channel}}$. This order is fixed.

11.8.2 Operation during Multiple DMA Requests

The DMAC determines the priority every time single operand transfer is performed.

When a DMA request with a higher priority is generated during transfer for one channel, the transfer for the higher-priority channel only starts after the end of the current operand transfer. Figure 11.10 shows overall operation when multiple DMA requests are generated. The thick lines in the figure indicate the periods over which the DMA request signals are at the low level. Here channels 0, 2 and 3 are set to a level sense and channel 1 is set to an edge sense.

- 1. Since the channel 2 request is masked, it is regarded as non-existent. Thus, transfer on channel 3 starts up.
- 2. Since channel 0 has the highest priority, transfer on this channel starts up.
- 3. Since channel 2 has the higher priority of the requests at this point, transfer on this channel restarts.
- 4. Transfer on channel 3 is restarted as there are no other requests at this point.
- 5. When the DMA requests are simultaneously generated for channels 0, 1, and 3, transfer on channel 0 starts up because it has the highest priority.
- 6. After the transfer on channel 0 is complete, transfer on channel 1 starts up because it has the second highest priority.
- 7. A further DMA request (the selected edge) is received on channel 1 while DMA transfer is in progress. Transfer on channel 1 is thus restarted after completion of the current round of transfer on channel 1. No masking period applies in the case of edge sensing.
- 8. On completion of the transfer on channel 1, transfer on channel 3 starts up since there are no other requests.
- 9. No transfer starts up immediately after the end of the unit transfer operation on channel, since channel 3 requests are masked and there are no other requests. Transfer on channel 3 only restarts after the end of the masking period.

Figure 11.10 Overall Operation during Multiple DMA Requests

11.8.3 Output of the DMA Acknowledge and DNA Active Signals

The settings of the DMA active signal output control bits for the source and destination (SACT or DACT) in the corresponding DMA mode register control the output of the DMA active signal (DACT) for a channel.

When SACT is set to 1, the DACT signal is activated in response to read access.

When DACT is set to 1, the DACT signal is activated in response to write access.

When both SACT and DACT are set to 1, the DACT signal is activated in response to read and write access.

However, DACT signals are not activated for DMA requests from external peripheral circuits, regardless of the setting of this bit.

The DMA acknowledge signal (DACK) is output throughout each single operand transfer.

Figure 11.11 is the timing chart for DMA acknowledge and DMA active signal output.

Note: The BSC is provided with a write buffer. Writing data to this buffer while writing to the external devices stops bus access in the chip. Because of this, in DMA transfer to or from external devices, the DACT or DACK signal become disabled ("H") before the end of external bus access. In this case, these signals are not synchronized with the external bus access.

Figure 11.11 Timing of DMA Acknowledge and DNA Active Signal Output

11.9 Units of Transfer and Positioning of Bytes for Transfer

The number of bits (transfer data size) for a single data transfer can be selected from among the byte (8 bits), word (16 bits), and the longword (32 bits).

Figure 11.12 is an example of DMA data-byte control for a 32-bit wide bus.

This transfer data size cannot exceed either of the data bus bit widths supported by the source and destination for DMA transfer. The data bus widths are fixed by the hardware.

Figure 11.12 Example of DMA Data-Byte Control for 32-bit Bus Width
11.10 Reload Function

Reloading can be set up for each transfer parameter (source address, destination address, or byte count) of a channel through the setting of the individual reload function enable bits in the corresponding DMA control register A (DMCNTAn). When the DMA transfer end condition is detected, DMA transfer parameters specified for reloading are automatically reloaded.

(1) Reload and Current Registers

If reloading is not in use, only place the data in the current register. When reloading is in use, place data in both the reload and current registers.

Do not write to the current register during single operand transfer. If data is written to the register during continuous operation, further operation is not guaranteed. Although the reload register can be set during single operand transfer, ensure that this is not the last single operand transfer of a DMA transfer. If the setting is executed after that point, the new setting may not be reloaded on completion of the DMA transfer.

(2) Continuous Transfer to Dispersed areas

The reload function enables continuous transfer to dispersed areas.

Writing to the DMA reload source/destination address register (DMRSADRn/ DMRDADRn) or the DMA reload byte count register (DMRBCTn) before the completion of transfer provides a way of preparing the parameters for the next transfer without affecting the current DMA transfer (current registers). This enables the use of a single channel for the continuous transfer of multiple transfer blocks consisting of different numbers of bytes to and from different transfer areas over a single channel.

Figure 11.13 shows an example of the transfer of blocks between dispersed areas with the aid of the reload function.

Figure 11.13 Example of Transferring Blocks between Dispersed Areas by Using the Reload Function.

11.11 Rotate Function

When rotation is selected as the address "indexing" mode, the address is incremented. On completion of single operand transfer, the value in a working source or working destination address register for which rotation has been selected returns to the value of the source or destination address register (DMCSADRn or DMCDADRn) for the corresponding channel.

Figure 11.14 is an example of transfer using the rotate function (source: rotation, destination: incrementation).

11.12 Transfer Speed

Transfer speeds are calculated as shown below.

(1) Conditions for Calculation

- DMA transfer mode: cycle-stealing transfer mode/pipelined transfer mode
- Transfer unit (one data size): properly aligned 32-bit data
- Operating clock: 60 MHz
- Number of cycles for access to external devices: four cycles for reading; and two cycles for writing.

(2) Formulae Used in Calculation

• Cycle-stealing transfer mode

(data size in unit data transfer) / (number of read cycles $+$ number of write cycles $+$ one idle cycle) \times operating clock

• Pipelined transfer mode

(data size in unit data transfer) / (whichever is larger of number of read or write cycles) \times operating clock

Note: During transfer in the pipelined transfer mode, most read and write cycles overlap.

An example of the calculation of transfer speed is given below.

(a) Transfer between On-chip RAM

Maximum speed of transfer between on-chip RAM (0 wait) and on-chip RAM (0 wait).

• Cycle-stealing transfer mode

4 bytes / (1 read cycle + 1 write cycle + 1 idle cycle) \times 60 MHz = 79.8 Mbytes/sec

• Pipelined transfer mode

Pipelined transfer through a single BIU is not possible. See section 11.4.1 (2), Pipelined Transfer Mode.

(b) Transfer to External Devices

Maximum transfer speed from an on-chip CPU block as the source (0 wait) to an external device (2 write cycles).

• Cycle-stealing transfer mode

```
4 bytes / (1 read cycle + 2 write cycles + 1 idle cycle) \times 60 MHz = 60 Mbytes/sec
```
• Pipelined transfer mode

```
4 bytes / (2 write cycles)\times 60 MHz = 120 Mbytes/sec
```
Maximum transfer speed from an external device (4 read cycles) to an on-chip CPU block source (0 wait)

• Cycle-stealing transfer mode

4 bytes / (4 read cycles $+1$ write cycle $+1$ idle cycle) \times 60 MHz = 39.6 Mbytes/sec

• Pipelined transfer mode

```
4 bytes / (4 read cycles)\times 60 MHz = 60 Mbytes/sec
```
Maximum transfer speed from an external device (4 read cycles) to an external device (2 write cycles)

• Cycle-stealing transfer mode

4 bytes / (4 read cycles $+ 2$ write cycles $+ 1$ idle cycle) \times 60 MHz = 34.2 Mbytes/sec

• Pipelined transfer mode

No pipelined transfer is possible between the external devices.

Note: Access to external devices is controlled by the settings of the BSC control registers. For details, see section 9, Bus State Controller (BSC).

11.13 Usage Note

11.13.1 Note on Making a Transition To Software Standby Mode or Deep Standby Mode

If the SLEEP instruction is executed to make a transition to software standby mode or deep standby mode during transfer by the DMAC, the DMAC stops its operation without waiting for the completion of the transfer. Thus, the DMA transfer is not guaranteed. Therefore, when making a transition to software standby mode or deep standby mode, wait for the completion of the DMA transfer or stop the DMA transfer to execute the SLEEP instruction.

Section 12 Multi-Function Timer Pulse Unit 2 (MTU2)

This LSI has an on-chip multi-function timer pulse unit 2 (MTU2) that comprises six 16-bit timer channels.

12.1 Features

- Up to 16 pulse input/output lines and three pulse input lines
- Selection of eight counter input clocks for each channel (four clocks for channel 5)
- The following operations can be set for channels 0 to 4:
	- Waveform output at compare match
	- Input capture function
	- Counter clear operation
	- Multiple timer counters (TCNT) can be written to simultaneously
	- Simultaneous clearing by compare match and input capture is possible
	- Register simultaneous input/output is possible by synchronous counter operation
	- \overline{A} A maximum 12-phase PWM output is possible in combination with synchronous operation
- Buffer operation settable for channels 0, 3, and 4
- Phase counting mode settable independently for each of channels 1 and 2
- Cascade connection operation
- Fast access via internal 16-bit bus
- 28 interrupt sources
- Automatic transfer of register data
- A/D converter start trigger can be generated
- Module standby mode can be settable
- A total of six-phase waveform output, which includes complementary PWM output, and positive and negative phases of reset PWM output by interlocking operation of channels 3 and 4, is possible.
- AC synchronous motor (brushless DC motor) drive mode using complementary PWM output and reset PWM output is settable by interlocking operation of channels 0, 3, and 4, and the selection of two types of waveform outputs (chopping and level) is possible.
- Dead time compensation counter available in channel 5
- In complementary PWM mode, interrupts at the crest and trough of the counter value and A/D converter start triggers can be skipped.

Table 12.1 MTU2 Functions

[Legend]

√: Possible

—: Not possible

Figure 12.1 Block Diagram of MTU2

12.2 Input/Output Pins

Table 12.2 Pin Configuration

Channel Pin Name I/O Function

Note: For the pin configuration in complementary PWM mode, see table 12.54.

12.3 Register Descriptions

The MTU2 has the following registers. For details on register addresses and register states during each process, refer to section 30, List of Registers. To distinguish registers in each channel, an underscore and the channel number are added as a suffix to the register name; TCR for channel 0 is expressed as TCR_0.

Table 12.3 Register Configuration

12.3.1 Timer Control Register (TCR)

The TCR registers are 8-bit readable/writable registers that control the TCNT operation for each channel. The MTU2 has a total of eight TCR registers, one each for channels 0 to 4 and three (TCRU_5, TCRV_5, and TCRW_5) for channel 5. TCR register settings should be conducted only when TCNT operation is stopped.

[Legend]

x: Don't care

Table 12.4 CCLR0 to CCLR2 (Channels 0, 3, and 4)

Notes: 1. Synchronous operation is set by setting the SYNC bit in TSYR to 1.

 2. When TGRC or TGRD is used as a buffer register, TCNT is not cleared because the buffer register setting has priority, and compare match/input capture does not occur.

Notes: 1. Synchronous operation is selected by setting the SYNC bit in TSYR to 1.

2. Bit 7 is reserved in channels 1 and 2. It is always read as 0 and cannot be modified.

Table 12.6 TPSC0 to TPSC2 (Channel 0)

Table 12.7 TPSC0 to TPSC2 (Channel 1)

Note: This setting is ignored when channel 1 is in phase counting mode.

Table 12.8 TPSC0 to TPSC2 (Channel 2)

Note: This setting is ignored when channel 2 is in phase counting mode.

Table 12.10 TPSC1 and TPSC0 (Channel 5)

Note: Bits 7 to 2 are reserved in channel 5. These bits are always read as 0. The write value should always be 0.

12.3.2 Timer Mode Register (TMDR)

The TMDR registers are 8-bit readable/writable registers that are used to set the operating mode of each channel. The MTU2 has five TMDR registers, one each for channels 0 to 4. TMDR register settings should be changed only when TCNT operation is stopped.

Table 12.11 Setting of Operation Mode by Bits MD0 to MD3

[Legend]

X: Don't care

Notes: 1. PWM mode 2 cannot be set for channels 3 and 4.

- 2. Phase counting mode cannot be set for channels 0, 3, and 4.
- 3. Reset synchronous PWM mode, complementary PWM mode can only be set for channel 3. When channel 3 is set to reset synchronous PWM mode or complementary PWM mode, the channel 4 settings become ineffective and automatically conform to the channel 3 settings. However, do not set channel 4 to reset synchronous PWM mode or complementary PWM mode. Reset synchronous PWM mode and complementary PWM mode cannot be set for channels 0, 1, and 2.

12.3.3 Timer I/O Control Register (TIOR)

The TIOR registers are 8-bit readable/writable registers that control the TGR registers. The MTU2 has a total of eleven TIOR registers, two each for channels 0, 3, and 4, one each for channels 1 and 2, and three (TIORU_5, TIORV_5, and TIORW_5) for channel 5.

TIOR should be set while TMDR is set in normal operation, PWM mode, or phase counting mode.

The initial output specified by TIOR is valid when the counter is stopped (the CST bit in TSTR is cleared to 0). Note also that, in PWM mode 2, the output at the point at which the counter is cleared to 0 is specified.

When TGRC or TGRD is designated for buffer operation, this setting is invalid and the register operates as a buffer register.

• TIORH_0, TIOR_1, TIOR_2, TIORH_3, TIORH_4

• TIORL_0, TIORL_3, TIORL_4

• TIORU_5, TIORV_5, TIORW_5

 Description

Table 12.12 TIORH_0 (Channel 0)

[Legend]

X: Don't care

Table 12.13 TIORL_0 (Channel 0)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFB bit in TMDR_0 is set to 1 and TGRD_0 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

 Description

Table 12.14 TIOR_1 (Channel 1)

X: Don't care

Table 12.15 TIOR_2 (Channel 2)

[Legend]

X: Don't care

 Description

Table 12.16 TIORH_3 (Channel 3)

[Legend]

X: Don't care

Table 12.17 TIORL_3 (Channel 3)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFB bit in TMDR_3 is set to 1 and TGRD_3 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

 Description

Table 12.18 TIORH_4 (Channel 4)

[Legend]

X: Don't care

Table 12.19 TIORL_4 (Channel 4)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

2. When the BFB bit in TMDR_4 is set to 1 and TGRD_4 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.
Description

Table 12.20 TIORH_0 (Channel 0)

[Legend]

X: Don't care

Table 12.21 TIORL_0 (Channel 0)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFA bit in TMDR_0 is set to 1 and TGRC_0 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

 Description

Table 12.22 TIOR_1 (Channel 1)

X: Don't care

Table 12.23 TIOR_2 (Channel 2)

[Legend]

X: Don't care

 Description

Table 12.24 TIORH_3 (Channel 3)

[Legend]

X: Don't care

Table 12.25 TIORL_3 (Channel 3)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

2. When the BFA bit in TMDR_3 is set to 1 and TGRC_3 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

 Description

Table 12.26 TIORH_4 (Channel 4)

[Legend]

X: Don't care

Table 12.27 TIORL_4 (Channel 4)

[Legend]

X: Don't care

Notes: 1. After power-on reset, 0 is output until TIOR is set.

 2. When the BFA bit in TMDR_4 is set to 1 and TGRC_4 is used as a buffer register, this setting is invalid and input capture/output compare is not generated.

Table 12.28 TIORU_5, TIORV_5, and TIORW_5 (Channel 5)

[Legend]

X: Don't care

12.3.4 Timer Compare Match Clear Register (TCNTCMPCLR)

TCNTCMPCLR is an 8-bit readable/writable register that specifies requests to clear TCNTU_5, TCNTV_5, and TCNTW_5. The MTU2 has one TCNTCMPCLR in channel 5.

12.3.5 Timer Interrupt Enable Register (TIER)

The TIER registers are 8-bit readable/writable registers that control enabling or disabling of interrupt requests for each channel. The MTU2 has seven TIER registers, two for channel 0 and one each for channels 1 to 5.

[•] TIER_0, TIER_1, TIER_2, TIER_3, TIER_4

Bit:		\mathcal{L}		4 3		
						TTGE TTGE2 TCIEU TCIEV TGIED TGIEC TGIEB TGIEA
Initial value: 0	$\sqrt{1}$	\sim 0	Ω			
	R/W: R/W R/W R/W R/W R/W R/W R/W R/W					

• TIER2_0

• TIER_5

12.3.6 Timer Status Register (TSR)

The TSR registers are 8-bit readable/writable registers that indicate the status of each channel. The MTU2 has seven TSR registers, two for channel 0 and one each for channels 1 to 5.

• TSR_0, TSR_1, TSR_2, TSR_3, TSR_4

Bit [.]					
	TCFD				$ TCFU TCFV TGFD TGFC TGFB TGFA $
Initial value:					
B/W		R.			R/(W)*1R/(W)*1R/(W)*1R/(W)*1R/(W)*1R/(W)*1R/(W)*1

Note: 1. Writing 0 to this bit after reading it as 1 clears the flag and is the only allowed way.

 2. When writing to the timer status register (TSR), write 0 to the bit to be cleared after reading 1. Write 1 to other bits. But 1 is not actually written and the previous value is held.

• TSR2_0

Note: 1. Writing 0 to this bit after reading it as 1 clears the flag and is the only allowed way.

 2. When writing to the timer status register (TSR), write 0 to the bit to be cleared after reading 1. Write 1 to other bits. But 1 is not actually written and the previous value is held.

• TSR_5

Bit ⁻							
					I CMF U5 ∣	$\begin{array}{ c c c }\n\hline\n\text{CMF} & \text{CMF} \\ \hline\n\text{V5} & \text{W5}\n\end{array}$	
Initial value:	O						
B/W		R	R	R	$R/(W) * 1R/(W) * 1R/(W) * 1$		

Note: 1. Writing 0 to this bit after reading it as 1 clears the flag and is the only allowed way.

 2. Timing to transfer is set by the IOC bit in the timer I/O control register U_5/V_5/W_5 (TIORU_5/V_5/W_5).

12.3.7 Timer Buffer Operation Transfer Mode Register (TBTM)

TBTM is an 8-bit readable/writable register that specifies the timing for transferring data from the buffer register to the timer general register in PWM mode. The MTU2 has three TBTM registers, one each for channels 0, 3, and 4.

12.3.8 Timer Input Capture Control Register (TICCR)

TICCR is an 8-bit readable/writable register that specifies input capture conditions when TCNT_1 and TCNT_2 are cascaded. The MTU2 has one TICCR in channel 1.

12.3.9 Timer A/D Converter Start Request Control Register (TADCR)

TADCR is a 16-bit readable/writable register that enables or disables A/D converter start requests and specifies whether to link A/D converter start requests with interrupt skipping operation. The MTU2 has one TADCR in channel 4.

Note: * Do not set to 1 when complementary PWM mode is not selected.

Notes: 1. TADCR must not be accessed in eight bits; it should always be accessed in 16 bits.

- 2. When interrupt skipping is disabled (the T3AEN and T4VEN bits in the timer interrupt skipping set register (TITCR) are cleared to 0 or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0), do not link A/D converter start requests with interrupt skipping operation (clear the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in the timer A/D converter start request control register (TADCR) to 0).
- 3. If link with interrupt skipping is enabled while interrupt skipping is disabled, A/D converter start requests will not be issued.
- * Do not set to 1 when complementary PWM mode is not selected.

Table 12.29 Setting of Transfer Timing by BF1 and BF0 Bits

match occurs between TCNT_3 and TGRA_3 in reset-synchronized PWM mode, or when compare match occurs between TCNT 4 and TGRA 4 in PWM mode 1 or normal operation mode.

2. These settings are prohibited when complementary PWM mode is not selected.

12.3.10 Timer A/D Converter Start Request Cycle Set Registers (TADCORA_4 and TADCORB_4)

TADCORA_4 and TADCORB_4 are 16-bit readable/writable registers. When the TCNT_4 count reaches the value in TADCORA_4 or TADCORB_4, a corresponding A/D converter start request will be issued.

TADCORA_4 and TADCORB_4 are initialized to H'FFFF.

Note: TADCORA 4 and TADCORB 4 must not be accessed in eight bits; they should always be accessed in 16 bits.

12.3.11 Timer A/D Converter Start Request Cycle Set Buffer Registers (TADCOBRA_4 and TADCOBRB_4)

TADCOBRA_4 and TADCOBRB_4 are 16-bit readable/writable registers. When the crest or trough of the TCNT_4 count is reached, these register values are transferred to TADCORA_4 and TADCORB_4, respectively.

TADCOBRA_4 and TADCOBRB_4 are initialized to H'FFFF.

Note: TADCOBRA 4 and TADCOBRB 4 must not be accessed in eight bits; they should always be accessed in 16 bits.

12.3.12 Timer Counter (TCNT)

The TCNT counters are 16-bit readable/writable counters. The MTU2 has eight TCNT counters, one each for channels 0 to 4 and three (TCNTU_5, TCNTV_5, and TCNTW_5) for channel 5.

The TCNT counters are initialized to H'0000 by a reset.

Note: The TCNT counters must not be accessed in eight bits; they should always be accessed in 16 bits.

12.3.13 Timer General Register (TGR)

The TGR registers are 16-bit readable/writable registers. The MTU2 has 21 TGR registers, six for channel 0, two each for channels 1 and 2, four each for channels 3 and 4, and three for channel 5.

TGRA, TGRB, TGRC, and TGRD function as either output compare or input capture registers. TGRC and TGRD for channels 0, 3, and 4 can also be designated for operation as buffer registers. TGR buffer register combinations are TGRA and TGRC, and TGRB and TGRD.

TGRE_0 and TGRF_0 function as compare registers. When the TCNT_0 count matches the TGRE_0 value, an A/D converter start request can be issued. TGRF can also be designated for operation as a buffer register. TGR buffer register combination is TGRE and TGRF.

TGRU_5, TGRV_5, and TGRW_5 function as compare match, input capture, or external pulse width measurement registers.

Note: The TGR registers must not be accessed in eight bits; they should always be accessed in 16 bits. TGR registers are initialized to H'FFFF.

12.3.14 Timer Start Register (TSTR)

TSTR is an 8-bit readable/writable register that selects operation/stoppage of TCNT for channels 0 to 4.

TSTR_5 is an 8-bit readable/writable register that selects operation/stoppage of TCNTU_5, TCNTV_5, and TCNTW_5 for channel 5.

When setting the operating mode in TMDR or setting the count clock in TCR, first stop the TCNT counter.

• TSTR

• TSTR_5

12.3.15 Timer Synchronous Register (TSYR)

TSYR is an 8-bit readable/writable register that selects independent operation or synchronous operation for the channel 0 to 4 TCNT counters. A channel performs synchronous operation when the corresponding bit in TSYR is set to 1.

12.3.16 Timer Counter Synchronous Start Register (TCSYSTR)

TCSYSTR is an 8-bit readable/writable register that specifies synchronous start of the MTU2 counters.

Bit:							
					SCHO SCH1 SCH2 SCH3 SCH4		
Initial value:							
$R/W: R/(W)*R/(W)*R/(W)*R/(W)*R/(W)*$							

Note: $*$ Only 1 can be written to set the register.

Note: * Only 1 can be written to set the register.
12.3.17 Timer Read/Write Enable Register (TRWER)

TRWER is an 8-bit readable/writable register that enables or disables access to the registers and counters which have write-protection capability against accidental modification in channels 3 and 4.

• Registers and counters having write-protection capability against accidental modification 22 registers: TCR_3, TCR_4, TMDR_3, TMDR_4, TIORH_3, TIORH_4, TIORL_3, TIORL_4, TIER_3, TIER_4, TGRA_3, TGRA_4, TGRB_3, TGRB_4, TOER, TOCR1, TOCR2, TGCR, TCDR, TDDR, TCNT_3, and TCNT_4.

12.3.18 Timer Output Master Enable Register (TOER)

TOER is an 8-bit readable/writable register that enables/disables output settings for output pins TIOC4D, TIOC4C, TIOC3D, TIOC4B, TIOC4A, and TIOC3B. These pins do not output correctly if the TOER bits have not been set. Set TOER of channel 3 and channel 4 prior to setting TIOR of channel 3 and channel 4.

Note: * The inactive level is determined by the settings in timer output control registers 1 and 2 (TOCR1 and TOCR2). For details, refer to section 12.3.19, Timer Output Control Register 1 (TOCR1), and section 12.3.20, Timer Output Control Register 2 (TOCR2). Set these bits to 1 to enable MTU2 output in other than complementary PWM or resetsynchronized PWM mode. If these bits are set to 0, low level is output.

12.3.19 Timer Output Control Register 1 (TOCR1)

TOCR1 is an 8-bit readable/writable register that enables/disables PWM synchronized toggle output in complementary PWM mode/reset synchronized PWM mode, and controls output level inversion of PWM output.

Note: 1. This bit can be set to 1 only once after a power-on reset. After 1 is written, 0 cannot be written to the bit.

written to the bit.

- 2. Setting the TOCL bit to 1 prevents accidental modification when the CPU goes out of control.
- 3. Clearing the TOCS bit to 0 makes this bit setting valid.

Table 12.30 Output Level Select Function

Note: The reverse phase waveform initial output value changes to active level after elapse of the dead time after count start.

Table 12.31 Output Level Select Function

Figure 12.2 shows an example of complementary PWM mode output (1 phase) when OLSN = 1 and $OLSP = 1$.

Figure 12.2 Complementary PWM Mode Output Level Example

12.3.20 Timer Output Control Register 2 (TOCR2)

TOCR2 is an 8-bit readable/writable register that controls output level inversion of PWM output in complementary PWM mode and reset-synchronized PWM mode.

Note: * Setting the TOCS bit in TOCR1 to 1 makes this bit setting valid.

Table 12.32 Setting of Bits BF1 and BF0

Table 12.33 TIOC4D Output Level Select Function

Note: The reverse phase waveform initial output value changes to the active level after elapse of the dead time after count start.

Table 12.34 TIOC4B Output Level Select Function

Table 12.35 TIOC4C Output Level Select Function

Note: The reverse phase waveform initial output value changes to the active level after elapse of the dead time after count start.

Table 12.36 TIOC4A Output Level Select Function

Table 12.37 TIOC3D Output Level Select Function

Note: The reverse phase waveform initial output value changes to the active level after elapse of the dead time after count start.

Table 12.38 TIOC3B Output Level Select Function

12.3.21 Timer Output Level Buffer Register (TOLBR)

TOLBR is an 8-bit readable/writable register that functions as a buffer for TOCR2 and specifies the PWM output level in complementary PWM mode and reset-synchronized PWM mode.

Figure 12.3 shows an example of the PWM output level setting procedure in buffer operation.

12.3.22 Timer Gate Control Register (TGCR)

TGCR is an 8-bit readable/writable register that controls the waveform output necessary for brushless DC motor control in reset-synchronized PWM mode/complementary PWM mode. These register settings are ineffective for anything other than complementary PWM mode/resetsynchronized PWM mode.

Table 12.39 Output level Select Function

12.3.23 Timer Subcounter (TCNTS)

TCNTS is a 16-bit read-only counter that is used only in complementary PWM mode.

The initial value of TCNTS is H'0000.

Note: Accessing the TCNTS in 8-bit units is prohibited. Always access in 16-bit units.

12.3.24 Timer Dead Time Data Register (TDDR)

TDDR is a 16-bit register, used only in complementary PWM mode that specifies the TCNT_3 and TCNT_4 counter offset values. In complementary PWM mode, when the TCNT_3 and TCNT_4 counters are cleared and then restarted, the TDDR register value is loaded into the TCNT_3 counter and the count operation starts.

The initial value of TDDR is H'FFFF.

Note: Accessing the TDDR in 8-bit units is prohibited. Always access in 16-bit units.

12.3.25 Timer Cycle Data Register (TCDR)

TCDR is a 16-bit register used only in complementary PWM mode. Set half the PWM carrier sync value as the TCDR register value. This register is constantly compared with the TCNTS counter in complementary PWM mode, and when a match occurs, the TCNTS counter switches direction (decrement to increment).

The initial value of TCDR is H'FFFF.

Note: Accessing the TCDR in 8-bit units is prohibited. Always access in 16-bit units.

12.3.26 Timer Cycle Buffer Register (TCBR)

TCBR is a 16-bit register used only in complementary PWM mode. It functions as a buffer register for the TCDR register. The TCBR register values are transferred to the TCDR register with the transfer timing set in the TMDR register.

The initial value of TCBR is H'FFFF.

12.3.27 Timer Interrupt Skipping Set Register (TITCR)

TITCR is an 8-bit readable/writable register that enables or disables interrupt skipping and specifies the interrupt skipping count. The MTU2 has one TITCR.

Note: * When 0 is specified for the interrupt skipping count, no interrupt skipping will be performed. Before changing the interrupt skipping count, be sure to clear the T3AEN and T4VEN bits to 0 to clear the skipping counter (TITCNT).

Table 12.40 Setting of Interrupt Skipping Count by Bits 3ACOR2 to 3ACOR0

Table 12.41 Setting of Interrupt Skipping Count by Bits 4VCOR2 to 4VCOR0

12.3.28 Timer Interrupt Skipping Counter (TITCNT)

TITCNT is an 8-bit readable/writable counter. The MTU2 has one TITCNT. TITCNT retains the value even after TCNT_3 or TCNT_4 stops counting.

Note: Clear the T3AEN and T4VEN bits in TITCR to 0, to clear the value of TITCNT.

12.3.29 Timer Buffer Transfer Set Register (TBTER)

TBTER is an 8-bit readable/writable register that enables or disables transfer from the buffer registers* used in complementary PWM mode to the temporary registers and specifies whether to link the transfer with interrupt skipping operation. The MTU2 has one TBTER.

R01UH0025EJ0300 Rev. 3.00
 RENESAS Sep 24, 2010

Table 12.42 Setting of Bits BTE1 and BTE0

 2. When interrupt skipping is disabled (the T3AEN and T4VEN bits are cleared to 0 in the timer interrupt skipping set register (TITCR) or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0)), be sure to disable link of buffer transfer with interrupt skipping (clear the BTE1 bit in the timer buffer transfer set register (TBTER) to 0). If link with interrupt skipping is enabled while interrupt skipping is disabled, buffer transfer will not be performed.

12.3.30 Timer Dead Time Enable Register (TDER)

TDER is an 8-bit readable/writable register that controls dead time generation in complementary PWM mode. The MTU2 has one TDER in channel 3. TDER must be modified only while TCNT stops.

Note: * TDDR must be set to 1 or a larger value.

12.3.31 Timer Waveform Control Register (TWCR)

TWCR is an 8-bit readable/writable register that controls the waveform when synchronous counter clearing occurs in TCNT_3 and TCNT_4 in complementary PWM mode and specifies whether to clear the counters at TGRA_3 compare match. The CCE bit and WRE bit in TWCR must be modified only while TCNT stops.

Bit ⁻		6	b					
	CCE							WRE
Initial value: 0*		O	O	O		\mathbf{U}		
	R/W: R/(W)	R	R	К	R	R	н	R/(W)

Note: $*$ Do not set to 1 when complementary PWM mode is not selected.

Note: * Do not set to 1 when complementary PWM mode is not selected.

12.3.32 Bus Master Interface

The timer counters (TCNT), general registers (TGR), timer subcounter (TCNTS), timer cycle buffer register (TCBR), timer dead time data register (TDDR), timer cycle data register (TCDR), timer A/D converter start request control register (TADCR), timer A/D converter start request cycle set registers (TADCOR), and timer A/D converter start request cycle set buffer registers (TADCOBR) are 16-bit registers. A 16-bit data bus to the bus master enables 16-bit read/writes. 8 bit read/write is not possible. Always access in 16-bit units.

All registers other than the above registers are 8-bit registers. These are connected to the CPU by a 16-bit data bus, so 16-bit read/writes and 8-bit read/writes are both possible.

12.4 Operation

12.4.1 Basic Functions

Each channel has a TCNT and TGR register. TCNT performs up-counting, and is also capable of free-running operation, cycle counting, and external event counting.

Each TGR can be used as an input capture register or output compare register.

Always select MTU2 external pins set function using the pin function controller (PFC).

(1) Counter Operation

When one of bits CST0 to CST4 in TSTR or bits CSTU5, CSTV5, and CSTW5 in TSTR 5 is set to 1, the TCNT counter for the corresponding channel begins counting. TCNT can operate as a free-running counter, periodic counter, for example.

(a) Example of Count Operation Setting Procedure

Figure 12.4 shows an example of the count operation setting procedure.

Figure 12.4 Example of Counter Operation Setting Procedure

(b) Free-Running Count Operation and Periodic Count Operation:

Immediately after a reset, the MTU2's TCNT counters are all designated as free-running counters. When the relevant bit in TSTR is set to 1 the corresponding TCNT counter starts up-count operation as a free-running counter. When TCNT overflows (from H'FFFF to H'0000), the TCFV bit in TSR is set to 1. If the value of the corresponding TCIEV bit in TIER is 1 at this point, the MTU2 requests an interrupt. After overflow, TCNT starts counting up again from H'0000.

Figure 12.5 illustrates free-running counter operation.

Figure 12.5 Free-Running Counter Operation

When compare match is selected as the TCNT clearing source, the TCNT counter for the relevant channel performs periodic count operation. The TGR register for setting the period is designated as an output compare register, and counter clearing by compare match is selected by means of bits CCLR0 to CCLR2 in TCR. After the settings have been made, TCNT starts up-count operation as a periodic counter when the corresponding bit in TSTR is set to 1. When the count value matches the value in TGR, the TGF bit in TSR is set to 1 and TCNT is cleared to H'0000.

If the value of the corresponding TGIE bit in TIER is 1 at this point, the MTU2 requests an interrupt. After a compare match, TCNT starts counting up again from H'0000.

Figure 12.6 illustrates periodic counter operation.

Figure 12.6 Periodic Counter Operation

(2) Waveform Output by Compare Match

The MTU2 can perform 0, 1, or toggle output from the corresponding output pin using compare match.

(a) Example of Setting Procedure for Waveform Output by Compare Match

Figure 12.7 shows an example of the setting procedure for waveform output by compare match

Figure 12.7 Example of Setting Procedure for Waveform Output by Compare Match

(b) Examples of Waveform Output Operation:

Figure 12.8 shows an example of 0 output/1 output.

In this example, TCNT has been designated as a free-running counter, and settings have been made such that 1 is output by compare match A, and 0 is output by compare match B. When the set level and the pin level coincide, the pin level does not change.

Figure 12.8 Example of 0 Output/1 Output Operation

Figure 12.9 shows an example of toggle output.

In this example, TCNT has been designated as a periodic counter (with counter clearing on compare match B), and settings have been made such that the output is toggled by both compare match A and compare match B.

Figure 12.9 Example of Toggle Output Operation

(3) Input Capture Function:

The TCNT value can be transferred to TGR on detection of the TIOC pin input edge.

Rising edge, falling edge, or both edges can be selected as the detected edge. For channels 0 and 1, it is also possible to specify another channel's counter input clock or compare match signal as the input capture source.

Note: When another channel's counter input clock is used as the input capture input for channels 0 and 1, Pφ/1 should not be selected as the counter input clock used for input capture input. Input capture will not be generated if P $\phi/1$ is selected.

(a) Example of Input Capture Operation Setting Procedure

Figure 12.10 shows an example of the input capture operation setting procedure.

Figure 12.10 Example of Input Capture Operation Setting Procedure

(b) Example of Input Capture Operation:

Figure 12.11 shows an example of input capture operation.

In this example both rising and falling edges have been selected as the TIOCA pin input capture input edge, the falling edge has been selected as the TIOCB pin input capture input edge, and counter clearing by TGRB input capture has been designated for TCNT.

Figure 12.11 Example of Input Capture Operation

12.4.2 Synchronous Operation

In synchronous operation, the values in a number of TCNT counters can be rewritten simultaneously (synchronous presetting). Also, a number of TCNT counters can be cleared simultaneously by making the appropriate setting in TCR (synchronous clearing).

Synchronous operation enables TGR to be incremented with respect to a single time base.

Channels 0 to 4 can all be designated for synchronous operation. Channel 5 cannot be used for synchronous operation.

(1) Example of Synchronous Operation Setting Procedure:

Figure 12.12 shows an example of the synchronous operation setting procedure.

[5] Set to 1 the CST bits in TSTR for the relevant channels, to start the count operation.

Figure 12.12 Example of Synchronous Operation Setting Procedure

(2) Example of Synchronous Operation

Figure 12.13 shows an example of synchronous operation.

In this example, synchronous operation and PWM mode 1 have been designated for channels 0 to 2, TGRB_0 compare match has been set as the channel 0 counter clearing source, and synchronous clearing has been set for the channel 1 and 2 counter clearing source.

Three-phase PWM waveforms are output from pins TIOC0A, TIOC1A, and TIOC2A. At this time, synchronous presetting, and synchronous clearing by TGRB_0 compare match, are performed for channel 0 to 2 TCNT counters, and the data set in TGRB_0 is used as the PWM cycle.

TCNT_0 to TCNT_2 values H'0000 TIOC0A TIOC1A TGRB_0 Synchronous clearing by TGRB_0 compare match TGRA_2 TGRA_1 TGRB_2 TGRA_0 TGRB_1 TIOC2A Time

For details of PWM modes, see section 12.4.5, PWM Modes.

Figure 12.13 Example of Synchronous Operation

12.4.3 Buffer Operation

Buffer operation, provided for channels 0, 3, and 4, enables TGRC and TGRD to be used as buffer registers. In channel 0, TGRF can also be used as a buffer register.

Buffer operation differs depending on whether TGR has been designated as an input capture register or as a compare match register.

Note: TGRE 0 cannot be designated as an input capture register and can only operate as a compare match register.

Table 12.43 shows the register combinations used in buffer operation.

Table 12.43 Register Combinations in Buffer Operation

Channel	Timer General Register	Buffer Register
$\mathbf 0$	TGRA 0	TGRC_0
	TGRB_0	TGRD_0
	TGRE 0	TGRF_0
3	TGRA_3	TGRC_3
	TGRB_3	TGRD 3
4	TGRA_4	TGRC_4
	TGRB_4	TGRD_4

When TGR is an output compare register

When a compare match occurs, the value in the buffer register for the corresponding channel is transferred to the timer general register.

This operation is illustrated in figure 12.14.

Figure 12.14 Compare Match Buffer Operation

When TGR is an input capture register

When input capture occurs, the value in TCNT is transferred to TGR and the value previously held in the timer general register is transferred to the buffer register.

This operation is illustrated in figure 12.15.

Figure 12.15 Input Capture Buffer Operation

(1) Example of Buffer Operation Setting Procedure

Figure 12.16 shows an example of the buffer operation setting procedure.

Figure 12.16 Example of Buffer Operation Setting Procedure

(2) Examples of Buffer Operation

(a) When TGR is an output compare register

Figure 12.17 shows an operation example in which PWM mode 1 has been designated for channel 0, and buffer operation has been designated for TGRA and TGRC. The settings used in this example are TCNT clearing by compare match B, 1 output at compare match A, and 0 output at compare match B. In this example, the TTSA bit in TBTM is cleared to 0.

As buffer operation has been set, when compare match A occurs the output changes and the value in buffer register TGRC is simultaneously transferred to timer general register TGRA. This operation is repeated each time that compare match A occurs.

For details of PWM modes, see section 12.4.5, PWM Modes.

Figure 12.17 Example of Buffer Operation (1)

(b) When TGR is an input capture register

Figure 12.18 shows an operation example in which TGRA has been designated as an input capture register, and buffer operation has been designated for TGRA and TGRC.

Counter clearing by TGRA input capture has been set for TCNT, and both rising and falling edges have been selected as the TIOCA pin input capture input edge.

As buffer operation has been set, when the TCNT value is stored in TGRA upon the occurrence of input capture A, the value previously stored in TGRA is simultaneously transferred to TGRC.

Figure 12.18 Example of Buffer Operation (2)

(3) Selecting Timing for Transfer from Buffer Registers to Timer General Registers in Buffer Operation:

The timing for transfer from buffer registers to timer general registers can be selected in PWM mode 1 or 2 for channel 0 or in PWM mode 1 for channels 3 and 4 by setting the buffer operation transfer mode registers (TBTM_0, TBTM_3, and TBTM_4). Either compare match (initial setting) or TCNT clearing can be selected for the transfer timing. TCNT clearing as transfer timing is one of the following cases.

- When TCNT overflows (H'FFFF to H'0000)
- When H'0000 is written to TCNT during counting
- When TCNT is cleared to H'0000 under the condition specified in the CCLR2 to CCLR0 bits in TCR

Note: TBTM must be modified only while TCNT stops.

Figure 12.19 shows an operation example in which PWM mode 1 is designated for channel 0 and buffer operation is designated for TGRA_0 and TGRC_0. The settings used in this example are TCNT_0 clearing by compare match B, 1 output at compare match A, and 0 output at compare match B. The TTSA bit in TBTM 0 is set to 1.

Figure 12.19 Example of Buffer Operation When TCNT_0 Clearing is Selected for TGRC_0 to TGRA_0 Transfer Timing

12.4.4 Cascaded Operation

In cascaded operation, two 16-bit counters for different channels are used together as a 32-bit counter.

This function works by counting the channel 1 counter clock upon overflow/underflow of TCNT_2 as set in bits TPSC0 to TPSC2 in TCR.

Underflow occurs only when the lower 16-bit TCNT is in phase-counting mode.

Table 12.44 shows the register combinations used in cascaded operation.

Note: When phase counting mode is set for channel 1, the counter clock setting is invalid and the counters operates independently in phase counting mode.

Table 12.44 Cascaded Combinations

For simultaneous input capture of TCNT_1 and TCNT_2 during cascaded operation, additional input capture input pins can be specified by the input capture control register (TICCR). For input capture in cascade connection, refer to section 12.7.22, Simultaneous Capture of TCNT_1 and TCNT_2 in Cascade Connection.

Table 12.45 show the TICCR setting and input capture input pins.

Table 12.45 TICCR Setting and Input Capture Input Pins

(1) Example of Cascaded Operation Setting Procedure

Figure 12.20 shows an example of the setting procedure for cascaded operation.

Figure 12.20 Cascaded Operation Setting Procedure

(2) Cascaded Operation Example (a)

Figure 12.21 illustrates the operation when TCNT_2 overflow/underflow counting has been set for TCNT 1 and phase counting mode has been designated for channel 2.

TCNT_1 is incremented by TCNT_2 overflow and decremented by TCNT_2 underflow.

Figure 12.21 Cascaded Operation Example (a)
(3) Cascaded Operation Example (b)

Figure 12.22 illustrates the operation when TCNT 1 and TCNT 2 have been cascaded and the I2AE bit in TICCR has been set to 1 to include the TIOC2A pin in the TGRA_1 input capture conditions. In this example, the IOA0 to IOA3 bits in TIOR_1 have selected the TIOC1A rising edge for the input capture timing while the IOA0 to IOA3 bits in TIOR_2 have selected the TIOC2A rising edge for the input capture timing.

Under these conditions, the rising edge of both TIOC1A and TIOC2A is used for the TGRA 1 input capture condition. For the TGRA_2 input capture condition, the TIOC2A rising edge is used.

Figure 12.22 Cascaded Operation Example (b)

(4) Cascaded Operation Example (c)

Figure 12.23 illustrates the operation when TCNT 1 and TCNT 2 have been cascaded and the I2AE and I1AE bits in TICCR have been set to 1 to include the TIOC2A and TIOC1A pins in the TGRA_1 and TGRA_2 input capture conditions, respectively. In this example, the IOA0 to IOA3 bits in both TIOR_1 and TIOR_2 have selected both the rising and falling edges for the input capture timing. Under these conditions, the ORed result of TIOC1A and TIOC2A input is used for the TGRA_1 and TGRA_2 input capture conditions.

Figure 12.23 Cascaded Operation Example (c)

(5) Cascaded Operation Example (d)

Figure 12.24 illustrates the operation when TCNT_1 and TCNT_2 have been cascaded and the I2AE bit in TICCR has been set to 1 to include the TIOC2A pin in the TGRA_1 input capture conditions. In this example, the IOA0 to IOA3 bits in TIOR_1 have selected TGRA_0 compare match or input capture occurrence for the input capture timing while the IOA0 to IOA3 bits in TIOR_2 have selected the TIOC2A rising edge for the input capture timing.

Under these conditions, as TIOR 1 has selected TGRA 0 compare match or input capture occurrence for the input capture timing, the TIOC2A edge is not used for TGRA_1 input capture condition although the I2AE bit in TICCR has been set to 1.

Figure 12.24 Cascaded Operation Example (d)

12.4.5 PWM Modes

In PWM mode, PWM waveforms are output from the output pins. The output level can be selected as 0, 1, or toggle output in response to a compare match of each TGR.

TGR registers settings can be used to output a PWM waveform in the range of 0% to 100% duty.

Designating TGR compare match as the counter clearing source enables the period to be set in that register. All channels can be designated for PWM mode independently. Synchronous operation is also possible.

There are two PWM modes, as described below.

1. PWM mode 1

PWM output is generated from the TIOCA and TIOCC pins by pairing TGRA with TGRB and TGRC with TGRD. The output specified by bits IOA0 to IOA3 and IOC0 to IOC3 in TIOR is output from the TIOCA and TIOCC pins at compare matches A and C, and the output specified by bits IOB0 to IOB3 and IOD0 to IOD3 in TIOR is output at compare matches B and D. The initial output value is the value set in TGRA or TGRC. If the set values of paired TGRs are identical, the output value does not change when a compare match occurs.

In PWM mode 1, a maximum 8-phase PWM output is possible.

2. PWM mode 2

PWM output is generated using one TGR as the cycle register and the others as duty registers. The output specified in TIOR is performed by means of compare matches. Upon counter clearing by a synchronization register compare match, the output value of each pin is the initial value set in TIOR. If the set values of the cycle and duty registers are identical, the output value does not change when a compare match occurs.

In PWM mode 2, a maximum 8-phase PWM output is possible in combination use with synchronous operation.

The correspondence between PWM output pins and registers is shown in table 12.46.

Table 12.46 PWM Output Registers and Output Pins

Note: In PWM mode 2, PWM output is not possible for the TGR register in which the period is set.

(1) Example of PWM Mode Setting Procedure:

Figure 12.25 shows an example of the PWM mode setting procedure.

Figure 12.25 Example of PWM Mode Setting Procedure

(2) Examples of PWM Mode Operation

Figure 12.26 shows an example of PWM mode 1 operation.

In this example, TGRA compare match is set as the TCNT clearing source, 0 is set for the TGRA initial output value and output value, and 1 is set as the TGRB output value.

In this case, the value set in TGRA is used as the period, and the values set in the TGRB registers are used as the duty levels.

Figure 12.26 Example of PWM Mode Operation (1)

Figure 12.27 shows an example of PWM mode 2 operation.

In this example, synchronous operation is designated for channels 0 and 1, TGRB_1 compare match is set as the TCNT clearing source, and 0 is set for the initial output value and 1 for the output value of the other TGR registers (TGRA_0 to TGRD_0, TGRA_1), outputting a 5-phase PWM waveform.

In this case, the value set in TGRB_1 is used as the cycle, and the values set in the other TGRs are used as the duty levels.

Figure 12.27 Example of PWM Mode Operation (2)

Figure 12.28 shows examples of PWM waveform output with 0% duty and 100% duty in PWM mode.

Figure 12.28 Example of PWM Mode Operation (3)

12.4.6 Phase Counting Mode

In phase counting mode, the phase difference between two external clock inputs is detected and TCNT is incremented/decremented accordingly. This mode can be set for channels 1 and 2.

When phase counting mode is set, an external clock is selected as the counter input clock and TCNT operates as an up/down-counter regardless of the setting of bits TPSC0 to TPSC2 and bits CKEG0 and CKEG1 in TCR. However, the functions of bits CCLR0 and CCLR1 in TCR, and of TIOR, TIER, and TGR, are valid, and input capture/compare match and interrupt functions can be used.

This can be used for two-phase encoder pulse input.

If overflow occurs when TCNT is counting up, the TCFV flag in TSR is set; if underflow occurs when TCNT is counting down, the TCFU flag is set.

The TCFD bit in TSR is the count direction flag. Reading the TCFD flag reveals whether TCNT is counting up or down.

Table 12.47 shows the correspondence between external clock pins and channels.

Table 12.47 Phase Counting Mode Clock Input Pins

(1) Example of Phase Counting Mode Setting Procedure

Figure 12.29 shows an example of the phase counting mode setting procedure.

Figure 12.29 Example of Phase Counting Mode Setting Procedure

(2) Examples of Phase Counting Mode Operation

In phase counting mode, TCNT counts up or down according to the phase difference between two external clocks. There are four modes, according to the count conditions.

(a) Phase counting mode 1

Figure 12.30 shows an example of phase counting mode 1 operation, and table 12.48 summarizes the TCNT up/down-count conditions.

Figure 12.30 Example of Phase Counting Mode 1 Operation

Table 12.48 Up/Down-Count Conditions in Phase Counting Mode 1

: Falling edge

(b) Phase counting mode 2

Figure 12.31 shows an example of phase counting mode 2 operation, and table 12.49 summarizes the TCNT up/down-count conditions.

Figure 12.31 Example of Phase Counting Mode 2 Operation

[Legend]

Rising edge

i: Falling edge

(c) Phase counting mode 3

Figure 12.32 shows an example of phase counting mode 3 operation, and table 12.50 summarizes the TCNT up/down-count conditions.

Figure 12.32 Example of Phase Counting Mode 3 Operation

[Legend]

: Rising edge

i: Falling edge

(d) Phase counting mode 4

Figure 12.33 shows an example of phase counting mode 4 operation, and table 12.51 summarizes the TCNT up/down-count conditions.

Figure 12.33 Example of Phase Counting Mode 4 Operation

[Legend]

Rising edge

1: Falling edge

(3) Phase Counting Mode Application Example:

Figure 12.34 shows an example in which channel 1 is in phase counting mode, and channel 1 is coupled with channel 0 to input servo motor 2-phase encoder pulses in order to detect position or speed.

Channel 1 is set to phase counting mode 1, and the encoder pulse A-phase and B-phase are input to TCLKA and TCLKB.

Channel 0 operates with TCNT counter clearing by TGRC_0 compare match; TGRA_0 and TGRC_0 are used for the compare match function and are set with the speed control period and position control period. TGRB_0 is used for input capture, with TGRB_0 and TGRD_0 operating in buffer mode. The channel 1 counter input clock is designated as the TGRB_0 input capture source, and the pulse widths of 2-phase encoder 4-multiplication pulses are detected.

TGRA_1 and TGRB_1 for channel 1 are designated for input capture, and channel 0 TGRA_0 and TGRC_0 compare matches are selected as the input capture source and store the up/down-counter values for the control periods.

This procedure enables the accurate detection of position and speed.

Figure 12.34 Phase Counting Mode Application Example

12.4.7 Reset-Synchronized PWM Mode

In the reset-synchronized PWM mode, three-phase output of positive and negative PWM waveforms that share a common wave transition point can be obtained by combining channels 3 and 4.

When set for reset-synchronized PWM mode, the TIOC3B, TIOC3D, TIOC4A, TIOC4C, TIOC4B, and TIOC4D pins function as PWM output pins and TCNT_3 functions as an upcounter.

Table 12.52 shows the PWM output pins used. Table 12.53 shows the settings of the registers.

Table 12.52 Output Pins for Reset-Synchronized PWM Mode

Channel	Output Pin	Description
3	TIOC3B	PWM output pin 1
	TIOC3D	PWM output pin 1' (negative-phase waveform of PWM output 1)
4	TIOC4A	PWM output pin 2
	TIOC4C	PWM output pin 2' (negative-phase waveform of PWM output 2)
	TIOC4B	PWM output pin 3
	TIOC4D	PWM output pin 3' (negative-phase waveform of PWM output 3)

Table 12.53 Register Settings for Reset-Synchronized PWM Mode

(1) Procedure for Selecting the Reset-Synchronized PWM Mode

Figure 12.35 shows an example of procedure for selecting the reset synchronized PWM mode.

Figure 12.35 Procedure for Selecting Reset-Synchronized PWM Mode

(2) Reset-Synchronized PWM Mode Operation

Figure 12.36 shows an example of operation in the reset-synchronized PWM mode. TCNT_3 and TCNT_4 operate as upcounters. The counter is cleared when a TCNT_3 and TGRA_3 comparematch occurs, and then begins incrementing from H'0000. The PWM output pin output toggles with each occurrence of a TGRB_3, TGRA_4, TGRB_4 compare-match, and upon counter clears.

Figure 12.36 Reset-Synchronized PWM Mode Operation Example (When TOCR's OLSN = 1 and OLSP = 1)

12.4.8 Complementary PWM Mode

In the complementary PWM mode, three-phase output of non-overlapping positive and negative PWM waveforms can be obtained by combining channels 3 and 4. PWM waveforms without nonoverlapping interval is also available.

In complementary PWM mode, TIOC3B, TIOC3D, TIOC4A, TIOC4B, TIOC4C, and TIOC4D pins function as PWM output pins, the TIOC3A pin can be set for toggle output synchronized with the PWM period. TCNT_3 and TCNT_4 function as up/down counters.

Table 12.54 shows the PWM output pins used. Table 12.55 shows the settings of the registers used.

Table 12.54 Output Pins for Complementary PWM Mode

Note: * Avoid setting the TIOC3C pin as a timer I/O pin in the complementary PWM mode.

Figure 12.37 Block Diagram of Channels 3 and 4 in Complementary PWM Mode

(1) Example of Complementary PWM Mode Setting Procedure

An example of the complementary PWM mode setting procedure is shown in figure 12.38.

Figure 12.38 Example of Complementary PWM Mode Setting Procedure

(2) Outline of Complementary PWM Mode Operation

In complementary PWM mode, 6-phase PWM output is possible. Figure 12.39 illustrates counter operation in complementary PWM mode, and figure 12.40 shows an example of complementary PWM mode operation.

(a) Counter Operation

In complementary PWM mode, three counters—TCNT_3, TCNT_4, and TCNTS—perform up/down-count operations.

TCNT_3 is automatically initialized to the value set in TDDR when complementary PWM mode is selected and the CST bit in TSTR is 0.

When the CST bit is set to 1, TCNT_{_3} counts up to the value set in TGRA_{_3}, then switches to down-counting when it matches TGRA_3. When the TCNT3 value matches TDDR, the counter switches to up-counting, and the operation is repeated in this way.

TCNT_4 is initialized to H'0000.

When the CST bit is set to 1, TCNT 4 counts up in synchronization with TCNT 3, and switches to down-counting when it matches TCDR. On reaching H'0000, TCNT4 switches to up-counting, and the operation is repeated in this way.

TCNTS is a read-only counter. It need not be initialized.

When TCNT_3 matches TCDR during TCNT_3 and TCNT_4 up/down-counting, down-counting is started, and when TCNTS matches TCDR, the operation switches to up-counting. When TCNTS matches TGRA_3, it is cleared to H'0000.

When TCNT₋₄ matches TDDR during TCNT₋₃ and TCNT₋₄ down-counting, up-counting is started, and when TCNTS matches TDDR, the operation switches to down-counting. When TCNTS reaches H'0000, it is set with the value in TGRA_3.

TCNTS is compared with the compare register and temporary register in which the PWM duty is set during the count operation only.

Figure 12.39 Complementary PWM Mode Counter Operation

(b) Register Operation

In complementary PWM mode, nine registers are used, comprising compare registers, buffer registers, and temporary registers. Figure 12.40 shows an example of complementary PWM mode operation.

The registers which are constantly compared with the counters to perform PWM output are TGRB_3, TGRA_4, and TGRB_4. When these registers match the counter, the value set in bits OLSN and OLSP in the timer output control register (TOCR) is output.

The buffer registers for these compare registers are TGRD_3, TGRC_4, and TGRD_4.

Between a buffer register and compare register there is a temporary register. The temporary registers cannot be accessed by the CPU.

Data in a compare register is changed by writing the new data to the corresponding buffer register. The buffer registers can be read or written at any time.

The data written to a buffer register is constantly transferred to the temporary register in the Ta interval. Data is not transferred to the temporary register in the Tb interval. Data written to a buffer register in this interval is transferred to the temporary register at the end of the Tb interval.

The value transferred to a temporary register is transferred to the compare register when TCNTS for which the Tb interval ends matches TGRA_3 when counting up, or H'0000 when counting down. The timing for transfer from the temporary register to the compare register can be selected with bits MD3 to MD0 in the timer mode register (TMDR). Figure 12.40 shows an example in which the mode is selected in which the change is made in the trough.

In the Tb interval (Tb1 in figure 12.40) in which data transfer to the temporary register is not performed, the temporary register has the same function as the compare register, and is compared

with the counter. In this interval, therefore, there are two compare match registers for one-phase output, with the compare register containing the pre-change data, and the temporary register containing the new data. In this interval, the three counters—TCNT 3, TCNT 4, and TCNTS and two registers—compare register and temporary register—are compared, and PWM output controlled accordingly.

Figure 12.40 Example of Complementary PWM Mode Operation

(c) Initialization

In complementary PWM mode, there are six registers that must be initialized. In addition, there is a register that specifies whether to generate dead time (it should be used only when dead time generation should be disabled).

Before setting complementary PWM mode with bits MD3 to MD0 in the timer mode register (TMDR), the following initial register values must be set.

TGRC_3 operates as the buffer register for TGRA_3, and should be set with 1/2 the PWM carrier cycle + dead time Td. The timer cycle buffer register (TCBR) operates as the buffer register for the timer cycle data register (TCDR), and should be set with 1/2 the PWM carrier cycle. Set dead time Td in the timer dead time data register (TDDR).

When dead time is not needed, the TDER bit in the timer dead time enable register (TDER) should be cleared to 0, TGRC_3 and TGRA_3 should be set to 1/2 the PWM carrier cycle + 1, and TDDR should be set to 1.

Set the respective initial PWM duty values in buffer registers TGRD_3, TGRC_4, and TGRD_4.

The values set in the five buffer registers excluding TDDR are transferred simultaneously to the corresponding compare registers when complementary PWM mode is set.

Set TCNT_4 to H'0000 before setting complementary PWM mode.

dead time Td set in TDDR. When dead time generation is disabled by TDER, TGRC_3 must be set to $1/2$ the PWM carrier cycle $+1$.

(d) PWM Output Level Setting

In complementary PWM mode, the PWM pulse output level is set with bits OLSN and OLSP in timer output control register 1 (TOCR1) or bits OLS1P to OLS3P and OLS1N to OLS3N in timer output control register 2 (TOCR2).

The output level can be set for each of the three positive phases and three negative phases of 6 phase output.

Complementary PWM mode should be cleared before setting or changing output levels.

(e) Dead Time Setting

In complementary PWM mode, PWM pulses are output with a non-overlapping relationship between the positive and negative phases. This non-overlap time is called the dead time.

The non-overlap time is set in the timer dead time data register (TDDR). The value set in TDDR is used as the TCNT_3 counter start value, and creates non-overlap between TCNT_3 and TCNT_4. Complementary PWM mode should be cleared before changing the contents of TDDR.

(f) Dead Time Suppressing

Dead time generation is suppressed by clearing the TDER bit in the timer dead time enable register (TDER) to 0. TDER can be cleared to 0 only when 0 is written to it after reading TDER $=$ 1.

TGRA β and TGRC β should be set to 1/2 PWM carrier cycle $+ 1$ and the timer dead time data register (TDDR) should be set to 1.

By the above settings, PWM waveforms without dead time can be obtained. Figure 12.41 shows an example of operation without dead time.

Figure 12.41 Example of Operation without Dead Time

(g) PWM Cycle Setting

In complementary PWM mode, the PWM pulse cycle is set in two registers—TGRA_3, in which the TCNT_3 upper limit value is set, and TCDR, in which the TCNT_4 upper limit value is set. The settings should be made so as to achieve the following relationship between these two registers:

With dead time: TGRA_3 set value = TCDR set value + TDDR set value Without dead time: TGRA 3 set value = TCDR set value $+1$

The TGRA_3 and TCDR settings are made by setting the values in buffer registers TGRC_3 and TCBR. The values set in TGRC_3 and TCBR are transferred simultaneously to TGRA_3 and TCDR in accordance with the transfer timing selected with bits MD3 to MD0 in the timer mode register (TMDR).

The updated PWM cycle is reflected from the next cycle when the data update is performed at the crest, and from the current cycle when performed in the trough. Figure 12.42 illustrates the operation when the PWM cycle is updated at the crest.

See the following section, Register Data Updating, for the method of updating the data in each buffer register.

Figure 12.42 Example of PWM Cycle Updating

(h) Register Data Updating

In complementary PWM mode, the buffer register is used to update the data in a compare register. The update data can be written to the buffer register at any time. There are five PWM duty and carrier cycle registers that have buffer registers and can be updated during operation.

There is a temporary register between each of these registers and its buffer register. When subcounter TCNTS is not counting, if buffer register data is updated, the temporary register value is also rewritten. Transfer is not performed from buffer registers to temporary registers when TCNTS is counting; in this case, the value written to a buffer register is transferred after TCNTS halts.

The temporary register value is transferred to the compare register at the data update timing set with bits MD3 to MD0 in the timer mode register (TMDR). Figure 12.43 shows an example of data updating in complementary PWM mode. This example shows the mode in which data updating is performed at both the counter crest and trough.

When rewriting buffer register data, a write to TGRD 4 must be performed at the end of the update. Data transfer from the buffer registers to the temporary registers is performed simultaneously for all five registers after the write to TGRD 4.

A write to TGRD 4 must be performed after writing data to the registers to be updated, even when not updating all five registers, or when updating the TGRD_4 data. In this case, the data written to TGRD_4 should be the same as the data prior to the write operation.

Figure 12.43 Example of Data Update in Complementary PWM Mode

SH7261 Group

(i) Initial Output in Complementary PWM Mode

In complementary PWM mode, the initial output is determined by the setting of bits OLSN and OLSP in timer output control register 1 (TOCR1) or bits OLS1N to OLS3N and OLS1P to OLS3P in timer output control register 2 (TOCR2).

This initial output is the PWM pulse non-active level, and is output from when complementary PWM mode is set with the timer mode register (TMDR) until TCNT_4 exceeds the value set in the dead time register (TDDR). Figure 12.44 shows an example of the initial output in complementary PWM mode.

An example of the waveform when the initial PWM duty value is smaller than the TDDR value is shown in figure 12.45.

Figure 12.44 Example of Initial Output in Complementary PWM Mode (1)

Figure 12.45 Example of Initial Output in Complementary PWM Mode (2)

(j) 10. Complementary PWM Mode PWM Output Generation Method

In complementary PWM mode, 3-phase output is performed of PWM waveforms with a nonoverlap time between the positive and negative phases. This non-overlap time is called the dead time.

A PWM waveform is generated by output of the output level selected in the timer output control register in the event of a compare-match between a counter and data register. While TCNTS is counting, data register and temporary register values are simultaneously compared to create consecutive PWM pulses from 0 to 100%. The relative timing of on and off compare-match occurrence may vary, but the compare-match that turns off each phase takes precedence to secure the dead time and ensure that the positive phase and negative phase on times do not overlap. Figures 12.46 to 12.48 show examples of waveform generation in complementary PWM mode.

The positive phase/negative phase off timing is generated by a compare-match with the solid-line counter, and the on timing by a compare-match with the dotted-line counter operating with a delay of the dead time behind the solid-line counter. In the T1 period, compare-match **a** that turns off the negative phase has the highest priority, and compare-matches occurring prior to **a** are ignored. In the T2 period, compare-match **c** that turns off the positive phase has the highest priority, and compare-matches occurring prior to **c** are ignored.

In normal cases, compare-matches occur in the order $\mathbf{a} \to \mathbf{b} \to \mathbf{c} \to \mathbf{d}$ (or $\mathbf{c} \to \mathbf{d} \to \mathbf{a'} \to \mathbf{b'}$), as shown in figure 12.46.

If compare-matches deviate from the $\mathbf{a} \to \mathbf{b} \to \mathbf{c} \to \mathbf{d}$ order, since the time for which the negative phase is off is less than twice the dead time, the figure shows the positive phase is not being turned on. If compare-matches deviate from the $c \rightarrow d \rightarrow a' \rightarrow b'$ order, since the time for which the positive phase is off is less than twice the dead time, the figure shows the negative phase is not being turned on.

If compare-match **c** occurs first following compare-match **a**, as shown in figure 12.47, comparematch **b** is ignored, and the negative phase is turned off by compare-match **d**. This is because turning off of the positive phase has priority due to the occurrence of compare-match **c** (positive phase off timing) before compare-match **b** (positive phase on timing) (consequently, the waveform does not change since the positive phase goes from off to off).

Similarly, in the example in figure 12.48, compare-match **a'** with the new data in the temporary register occurs before compare-match **c**, but other compare-matches occurring up to **c**, which turns off the positive phase, are ignored. As a result, the negative phase is not turned on.

Thus, in complementary PWM mode, compare-matches at turn-off timings take precedence, and turn-on timing compare-matches that occur before a turn-off timing compare-match are ignored.

Figure 12.46 Example of Complementary PWM Mode Waveform Output (1)

Figure 12.47 Example of Complementary PWM Mode Waveform Output (2)

Figure 12.48 Example of Complementary PWM Mode Waveform Output (3)

Figure 12.49 Example of Complementary PWM Mode 0% and 100% Waveform Output (1)

Figure 12.50 Example of Complementary PWM Mode 0% and 100% Waveform Output (2)

Figure 12.51 Example of Complementary PWM Mode 0% and 100% Waveform Output (3)

Figure 12.52 Example of Complementary PWM Mode 0% and 100% Waveform Output (4)

Figure 12.53 Example of Complementary PWM Mode 0% and 100% Waveform Output (5)

(k) 11. Complementary PWM Mode 0% and 100% Duty Output

In complementary PWM mode, 0% and 100% duty cycles can be output as required. Figures 12.49 to 12.53 show output examples.

100% duty output is performed when the data register value is set to H'0000. The waveform in this case has a positive phase with a 100% on-state. 0% duty output is performed when the data register value is set to the same value as TGRA_3. The waveform in this case has a positive phase with a 100% off-state.

On and off compare-matches occur simultaneously, but if a turn-on compare-match and turn-off compare-match for the same phase occur simultaneously, both compare-matches are ignored and the waveform does not change.

(l) 12. Toggle Output Synchronized with PWM Cycle

In complementary PWM mode, toggle output can be performed in synchronization with the PWM carrier cycle by setting the PSYE bit to 1 in the timer output control register (TOCR). An example of a toggle output waveform is shown in figure 12.54.

This output is toggled by a compare-match between TCNT 3 and TGRA 3 and a compare-match between TCNT4 and H'0000.

The output pin for this toggle output is the TIOC3A pin. The initial output is 1.

Figure 12.54 Example of Toggle Output Waveform Synchronized with PWM Output

(m) Counter Clearing by Another Channel

In complementary PWM mode, by setting a mode for synchronization with another channel by means of the timer synchronous register (TSYR), and selecting synchronous clearing with bits CCLR2 to CCLR0 in the timer control register (TCR), it is possible to have TCNT_3, TCNT_4, and TCNTS cleared by another channel.

Figure 12.55 illustrates the operation.

Use of this function enables counter clearing and restarting to be performed by means of an external signal.

Figure 12.55 Counter Clearing Synchronized with Another Channel

(n) Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

Setting the WRE bit in TWCR to 1 suppresses initial output when synchronous counter clearing occurs in the Tb interval at the trough in complementary PWM mode and controls abrupt change in duty cycle at synchronous counter clearing.

Initial output suppression is applicable only when synchronous clearing occurs in the Tb interval at the trough as indicated by (10) or (11) in figure 12.56. When synchronous clearing occurs outside that interval, the initial value specified by the OLS bits in TOCR is output. Even in the Tb interval at the trough, if synchronous clearing occurs in the initial value output period (indicated by (1) in figure 12.56) immediately after the counters start operation, initial value output is not suppressed.

When using the initial output suppression function, make sure to set compare registers TGRB 3, TGRA_4, and TGRB_4 to a value twice or more the setting of dead time data register TDDR. If synchronous clearing occurs with the compare registers set to a value less than twice the setting of TDDR, the PWM output dead time may be too short (or nonexistent) or illegal active-level PWM negative-phase output may occur during the initial output suppression interval. For details, see section 12.7.23, Notes on Output Waveform Control During Synchronous Counter Clearing in Complementary PWM Mode.

In the MTU2, synchronous clearing generated in channels 0 to 2 in the MTU2 can cause counter clearing.

Figure 12.56 Timing for Synchronous Counter Clearing

• Example of Procedure for Setting Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

An example of the procedure for setting output waveform control at synchronous counter clearing in complementary PWM mode is shown in figure 12.57.

Figure 12.57 Example of Procedure for Setting Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

• Examples of Output Waveform Control at Synchronous Counter Clearing in Complementary PWM Mode

Figures 12.58 to 12.61 show examples of output waveform control in which the MTU2 operates in complementary PWM mode and synchronous counter clearing is generated while the WRE bit in TWCR is set to 1. In the examples shown in figures 12.58 to 12.61, synchronous counter clearing occurs at timing (3), (6), (8), and (11) shown in figure 12.56, respectively.

Figure 12.58 Example of Synchronous Clearing in Dead Time during Up-Counting (Timing (3) in Figure 12.56; Bit WRE of TWCR in MTU2 is 1)

Figure 12.59 Example of Synchronous Clearing in Interval Tb at Crest (Timing (6) in Figure 12.56; Bit WRE of TWCR in MTU2 is 1)

Figure 12.60 Example of Synchronous Clearing in Dead Time during Down-Counting (Timing (8) in Figure 12.56; Bit WRE of TWCR is 1)

(o) Counter Clearing by TGRA_3 Compare Match

In complementary PWM mode, by setting the CCE bit in the timer waveform control register (TWCR), it is possible to have TCNT_3, TCNT_4, and TCNTS cleared by TGRA_3 compare match.

Figure 12.62 illustrates an operation example.

Notes: 1. Use this function only in complementary PWM mode 1 (transfer at crest).

- 2. Do not specify synchronous clearing by another channel (do not set the SYNC0 to SYNC4 bits in the timer synchronous register (TSYR) to 1).
- 3. Do not set the PWM duty value to H'0000.
- 4. Do not set the PSYE bit in timer output control register 1 (TOCR1) to 1.

Figure 12.62 Example of Counter Clearing Operation by TGRA_3 Compare Match

(p) Example of AC Synchronous Motor (Brushless DC Motor) Drive Waveform Output

In complementary PWM mode, a brushless DC motor can easily be controlled using the timer gate control register (TGCR). Figures 12.63 to 12.66 show examples of brushless DC motor drive waveforms created using TGCR.

When output phase switching for a 3-phase brushless DC motor is performed by means of external signals detected with a Hall element, etc., clear the FB bit in TGCR to 0. In this case, the external signals indicating the polarity position are input to channel 0 timer input pins TIOC0A, TIOC0B, and TIOC0C (set with PFC). When an edge is detected at pin TIOC0A, TIOC0B, or TIOC0C, the output on/off state is switched automatically.

When the FB bit is 1, the output on/off state is switched when the UF, VF, or WF bit in TGCR is cleared to 0 or set to 1.

The drive waveforms are output from the complementary PWM mode 6-phase output pins. With this 6-phase output, in the case of on output, it is possible to use complementary PWM mode output and perform chopping output by setting the N bit or P bit to 1. When the N bit or P bit is 0, level output is selected.

The 6-phase output active level (on output level) can be set with the OLSN and OLSP bits in the timer output control register (TOCR) regardless of the setting of the N and P bits.

Figure 12.63 Example of Output Phase Switching by External Input (1)

Figure 12.64 Example of Output Phase Switching by External Input (2)

Figure 12.65 Example of Output Phase Switching by Means of UF, VF, WF Bit Settings (1)

Figure 12.66 Example of Output Phase Switching by Means of UF, VF, WF Bit Settings (2)

(q) A/D Converter Start Request Setting

In complementary PWM mode, an A/D converter start request can be issued using a TGRA_3 compare-match, TCNT_4 underflow (trough), or compare-match on a channel other than channels 3 and 4.

When start requests using a TGRA 3 compare-match are specified, A/D conversion can be started at the crest of the TCNT_3 count.

A/D converter start requests can be set by setting the TTGE bit to 1 in the timer interrupt enable register (TIER). To issue an A/D converter start request at a TCNT_4 underflow (trough), set the TTGE2 bit in TIER_4 to 1.

(3) Interrupt Skipping in Complementary PWM Mode:

Interrupts TGIA 3 (at the crest) and TCIV 4 (at the trough) in channels 3 and 4 can be skipped up to seven times by making settings in the timer interrupt skipping set register (TITCR).

Transfers from a buffer register to a temporary register or a compare register can be skipped in coordination with interrupt skipping by making settings in the timer buffer transfer register (TBTER). For the linkage with buffer registers, refer to description (c), Buffer Transfer Control Linked with Interrupt Skipping, below.

A/D converter start requests generated by the A/D converter start request delaying function can also be skipped in coordination with interrupt skipping by making settings in the timer A/D converter request control register (TADCR). For the linkage with the A/D converter start request delaying function, refer to section 12.4.9, A/D Converter Start Request Delaying Function.

The setting of the timer interrupt skipping setting register (TITCR) must be done while the TGIA_3 and TCIV_4 interrupt requests are disabled by the settings of registers TIER_3 and TIER_4 along with under the conditions in which TGFA_3 and TCFV_4 flag settings by compare match never occur. Before changing the skipping count, be sure to clear the T3AEN and T4VEN bits to 0 to clear the skipping counter.

(a) Example of Interrupt Skipping Operation Setting Procedure

Figure 12.67 shows an example of the interrupt skipping operation setting procedure. Figure 12.68 shows the periods during which interrupt skipping count can be changed.

Figure 12.67 Example of Interrupt Skipping Operation Setting Procedure

Figure 12.68 Periods during which Interrupt Skipping Count can be Changed

(b) Example of Interrupt Skipping Operation

Figure 12.69 shows an example of TGIA_3 interrupt skipping in which the interrupt skipping count is set to three by the 3ACOR bit and the T3AEN bit is set to 1 in the timer interrupt skipping set register (TITCR).

Figure 12.69 Example of Interrupt Skipping Operation

(c) Buffer Transfer Control Linked with Interrupt Skipping

In complementary PWM mode, whether to transfer data from a buffer register to a temporary register and whether to link the transfer with interrupt skipping can be specified with the BTE1 and BTE0 bits in the timer buffer transfer set register (TBTER).

Figure 12.70 shows an example of operation when buffer transfer is suppressed (BTE1 = 0 and $BTE0 = 1$). While this setting is valid, data is not transferred from the buffer register to the temporary register.

Figure 12.71 shows an example of operation when buffer transfer is linked with interrupt skipping $(BTE1 = 1$ and $BET0 = 0)$. While this setting is valid, data is not transferred from the buffer register outside the buffer transfer-enabled period.

Note that the buffer transfer-enabled period depends on the T3AEN and T4VEN bit settings in the timer interrupt skipping set register (TITCR). Figure 12.72 shows the relationship between the T3AEN and T4VEN bit settings in TITCR and buffer transfer-enabled period.

Note: This function must always be used in combination with interrupt skipping. When interrupt skipping is disabled (the T3AEN and T4VEN bits in the timer interrupt skipping set register (TITCR) are cleared to 0 or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0), make sure that buffer transfer is not linked with interrupt skipping (clear the BTE1 bit in the timer buffer transfer set register (TBTER) to 0). If buffer transfer is linked with interrupt skipping while interrupt skipping is disabled, buffer transfer is never performed.

Figure 12.70 Example of Operation when Buffer Transfer is Suppressed (BTE1 = 0 and BTE0 = 1)

(2)When rewriting the buffer register after passing 1 carrier cycle from TGIA_3 interrupt

Figure 12.71 Example of Operation when Buffer Transfer is Linked with Interrupt Skipping (BTE1 = 1 and BTE0 = 0)

Figure 12.72 Relationship between Bits T3AEN and T4VEN in TITCR and Buffer Transfer-Enabled Period

(4) Complementary PWM Mode Output Protection Function

Complementary PWM mode output has the following protection functions.

(a) Register and counter miswrite prevention function

With the exception of the buffer registers, which can be rewritten at any time, access by the CPU can be enabled or disabled for the mode registers, control registers, compare registers, and counters used in complementary PWM mode by means of the RWE bit in the timer read/write enable register (TRWER). The applicable registers are some (21 in total) of the registers in channels 3 and 4 shown in the following:

• TCR 3 and TCR 4, TMDR 3 and TMDR 4, TIORH 3 and TIORH 4, TIORL 3 and TIORL_4, TIER_3 and TIER_4, TCNT_3 and TCNT_4, TGRA_3 and TGRA_4, TGRB_3 and TGRB_4, TOER, TOCR, TGCR, TCDR, and TDDR.

This function enables miswriting due to CPU runaway to be prevented by disabling CPU access to the mode registers, control registers, and counters. When the applicable registers are read in the access-disabled state, undefined values are returned. Writing to these registers is ignored.

12.4.9 A/D Converter Start Request Delaying Function

A/D converter start requests can be issued in channel 4 by making settings in the timer A/D converter start request control register (TADCR), timer A/D converter start request cycle set registers (TADCORA_4 and TADCORB_4), and timer A/D converter start request cycle set buffer registers (TADCOBRA 4 and TADCOBRB 4).

The A/D converter start request delaying function compares TCNT_4 with TADCORA_4 or TADCORB_4, and when their values match, the function issues a respective A/D converter start request (TRG4AN or TRG4BN).

A/D converter start requests (TRG4AN and TRG4BN) can be skipped in coordination with interrupt skipping by making settings in the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in **TADCR**

(a) Example of Procedure for Specifying A/D Converter Start Request Delaying Function

Figure 12.73 shows an example of procedure for specifying the A/D converter start request delaying function.

Figure 12.73 Example of Procedure for Specifying A/D Converter Start Request Delaying Function

(b) Basic Operation Example of A/D Converter Start Request Delaying Function

Figure 12.74 shows a basic example of A/D converter request signal (TRG4AN) operation when the trough of TCNT 4 is specified for the buffer transfer timing and an A/D converter start request signal is output during TCNT 4 down-counting.

Figure 12.74 Basic Example of A/D Converter Start Request Signal (TRG4AN) Operation

(c) Buffer Transfer

The data in the timer A/D converter start request cycle set registers (TADCORA_4 and TADCORB_4) is updated by writing data to the timer A/D converter start request cycle set buffer registers (TADCOBRA_4 and TADCOBRB_4). Data is transferred from the buffer registers to the respective cycle set registers at the timing selected with the BF1 and BF0 bits in the timer A/D converter start request control register (TADCR_4).

(d) A/D Converter Start Request Delaying Function Linked with Interrupt Skipping

A/D converter start requests (TRG4AN and TRG4BN) can be issued in coordination with interrupt skipping by making settings in the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in the timer A/D converter start request control register (TADCR).

Figure 12.75 shows an example of A/D converter start request signal (TRG4AN) operation when TRG4AN output is enabled during TCNT_4 up-counting and down-counting and A/D converter start requests are linked with interrupt skipping.

Figure 12.76 shows another example of A/D converter start request signal (TRG4AN) operation when TRG4AN output is enabled during TCNT_4 up-counting and A/D converter start requests are linked with interrupt skipping.

Note: This function must be used in combination with interrupt skipping. When interrupt skipping is disabled (the T3AEN and T4VEN bits in the timer interrupt skipping set register (TITCR) are cleared to 0 or the skipping count set bits (3ACOR and 4VCOR) in TITCR are cleared to 0), make sure that A/D converter start requests are not linked with interrupt skipping (clear the ITA3AE, ITA4VE, ITB3AE, and ITB4VE bits in the timer A/D converter start request control register (TADCR) to 0).

Figure 12.75 Example of A/D Converter Start Request Signal (TRG4AN) Operation Linked with Interrupt Skipping

Figure 12.76 Example of A/D Converter Start Request Signal (TRG4AN) Operation Linked with Interrupt Skipping

12.4.10 External Pulse Width Measurement

The pulse widths of up to three external input lines can be measured in channel 5.

(1) Example of External Pulse Width Measurement Setting Procedure

Figure 12.77 Example of External Pulse Width Measurement Setting Procedure

(2) Example of External Pulse Width Measurement

Figure 12.78 Example of External Pulse Width Measurement (Measuring High Pulse Width)

12.4.11 Dead Time Compensation

By measuring the delay of the output waveform and reflecting it to duty, the external pulse width measurement function can be used as the dead time compensation function while the complementary PWM is in operation.

Figure 12.79 Delay in Dead Time in Complementary PWM Operation

(1) Example of Dead Time Compensation Setting Procedure

Figure 12.80 shows an example of dead time compensation setting procedure by using three counters in channel 5.

Figure 12.80 Example of Dead Time Compensation Setting Procedure

Figure 12.81 Example of Motor Control Circuit Configuration

12.4.12 TCNT Capture at Crest and/or Trough in Complementary PWM Operation

The TCNT value is captured in TGR at either the crest or trough or at both the crest and trough during complementary PWM operation. The timing for capturing in TGR can be selected by TIOR.

Figure 12.82 is an operating example in which TCNT is used as a free-running counter without being cleared, and the TCNT value is captured in TGR at the specified timing (either crest or trough, or both crest and trough).

Figure 12.82 TCNT Capturing at Crest and/or Trough in Complementary PWM Operation

12.5 Interrupt Sources

12.5.1 Interrupt Sources and Priorities

There are three kinds of MTU2 interrupt source; TGR input capture/compare match, TCNT overflow, and TCNT underflow. Each interrupt source has its own status flag and enable/disabled bit, allowing the generation of interrupt request signals to be enabled or disabled individually.

When an interrupt request is generated, the corresponding status flag in TSR is set to 1. If the corresponding enable/disable bit in TIER is set to 1 at this time, an interrupt is requested. The interrupt request is cleared by clearing the status flag to 0.

Relative channel priorities can be changed by the interrupt controller, however the priority order within a channel is fixed. For details, see section 6, Interrupt Controller (INTC).

Table 12.57 lists the MTU2 interrupt sources.

Table 12.57 MTU2 Interrupts

Note: This table shows the initial state immediately after a reset. The relative channel priorities can be changed by the interrupt controller.

(1) Input Capture/Compare Match Interrupt

An interrupt is requested if the TGIE bit in TIER is set to 1 when the TGF flag in TSR is set to 1 by the occurrence of a TGR input capture/compare match on a particular channel. The interrupt request is cleared by clearing the TGF flag to 0. The MTU2 has 21 input capture/compare match interrupts, six for channel 0, four each for channels 3 and 4, two each for channels 1 and 2, and three for channel 5. The TGFE_0 and TGFF_0 flags in channel 0 are not set by the occurrence of an input capture.

(2) Overflow Interrupt

An interrupt is requested if the TCIEV bit in TIER is set to 1 when the TCFV flag in TSR is set to 1 by the occurrence of TCNT overflow on a channel. The interrupt request is cleared by clearing the TCFV flag to 0. The MTU2 has five overflow interrupts, one for each channel.

(3) Underflow Interrupt

An interrupt is requested if the TCIEU bit in TIER is set to 1 when the TCFU flag in TSR is set to 1 by the occurrence of TCNT underflow on a channel. The interrupt request is cleared by clearing the TCFU flag to 0. The MTU2 has two underflow interrupts, one each for channels 1 and 2.

12.5.2 DMAC Activation

The DMAC can be activated by the TGRA input capture/compare match interrupt in each channel. For details, see section 11, Direct Memory Access Controller (DMAC).

In the MTU2, a total of five TGRA input capture/compare match interrupts can be used as DMAC activation sources, one each for channels 0 to 4.

12.5.3 A/D Converter Activation

The A/D converter can be activated by one of the following three methods in the MTU2. Table 12.58 shows the relationship between interrupt sources and A/D converter start request signals.

(1) A/D Converter Activation by TGRA Input Capture/Compare Match or at TCNT_4 Trough in Complementary PWM Mode

The A/D converter can be activated by the occurrence of a TGRA input capture/compare match in each channel. In addition, if complementary PWM operation is performed while the TTGE2 bit in TIER_4 is set to 1, the A/D converter can be activated at the trough of $TCNT_4$ count (TCNT_4 = H'0000).

A/D converter start request signal TRGAN is issued to the A/D converter under either one of the following conditions.

- When the TGFA flag in TSR is set to 1 by the occurrence of a TGRA input capture/compare match on a particular channel while the TTGE bit in TIER is set to 1
- When the TCNT 4 count reaches the trough (TCNT $4 = H'0000$) during complementary PWM operation while the TTGE2 bit in TIER_4 is set to 1

When either condition is satisfied, if A/D converter start signal TRGAN from the MTU2 is selected as the trigger in the A/D converter, A/D conversion will start.

(2) A/D Converter Activation by Compare Match between TCNT_0 and TGRE_0

The A/D converter can be activated by generating A/D converter start request signal TRG0N when a compare match occurs between TCNT 0 and TGRE 0 in channel 0.

When the TGFE flag in TSR2 0 is set to 1 by the occurrence of a compare match between TCNT 0 and TGRE 0 in channel 0 while the TTGE2 bit in TIER2_0 is set to 1, A/D converter start request TGR0N is issued to the A/D converter. If A/D converter start signal TGR0N from the MTU2 is selected as the trigger in the A/D converter, A/D conversion will start.

(3) A/D Converter Activation by A/D Converter Start Request Delaying Function

The A/D converter can be activated by generating A/D converter start request signal TRG4AN or TRG4BN when the TCNT_4 count matches the TADCORA or TADCORB value if the UT4AE, DT4AE, UT4BE, or DT4BE bit in the A/D converter start request control register (TADCR) is set to 1. For details, refer to section 12.4.9, A/D Converter Start Request Delaying Function.

A/D conversion will start if A/D converter start signal TRG4AN from the MTU2 is selected as the trigger in the A/D converter when TRG4AN is generated or if TRG4BN from the MTU2 is selected as the trigger in the A/D converter when TRG4BN is generated.

Table 12.58 Interrupt Sources and A/D Converter Start Request Signals

12.6 Operation Timing

12.6.1 Input/Output Timing

(1) TCNT Count Timing

Figures 12.83 and 12.84 show TCNT count timing in internal clock operation, and figure 12.85 shows TCNT count timing in external clock operation (normal mode), and figure 12.86 shows TCNT count timing in external clock operation (phase counting mode).

Figure 12.85 Count Timing in External Clock Operation (Channels 0 to 4)

Figure 12.86 Count Timing in External Clock Operation (Phase Counting Mode)

(2) Output Compare Output Timing

A compare match signal is generated in the final state in which TCNT and TGR match (the point at which the count value matched by TCNT is updated). When a compare match signal is generated, the output value set in TIOR is output at the output compare output pin (TIOC pin). After a match between TCNT and TGR, the compare match signal is not generated until the TCNT input clock is generated.

Figure 12.87 shows output compare output timing (normal mode and PWM mode) and figure 12.88 shows output compare output timing (complementary PWM mode and reset synchronous PWM mode).

Figure 12.87 Output Compare Output Timing (Normal Mode/PWM Mode)

Figure 12.88 Output Compare Output Timing (Complementary PWM Mode/Reset Synchronous PWM Mode)

(3) Input Capture Signal Timing

Figure 12.89 shows input capture signal timing.

Figure 12.89 Input Capture Input Signal Timing

(4) Timing for Counter Clearing by Compare Match/Input Capture

Figures 12.90 and 12.91 show the timing when counter clearing on compare match is specified, and figure 12.92 shows the timing when counter clearing on input capture is specified.

Figure 12.90 Counter Clear Timing (Compare Match) (Channel 0 to Channel 4)

Figure 12.91 Counter Clear Timing (Compare Match) (Channel 5)

Figure 12.92 Counter Clear Timing (Input Capture) (Channel 0 to Channel 5)

(5) Buffer Operation Timing

Figures 12.93 to 12.95 show the timing in buffer operation.

Figure 12.94 Buffer Operation Timing (Input Capture)

(6) Buffer Transfer Timing (Complementary PWM Mode)

Figures 12.96 to 12.98 show the buffer transfer timing in complementary PWM mode.

Figure 12.96 Transfer Timing from Buffer Register to Temporary Register (TCNTS Stop)

Figure 12.97 Transfer Timing from Buffer Register to Temporary Register (TCNTS Operating)

Figure 12.98 Transfer Timing from Temporary Register to Compare Register

12.6.2 Interrupt Signal Timing

(1) TGF Flag Setting Timing in Case of Compare Match

Figures 12.99 and 12.100 show the timing for setting of the TGF flag in TSR on compare match, and TGI interrupt request signal timing.

Figure 12.99 TGI Interrupt Timing (Compare Match) (Channels 0 to 4)

Figure 12.100 TGI Interrupt Timing (Compare Match) (Channel 5)

(2) TGF Flag Setting Timing in Case of Input Capture

Figures 12.101 and 12.102 show the timing for setting of the TGF flag in TSR on input capture, and TGI interrupt request signal timing.

Figure 12.101 TGI Interrupt Timing (Input Capture) (Channels 0 to 4)

Figure 12.102 TGI Interrupt Timing (Input Capture) (Channel 5)

(3) TCFV Flag/TCFU Flag Setting Timing

Figure 12.103 shows the timing for setting of the TCFV flag in TSR on overflow, and TCIV interrupt request signal timing.

Figure 12.104 shows the timing for setting of the TCFU flag in TSR on underflow, and TCIU interrupt request signal timing.

Figure 12.103 TCIV Interrupt Setting Timing

Figure 12.104 TCIU Interrupt Setting Timing

(4) Status Flag Clearing Timing

After a status flag is read as 1 by the CPU, it is cleared by writing 0 to it. Figures 12.105 and 12.106 show the timing for status flag clearing by the CPU, and figure 12.107 show the timing for status flag clearing by the DMAC.

Figure 12.105 Timing for Status Flag Clearing by CPU (Channels 0 to 4)

Figure 12.106 Timing for Status Flag Clearing by CPU (Channel 5)

Figure 12.107 Timing for Status Flag Clearing by DMAC Activation (Channels 0 to 4)

12.7 Usage Notes

12.7.1 Module Standby Mode Setting

MTU2 operation can be disabled or enabled using the standby control register. The initial setting is for MTU2 operation to be halted. Register access is enabled by clearing module standby mode. For details, refer to section 27, Power-Down Modes.

12.7.2 Input Clock Restrictions

The input clock pulse width must be at least 1.5 states in the case of single-edge detection, and at least 2.5 states in the case of both-edge detection. The MTU2 will not operate properly at narrower pulse widths.

In phase counting mode, the phase difference and overlap between the two input clocks must be at least 1.5 states, and the pulse width must be at least 2.5 states. Figure 12.108 shows the input clock conditions in phase counting mode.

Figure 12.108 Phase Difference, Overlap, and Pulse Width in Phase Counting Mode

12.7.3 Caution on Period Setting

When counter clearing on compare match is set, TCNT is cleared in the final state in which it matches the TGR value (the point at which the count value matched by TCNT is updated). Consequently, the actual counter frequency is given by the following formula:

• Channels 0 to 4

$$
f = \frac{P\phi}{(N+1)}
$$

• Channel 5

$$
f = \frac{P\phi}{N}
$$

Where f: Counter frequency

Pφ: MTU2 peripheral clock operating frequency

N: TGR set value

12.7.4 Contention between TCNT Write and Clear Operations

If the counter clear signal is generated in the T2 state of a TCNT write cycle, TCNT clearing takes precedence and the TCNT write is not performed.

Figure 12.109 shows the timing in this case.

Figure 12.109 Contention between TCNT Write and Clear Operations

12.7.5 Contention between TCNT Write and Increment Operations

If incrementing occurs in the T2 state of a TCNT write cycle, the TCNT write takes precedence and TCNT is not incremented.

Figure 12.110 shows the timing in this case.

Figure 12.110 Contention between TCNT Write and Increment Operations

12.7.6 Contention between TGR Write and Compare Match

If a compare match occurs in the T2 state of a TGR write cycle, the TGR write is executed and the compare match signal is also generated.

Figure 12.111 shows the timing in this case.

Figure 12.111 Contention between TGR Write and Compare Match

12.7.7 Contention between Buffer Register Write and Compare Match

If a compare match occurs in the T2 state of a TGR write cycle, the data that is transferred to TGR by the buffer operation is the data before write.

Figure 12.112 shows the timing in this case.

Figure 12.112 Contention between Buffer Register Write and Compare Match

12.7.8 Contention between Buffer Register Write and TCNT Clear

When the buffer transfer timing is set at the TCNT clear by the buffer transfer mode register (TBTM), if TCNT clear occurs in the T2 state of a TGR write cycle, the data that is transferred to TGR by the buffer operation is the data before write.

Figure 12.113 Contention between Buffer Register Write and TCNT Clear

12.7.9 Contention between TGR Read and Input Capture

If an input capture signal is generated in the T1 state of a TGR read cycle, the data that is read will be the data in the buffer before input capture transfer for channels 0 to 4, and the data after input capture transfer for channel 5.

Figures 12.114 and 12.115 show the timing in this case.

Figure 12.114 Contention between TGR Read and Input Capture (Channels 0 to 4)

Figure 12.115 Contention between TGR Read and Input Capture (Channel 5)

12.7.10 Contention between TGR Write and Input Capture

If an input capture signal is generated in the T2 state of a TGR write cycle, the input capture operation takes precedence and the write to TGR is not performed for channels 0 to 4. For channel 5, write to TGR is performed and the input capture signal is generated.

Figure 12.116 Contention between TGR Write and Input Capture (Channels 0 to 4)

Figure 12.117 Contention between TGR Write and Input Capture (Channel 5)

12.7.11 Contention between Buffer Register Write and Input Capture

If an input capture signal is generated in the T2 state of a buffer register write cycle, the buffer operation takes precedence and the write to the buffer register is not performed.

Figure 12.118 shows the timing in this case.

Figure 12.118 Contention between Buffer Register Write and Input Capture

12.7.12 TCNT_2 Write and Overflow/Underflow Contention in Cascade Connection

With timer counters TCNT 1 and TCNT 2 in a cascade connection, when a contention occurs during TCNT 1 count (during a TCNT 2 overflow/underflow) in the T2 state of the TCNT 2 write cycle, the write to TCNT 2 is conducted, and the TCNT 1 count signal is disabled. At this point, if there is match with TGRA_1 and the TCNT_1 value, a compare signal is issued. Furthermore, when the TCNT 1 count clock is selected as the input capture source of channel 0, TGRA_0 to TGRD_0 carry out the input capture operation. In addition, when the compare match/input capture is selected as the input capture source of TGRB_1, TGRB_1 carries out input capture operation. The timing is shown in figure 12.119.

For cascade connections, be sure to synchronize settings for channels 1 and 2 when setting TCNT clearing.

Figure 12.119 TCNT_2 Write and Overflow/Underflow Contention with Cascade Connection

12.7.13 Counter Value during Complementary PWM Mode Stop

When counting operation is suspended with $TCNT$ 3 and $TCNT$ 4 in complementary PWM mode, TCNT 3 has the timer dead time register (TDDR) value, and TCNT 4 is held at H'0000.

When restarting complementary PWM mode, counting begins automatically from the initialized state. This explanatory diagram is shown in figure 12.120.

When counting begins in another operating mode, be sure that TCNT 3 and TCNT 4 are set to the initial values.

Figure 12.120 Counter Value during Complementary PWM Mode Stop

12.7.14 Buffer Operation Setting in Complementary PWM Mode

In complementary PWM mode, conduct rewrites by buffer operation for the PWM cycle setting register (TGRA_3), timer cycle data register (TCDR), and duty setting registers (TGRB_3, TGRA_4, and TGRB_4).

In complementary PWM mode, channel 3 and channel 4 buffers operate in accordance with bit settings BFA and BFB of TMDR_3. When the BFA bit in TMDR_3 is set to 1, TGRC_3 functions as a buffer register for TGRA_3. At the same time, TGRC_4 functions as the buffer register for TGRA_4, and TCBR functions as the TCDR's buffer register.

12.7.15 Reset Sync PWM Mode Buffer Operation and Compare Match Flag

When setting buffer operation for reset sync PWM mode, set the BFA and BFB bits in TMDR 4 to 0. The TIOC4C pin will be unable to produce its waveform output if the BFA bit in TMDR_4 is set to 1

In reset sync PWM mode, the channel 3 and channel 4 buffers operate in accordance with the BFA and BFB bit settings of TMDR_3. For example, if the BFA bit in TMDR_3 is set to 1, TGRC_3 functions as the buffer register for TGRA_3. At the same time, TGRC_4 functions as the buffer register for TGRA_4.

The TGFC bit and TGFD bit in TSR_3 and TSR_4 are not set when TGRC_3 and TGRD_3 are operating as buffer registers.

Figure 12.121 shows an example of operations for TGR_3, TGR_4, TIOC3, and TIOC4, with TMDR_3's BFA and BFB bits set to 1, and TMDR_4's BFA and BFB bits set to 0.

Figure 12.121 Buffer Operation and Compare-Match Flags in Reset Synchronous PWM Mode

12.7.16 Overflow Flags in Reset Synchronous PWM Mode

When set to reset synchronous PWM mode, TCNT_3 and TCNT_4 start counting when the CST3 bit in TSTR is set to 1. At this point, TCNT_4's count clock source and count edge obey the TCR_3 setting.

In reset synchronous PWM mode, with cycle register TGRA_3's set value at H'FFFF, when specifying TGR3A compare-match for the counter clear source, TCNT_3 and TCNT_4 count up to H'FFFF, then a compare-match occurs with TGRA_3, and TCNT_3 and TCNT_4 are both cleared. At this point, TSR's overflow flag TCFV bit is not set.

Figure 12.122 shows a TCFV bit operation example in reset synchronous PWM mode with a set value for cycle register TGRA_3 of H'FFFF, when a TGRA_3 compare-match has been specified without synchronous setting for the counter clear source.

Figure 12.122 Reset Synchronous PWM Mode Overflow Flag

12.7.17 Contention between Overflow/Underflow and Counter Clearing

If overflow/underflow and counter clearing occur simultaneously, the TCFV/TCFU flag in TSR is not set and TCNT clearing takes precedence.

Figure 12.123 shows the operation timing when a TGR compare match is specified as the clearing source, and when H'FFFF is set in TGR.

Figure 12.123 Contention between Overflow and Counter Clearing

12.7.18 Contention between TCNT Write and Overflow/Underflow

If there is an up-count or down-count in the T2 state of a TCNT write cycle, and overflow/underflow occurs, the TCNT write takes precedence and the TCFV/TCFU flag in TSR is not set.

Figure 12.124 shows the operation timing when there is contention between TCNT write and overflow.

Figure 12.124 Contention between TCNT Write and Overflow

12.7.19 Cautions on Transition from Normal Operation or PWM Mode 1 to Reset-Synchronized PWM Mode

When making a transition from channel 3 or 4 normal operation or PWM mode 1 to resetsynchronized PWM mode, if the counter is halted with the output pins (TIOC3B, TIOC3D, TIOC4A, TIOC4C, TIOC4B, TIOC4D) in the high-level state, followed by the transition to resetsynchronized PWM mode and operation in that mode, the initial pin output will not be correct.

When making a transition from normal operation to reset-synchronized PWM mode, write H'11 to registers TIORH_3, TIORL_3, TIORH_4, and TIORL_4 to initialize the output pins to low level output, then set an initial register value of H'00 before making the mode transition.

When making a transition from PWM mode 1 to reset-synchronized PWM mode, first switch to normal operation, then initialize the output pins to low level output and set an initial register value of H'00 before making the transition to reset-synchronized PWM mode.

12.7.20 Output Level in Complementary PWM Mode and Reset-Synchronized PWM Mode

When channels 3 and 4 are in complementary PWM mode or reset-synchronized PWM mode, the PWM waveform output level is set with the OLSP and OLSN bits in the timer output control register (TOCR). In the case of complementary PWM mode or reset-synchronized PWM mode, TIOR should be set to H'00.

12.7.21 Interrupts in Module Standby Mode

If module standby mode is entered when an interrupt has been requested, it will not be possible to clear the CPU interrupt source or the DMAC activation source. Interrupts should therefore be disabled before entering module standby mode.

12.7.22 Simultaneous Capture of TCNT_1 and TCNT_2 in Cascade Connection

When timer counters 1 and 2 (TCNT 1 and TCNT 2) are operated as a 32-bit counter in cascade connection, the cascade counter value cannot be captured successfully even if input-capture input is simultaneously done to TIOC1A and TIOC2A or to TIOC1B and TIOC2B. This is because the input timing of TIOC1A and TIOC2A or of TIOC1B and TIOC2B may not be the same when external input-capture signals to be input into TCNT_1 and TCNT_2 are taken in synchronization with the internal clock. For example, TCNT 1 (the counter for upper 16 bits) does not capture the count-up value by overflow from TCNT_2 (the counter for lower 16 bits) but captures the count value before the count-up. In this case, the values of TCNT $1 = H'FFF1$ and TCNT $2 = H'0000$ should be transferred to TGRA_1 and TGRA_2 or to TGRB_1 and TGRB_2, but the values of TCNT_1 = H'FFF0 and TCNT_2 = H'0000 are erroneously transferred.

12.7.23 Notes on Output Waveform Control During Synchronous Counter Clearing in Complementary PWM Mode

In complementary PWM mode, when output waveform control during synchronous counter clearing is enabled (WRE in the TWCR register set to 1), the following problems may occur when condition (1) or condition (2), below, is satisfied.

- Dead time for the PWM output pins may be too short (or nonexistent).
- Active-level output from the PWM negative-phase pins may occur outside the correct activelevel output interval
- Condition (1): When synchronous clearing occurs in the PWM output dead time interval within initial output suppression interval (10) (figure 12.125).
- Condition (2): When synchronous clearing occurs within initial output suppression interval (10) or (11) and TGRB $3 \leq$ TDDR, TGRA $4 \leq$ TDDR, or TGRB $4 \leq$ TDDR is true (figure 12.126)

Figure 12.125 Condition (1) Synchronous Clearing Example

Figure 12.126 Condition (2) Synchronous Clearing Example

The following workaround can be used to avoid these problems.

When using synchronous clearing, make sure to set compare registers TGRB_3, TGRA_4, and TGRB_4 to a value twice or more the setting of dead time data register TDDR.

12.8 MTU2 Output Pin Initialization

12.8.1 Operating Modes

The MTU2 has the following six operating modes. Waveform output is possible in all of these modes.

- Normal mode (channels 0 to 4)
- PWM mode 1 (channels 0 to 4)
- PWM mode 2 (channels 0 to 2)
- Phase counting modes 1 to 4 (channels 1 and 2)
- Complementary PWM mode (channels 3 and 4)
- Reset-synchronized PWM mode (channels 3 and 4)

The MTU2 output pin initialization method for each of these modes is described in this section.

12.8.2 Reset Start Operation

The MTU2 output pins (TIOC*) are initialized low by a power-on reset. Since MTU2 pin function selection is performed by the pin function controller (PFC), when the PFC is set, the MTU2 pin states at that point are output to the ports. When MTU2 output is selected by the PFC immediately after a power-on reset, the MTU2 output initial level, low, is output directly at the port. When the active level is low, the system will operate at this point, and therefore the PFC setting should be made after initialization of the MTU2 output pins is completed.

Note: Channel number and port notation are substituted for $*$.

12.8.3 Operation in Case of Re-Setting Due to Error During Operation, etc.

If an error occurs during MTU2 operation, MTU2 output should be cut by the system. Cutoff is performed by switching the pin output to port output with the PFC and outputting the inverse of the active level. The pin initialization procedures for re-setting due to an error during operation, etc., and the procedures for restarting in a different mode after re-setting, are shown below.

The MTU2 has six operating modes, as stated above. There are thus 36 mode transition combinations, but some transitions are not available with certain channel and mode combinations. Possible mode transition combinations are shown in table 12.59.

Table 12.59 Mode Transition Combinations

[Legend]

Normal: Normal mode

PWM1: PWM mode 1

PWM2: PWM mode 2

PCM: Phase counting modes 1 to 4

CPWM: Complementary PWM mode

RPWM: Reset-synchronized PWM mode

12.8.4 Overview of Initialization Procedures and Mode Transitions in Case of Error during Operation, etc.

- When making a transition to a mode (Normal, PWM1, PWM2, PCM) in which the pin output level is selected by the timer I/O control register (TIOR) setting, initialize the pins by means of a TIOR setting.
- In PWM mode 1, since a waveform is not output to the TIOC*B (TIOC *D) pin, setting TIOR will not initialize the pins. If initialization is required, carry it out in normal mode, then switch to PWM mode 1.
- In PWM mode 2, since a waveform is not output to the cycle register pin, setting TIOR will not initialize the pins. If initialization is required, carry it out in normal mode, then switch to PWM mode 2.
- In normal mode or PWM mode 2, if TGRC and TGRD operate as buffer registers, setting TIOR will not initialize the buffer register pins. If initialization is required, clear buffer mode, carry out initialization, then set buffer mode again.
- In PWM mode 1, if either TGRC or TGRD operates as a buffer register, setting TIOR will not initialize the TGRC pin. To initialize the TGRC pin, clear buffer mode, carry out initialization, then set buffer mode again.
- When making a transition to a mode (CPWM, RPWM) in which the pin output level is selected by the timer output control register (TOCR) setting, switch to normal mode and perform initialization with TIOR, then restore TIOR to its initial value, and temporarily disable channel 3 and 4 output with the timer output master enable register (TOER). Then operate the unit in accordance with the mode setting procedure (TOCR setting, TMDR setting, TOER setting).

Note: Channel number is substituted for * indicated in this article.

Pin initialization procedures are described below for the numbered combinations in table 12.59. The active level is assumed to be low.

(1) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Normal Mode

 Figure 12.127 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in normal mode after re-setting.

Figure 12.127 Error Occurrence in Normal Mode, Recovery in Normal Mode

- 1. After a power-on reset, MTU2 output is low and ports are in the high-impedance state.
- 2. After a power-on reset, the TMDR setting is for normal mode.
- 3. For channels 3 and 4, enable output with TOER before initializing the pins with TIOR.
- 4. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence.)
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. Output goes low on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR.
- 11. Not necessary when restarting in normal mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(2) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.128 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in PWM mode 1 after re-setting.

Figure 12.128 Error Occurrence in Normal Mode, Recovery in PWM Mode 1

- 11. Set PWM mode 1.
- 12. Initialize the pins with TIOR. (In PWM mode 1, the $TIOC*B$ side is not initialized. If initialization is required, initialize in normal mode, then switch to PWM mode 1.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(3) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in PWM Mode 2

Figure 12.129 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in PWM mode 2 after re-setting.

- 11. Set PWM mode 2.
- 12. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized. If initialization is required, initialize in normal mode, then switch to PWM mode 2.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: PWM mode 2 can only be set for channels 0 to 2, and therefore TOER setting is not necessary.

(4) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.130 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in phase counting mode after re-setting.

Figure 12.130 Error Occurrence in Normal Mode, Recovery in Phase Counting Mode

- 11. Set phase counting mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: Phase counting mode can only be set for channels 1 and 2, and therefore TOER setting is not necessary.

(5) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.131 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in complementary PWM mode after re-setting.

- 11. Initialize the normal mode waveform generation section with TIOR.
- 12. Disable operation of the normal mode waveform generation section with TIOR.
- 13. Disable channel 3 and 4 output with TOER.
- 14. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 15. Set complementary PWM.
- 16. Enable channel 3 and 4 output with TOER.
- 17. Set MTU2 output with the PFC.
- 18. Operation is restarted by TSTR.

(6) Operation when Error Occurs during Normal Mode Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.132 shows an explanatory diagram of the case where an error occurs in normal mode and operation is restarted in reset-synchronized PWM mode after re-setting.

- 1 to 13 are the same as in figure 12.127.
- 14. Select the reset-synchronized PWM output level and cyclic output enabling/disabling with TOCR.
- 15. Set reset-synchronized PWM.
- 16. Enable channel 3 and 4 output with TOER.
- 17. Set MTU2 output with the PFC.
- 18. Operation is restarted by TSTR.

(7) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Normal Mode

Figure 12.133 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in normal mode after re-setting.

Figure 12.133 Error Occurrence in PWM Mode 1, Recovery in Normal Mode

- 1. After a power-on reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Set PWM mode 1.
- 3. For channels 3 and 4, enable output with TOER before initializing the pins with TIOR.
- 4. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence. In PWM mode 1, the TIOC*B side is not initialized.)
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. Output goes low on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR.
- 11. Set normal mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(8) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in PWM Mode 1

Figure 12.134 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in PWM mode 1 after re-setting.

- 11. Not necessary when restarting in PWM mode 1.
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC*B side is not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
(9) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in PWM Mode 2

Figure 12.135 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in PWM mode 2 after re-setting.

- 11. Set PWM mode 2.
- 12. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: PWM mode 2 can only be set for channels 0 to 2, and therefore TOER setting is not necessary.

(10) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.136 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in phase counting mode after re-setting.

- 11. Set phase counting mode.
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.
- Note: Phase counting mode can only be set for channels 1 and 2, and therefore TOER setting is not necessary.

(11) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.137 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in complementary PWM mode after re-setting.

- 11. Set normal mode for initialization of the normal mode waveform generation section.
- 12. Initialize the PWM mode 1 waveform generation section with TIOR.
- 13. Disable operation of the PWM mode 1 waveform generation section with TIOR.
- 14. Disable channel 3 and 4 output with TOER.
- 15. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 16. Set complementary PWM.
- 17. Enable channel 3 and 4 output with TOER.
- 18. Set MTU2 output with the PFC.
- 19. Operation is restarted by TSTR.

(12) Operation when Error Occurs during PWM Mode 1 Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.138 shows an explanatory diagram of the case where an error occurs in PWM mode 1 and operation is restarted in reset-synchronized PWM mode after re-setting.

Figure 12.138 Error Occurrence in PWM Mode 1, Recovery in Reset-Synchronized PWM Mode

- 15. Select the reset-synchronized PWM output level and cyclic output enabling/disabling with TOCR.
- 16. Set reset-synchronized PWM.
- 17. Enable channel 3 and 4 output with TOER.
- 18. Set MTU2 output with the PFC.
- 19. Operation is restarted by TSTR.

(13) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in Normal Mode

Figure 12.139 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in normal mode after re-setting.

Figure 12.139 Error Occurrence in PWM Mode 2, Recovery in Normal Mode

- 1. After a reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Set PWM mode 2.
- 3. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence. In PWM mode 2, the cycle register pins are not initialized. In the example, TIOC *A is the cycle register.)
- 4. Set MTU2 output with the PFC.
- 5. The count operation is started by TSTR.
- 6. Output goes low on compare-match occurrence.
- 7. An error occurs.
- 8. Set port output with the PFC and output the inverse of the active level.
- 9. The count operation is stopped by TSTR.
- 10. Set normal mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(14) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in PWM Mode 1

Figure 12.140 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in PWM mode 1 after re-setting.

Figure 12.140 Error Occurrence in PWM Mode 2, Recovery in PWM Mode 1

- 10. Set PWM mode 1.
- 11. Initialize the pins with TIOR. (In PWM mode 1, the TIOC*B side is not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(15) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in PWM Mode 2

Figure 12.141 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in PWM mode 2 after re-setting.

- 10. Not necessary when restarting in PWM mode 2.
- 11. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(16) Operation when Error Occurs during PWM Mode 2 Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.142 shows an explanatory diagram of the case where an error occurs in PWM mode 2 and operation is restarted in phase counting mode after re-setting.

Figure 12.142 Error Occurrence in PWM Mode 2, Recovery in Phase Counting Mode

- 10. Set phase counting mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(17) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.143 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in normal mode after re-setting.

Figure 12.143 Error Occurrence in Phase Counting Mode, Recovery in Normal Mode

- 1. After a power-on reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Set phase counting mode.
- 3. Initialize the pins with TIOR. (The example shows initial high output, with low output on compare-match occurrence.)
- 4. Set MTU2 output with the PFC.
- 5. The count operation is started by TSTR.
- 6. Output goes low on compare-match occurrence.
- 7. An error occurs.
- 8. Set port output with the PFC and output the inverse of the active level.
- 9. The count operation is stopped by TSTR.
- 10. Set in normal mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(18) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.144 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in PWM mode 1 after re-setting.

Figure 12.144 Error Occurrence in Phase Counting Mode, Recovery in PWM Mode 1

- 10. Set PWM mode 1.
- 11. Initialize the pins with TIOR. (In PWM mode 1, the TIOC *B side is not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(19) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in PWM Mode 2

Figure 12.145 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in PWM mode 2 after re-setting.

Figure 12.145 Error Occurrence in Phase Counting Mode, Recovery in PWM Mode 2

- 10. Set PWM mode 2.
- 11. Initialize the pins with TIOR. (In PWM mode 2, the cycle register pins are not initialized.)
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(20) Operation when Error Occurs during Phase Counting Mode Operation, and Operation is Restarted in Phase Counting Mode

Figure 12.146 shows an explanatory diagram of the case where an error occurs in phase counting mode and operation is restarted in phase counting mode after re-setting.

- 10. Not necessary when restarting in phase counting mode.
- 11. Initialize the pins with TIOR.
- 12. Set MTU2 output with the PFC.
- 13. Operation is restarted by TSTR.

(21) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.147 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in normal mode after re-setting.

- 1. After a power-on reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 3. Set complementary PWM.
- 4. Enable channel 3 and 4 output with TOER.
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. The complementary PWM waveform is output on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR. (MTU2 output becomes the complementary PWM output initial value.)
- 11. Set normal mode. (MTU2 output goes low.)
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(22) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.148 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in PWM mode 1 after re-setting.

- 11. Set PWM mode 1. (MTU2 output goes low.)
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC *B side is not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(23) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.149 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in complementary PWM mode after re-setting (when operation is restarted using the cycle and duty settings at the time the counter was stopped).

Figure 12.149 Error Occurrence in Complementary PWM Mode, Recovery in Complementary PWM Mode

- 11. Set MTU2 output with the PFC.
- 12. Operation is restarted by TSTR.
- 13. The complementary PWM waveform is output on compare-match occurrence.

(24) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.150 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in complementary PWM mode after re-setting (when operation is restarted using completely new cycle and duty settings).

- 11. Set normal mode and make new settings. (MTU2 output goes low.)
- 12. Disable channel 3 and 4 output with TOER.
- 13. Select the complementary PWM mode output level and cyclic output enabling/disabling with TOCR.
- 14. Set complementary PWM.
- 15. Enable channel 3 and 4 output with TOER.
- 16. Set MTU2 output with the PFC.
- 17. Operation is restarted by TSTR.

(25) Operation when Error Occurs during Complementary PWM Mode Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.151 shows an explanatory diagram of the case where an error occurs in complementary PWM mode and operation is restarted in reset-synchronized PWM mode.

- 11. Set normal mode. (MTU2 output goes low.)
- 12. Disable channel 3 and 4 output with TOER.
- 13. Select the reset-synchronized PWM mode output level and cyclic output enabling/disabling with TOCR.
- 14. Set reset-synchronized PWM.
- 15. Enable channel 3 and 4 output with TOER.
- 16. Set MTU2 output with the PFC.
- 17. Operation is restarted by TSTR.

SH7261 Group

(26) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in Normal Mode

Figure 12.152 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in normal mode after re-setting.

- 1. After a power-on reset, MTU2 output is low and ports are in the high-impedance state.
- 2. Select the reset-synchronized PWM output level and cyclic output enabling/disabling with TOCR.
- 3. Set reset-synchronized PWM.
- 4. Enable channel 3 and 4 output with TOER.
- 5. Set MTU2 output with the PFC.
- 6. The count operation is started by TSTR.
- 7. The reset-synchronized PWM waveform is output on compare-match occurrence.
- 8. An error occurs.
- 9. Set port output with the PFC and output the inverse of the active level.
- 10. The count operation is stopped by TSTR. (MTU2 output becomes the reset-synchronized PWM output initial value.)
- 11. Set normal mode. (MTU2 positive phase output is low, and negative phase output is high.)
- 12. Initialize the pins with TIOR.
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(27) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in PWM Mode 1

Figure 12.153 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in PWM mode 1 after re-setting.

- 11. Set PWM mode 1. (MTU2 positive phase output is low, and negative phase output is high.)
- 12. Initialize the pins with TIOR. (In PWM mode 1, the TIOC *B side is not initialized.)
- 13. Set MTU2 output with the PFC.
- 14. Operation is restarted by TSTR.

(28) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in Complementary PWM Mode

Figure 12.154 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in complementary PWM mode after resetting.

- 11. Disable channel 3 and 4 output with TOER.
- 12. Select the complementary PWM output level and cyclic output enabling/disabling with TOCR.
- 13. Set complementary PWM. (The MTU2 cyclic output pin goes low.)
- 14. Enable channel 3 and 4 output with TOER.
- 15. Set MTU2 output with the PFC.
- 16. Operation is restarted by TSTR.

(29) Operation when Error Occurs during Reset-Synchronized PWM Mode Operation, and Operation is Restarted in Reset-Synchronized PWM Mode

Figure 12.155 shows an explanatory diagram of the case where an error occurs in resetsynchronized PWM mode and operation is restarted in reset-synchronized PWM mode after resetting.

- 11. Set MTU2 output with the PFC.
- 12. Operation is restarted by TSTR.
- 13. The reset-synchronized PWM waveform is output on compare-match occurrence.

Section 13 8-Bit Timers (TMR)

This LSI has an on-chip 2-channel 8-bit timer based on an 8-bit counter. It can be used to count external events and, using compare-match signals with two registers, as a multifunction timer in a variety of applications, such as the generation of counter resets, interrupt requests, and pulse output with a user-defined duty cycle.

Figure 13.1 shows a block diagram of the 8-bit timer.

13.1 Features

• Selection of seven clock sources

The counters can be driven by one of six internal clock signals (Pφ/8, Pφ/64, Pφ/8192, Pφ/2, Pφ/32, or Pφ/1024) or an external clock input.

- Selection of three ways to clear the counters The counters can be cleared on compare match A or B, or by an external reset signal.
- Timer output control by a combination of two compare match signals

The timer output signal in each channel is controlled by a combination of two independent compare match signals, enabling the timer to output pulses with a desired duty cycle or PWM output.

• Cascading of two channels (TMR_0 and TMR_1)

Operation as a 16-bit timer is possible, using TMR_0 for the upper 8 bits and TMR_1 for the lower 8 bits (16-bit count mode).

TMR_1 can be used to count TMR_0 compare matches (compare match count mode).

• Three interrupt sources

Compare match A, compare match B, and overflow interrupts can be requested independently.

• Generation of trigger to start A/D converter conversion

Figure 13.1 Block Diagram of 8-Bit Timer

13.2 Input/Output Pins

Table 13.1 shows the pin configuration of the TMR.

Table 13.1 Pin Configuration

13.3 Register Descriptions

The TMR has the following registers.

Channel 0:

- Timer counter_0 (TCNT_0)
- Time constant register A_0 (TCORA_0)
- Time constant register B_0 (TCORB_0)
- Timer control register 0 (TCR 0)
- Timer counter control register 0 (TCCR 0)
- Timer control/status register_0 (TCSR_0)

Channel 1:

- Timer counter_1 (TCNT_1)
- Time constant register A_1 (TCORA_1)
- Time constant register B 1 (TCORB 1)
- Timer control register_1 (TCR_1)
- Timer counter control register 1 (TCCR 1)
- Timer control/status register_1 (TCSR_1)

13.3.1 Timer Counter (TCNT)

TCNT is an 8-bit readable/writable up-counter. TCNT_0 and TCNT_1 comprise a single 16-bit register so they can be accessed together by a word transfer instruction. Bits CKS2 to CKS0 in TCR and bits ICKS1 and ICKS0 in TCCR are used to select a clock. TCNT can be cleared by an external reset input signal, compare match A signal, or compare match B signal. Which signal is to be used for clearing is selected by bits CCLR1 and CCLR0 in TCR. When TCNT overflows from H'FF to H'00, bit OVF in TCSR is set to 1. TCNT is initialized to H'00.

13.3.2 Time Constant Register A (TCORA)

TCORA is an 8-bit readable/writable register. TCORA_0 and TCORA_1 comprise a single 16-bit register so they can be accessed together by a word transfer instruction. The value in TCORA is continually compared with the value in TCNT. When a match is detected, the corresponding CMFA flag in TCSR is set to 1. Note however that comparison is disabled during the T2 state of a TCORA write cycle. The timer output from the TMO pin can be freely controlled by this compare match signal (compare match A) and the settings of bits OS1 and OS0 in TCSR. TCORA is initialized to H'FF.

13.3.3 Time Constant Register B (TCORB)

TCORB is an 8-bit readable/writable register. TCORB_0 and TCORB_1 comprise a single 16-bit register so they can be accessed together by a word transfer instruction. TCORB is continually compared with the value in TCNT. When a match is detected, the corresponding CMFB flag in TCSR is set to 1. Note however that comparison is disabled during the T2 state of a TCORB write cycle. The timer output from the TMO pin can be freely controlled by this compare match signal (compare match B) and the settings of bits OS3 and OS2 in TCSR. TCORB is initialized to H'FF.

13.3.4 Timer Control Register (TCR)

TCR selects the TCNT clock source and the condition for clearing TCNT, and enables/disables interrupt requests.

Function Controller (PFC).

13.3.5 Timer Counter Control Register (TCCR)

TCCR selects the TCNT internal clock source and controls external reset input.

Table 13.2 Clock Input to TCNT and Count Condition

Notes: 1. If the clock input of TMR_0 is the TCNT_1 overflow signal and that of TMR_1 is the TCNT_0 compare match signal, no incrementing clock is generated. Do not use this setting.

 2. To use the external clock, the function of the corresponding pin should be selected using the pin function controller (PFC). For details, see section 25, Pin Function Controller (PFC).

13.3.6 Timer Control/Status Register (TCSR)

TCSR displays status flags, and controls compare match output.

• TCSR_0

• TCSR_0

Notes: 1. Only 0 can be written to bits 7 to 5, to clear these flags.

- 2. Timer output is disabled when bits OS3 to OS0 are all 0. Timer output is 0 until the first compare match occurs after resetting.
- TCSR_1

Notes: 1. Only 0 can be written to bits 7 to 5, to clear these flags.

 2. Timer output is disabled when bits OS3 to OS0 are all 0. Timer output is 0 until the first compare match occurs after resetting.

13.4 Operation

13.4.1 Pulse Output

Figure 13.2 shows an example of the 8-bit timer being used to generate a pulse output with a desired duty cycle. The control bits are set as follows:

- 1. In TCR, clear bit CCLR1 to 0 and set bit CCLR0 to 1 so that TCNT is cleared at a TCORA compare match.
- 2. In TCSR, set bits OS3 to OS0 to B'0110, causing the output to change to 1 at a TCORA compare match and to 0 at a TCORB compare match.

With these settings, the 8-bit timer provides pulses output at a cycle determined by TCORA with a pulse width determined by TCORB. No software intervention is required. The output level of the 8-bit timer holds 0 until the first compare match occurs after a reset.

Figure 13.2 Example of Pulse Output

13.4.2 Reset Input

Figure 13.3 shows an example of the 8-bit timer being used to generate a pulse which is output after a desired delay time from a TMRI input. The control bits are set as follows:

- 1. Set both bits CCLR1 and CCLR0 in TCR to 1 and set the TMRIS bit in TCCR to 1 so that TCNT is cleared at the high level input of the TMRI signal.
- 2. In TCSR, set bits OS3 to OS0 to B'0110, causing the output to change to 1 at a TCORA compare match and to 0 at a TCORB compare match.

With these settings, the 8-bit timer provides pulses output at a desired delay time from a TMRI input determined by TCORA and with a pulse width determined by TCORB and TCORA.

Figure 13.3 Example of Reset Input
13.5 Operation Timing

13.5.1 TCNT Count Timing

Figure 13.4 shows the TCNT count timing for internal clock input. Figure 13.5 shows the TCNT count timing for external clock input. Note that the external clock pulse width must be at least 1.5 states for incrementation at a single edge, and at least 2.5 states for incrementation at both edges. The counter will not increment correctly if the pulse width is less than these values.

Figure 13.4 Count Timing for Internal Clock Input at Falling Edge

Figure 13.5 Count Timing for External Clock Input at Falling and Rising Edges

13.5.2 Timing of CMFA and CMFB Setting at Compare Match

The CMFA and CMFB flags in TCSR are set to 1 by a compare match signal generated when the TCOR and TCNT values match. The compare match signal is generated at the last state in which the match is true, just before the timer counter is updated. Therefore, when the TCOR and TCNT values match, the compare match signal is not generated until the next TCNT clock input. Figure 13.6 shows this timing.

Figure 13.6 Timing of CMF Setting at Compare Match

13.5.3 Timing of Timer Output at Compare Match

When a compare match signal is generated, the timer output changes as specified by bits OS3 to OS0 in TCSR. Figure 13.7 shows the timing when the timer output is toggled by the compare match A signal.

Figure 13.7 Timing of Toggled Timer Output at Compare Match A

13.5.4 Timing of Counter Clear by Compare Match

TCNT is cleared when compare match A or B occurs, depending on the settings of bits CCLR1 and CCLR0 in TCR. Figure 13.8 shows the timing of this operation.

Figure 13.8 Timing of Counter Clear by Compare Match

13.5.5 Timing of TCNT External Reset

TCNT is cleared at the rising edge or high level of an external reset input, depending on the settings of bits CCLR1 and CCLR0 in TCR. The clear pulse width must be at least 2 states. Figures 13.9 and 13.10 show the timing of this operation.

Figure 13.9 Timing of Clearance by External Reset (Rising Edge)

Figure 13.10 Timing of Clearance by External Reset (High Level)

13.5.6 Timing of Overflow Flag (OVF) Setting

The OVF bit in TCSR is set to 1 when TCNT overflows (changes from H'FF to H'00). Figure 13.11 shows the timing of this operation.

Figure 13.11 Timing of OVF Setting

13.6 Operation with Cascaded Connection

If bits CKS2 to CKS0 in either TCR_0 or TCR_1 are set to B'100, the 8-bit timers of the two channels are cascaded. With this configuration, a single 16-bit timer could be used (16-bit counter mode) or compare matches of the 8-bit channel 0 could be counted by the timer of channel 1 (compare match count mode).

13.6.1 16-Bit Counter Mode

When bits CKS2 to CKS0 in TCR_0 are set to B'100, the timer functions as a single 16-bit timer with channel 0 occupying the upper 8 bits and channel 1 occupying the lower 8 bits.

(1) Setting of Compare Match Flags

- The CMF flag in TCSR 0 is set to 1 when a 16-bit compare match event occurs.
- The CMF flag in TCSR 1 is set to 1 when a lower 8-bit compare match event occurs.

(2) Counter Clear Specification

- If the CCLR1 and CCLR0 bits in TCR_{_0} have been set for counter clear at compare match, the 16-bit counter (TCNT_0 and TCNT_1 together) is cleared when a 16-bit compare match event occurs. The 16-bit counter (TCNT0 and TCNT1 together) is cleared even if counter clear by the TMRI0 pin has been set.
- The settings of the CCLR1 and CCLR0 bits in TCR 1 are ignored. The lower 8 bits cannot be cleared independently.

(3) Pin Output

- Control of output from the TMO0 pin by bits OS3 to OS0 in TCSR 0 is in accordance with the 16-bit compare match conditions.
- Control of output from the TMO1 pin by bits OS3 to OS0 in TCSR_1 is in accordance with the lower 8-bit compare match conditions.

13.6.2 Compare Match Count Mode

When bits CKS2 to CKS0 in TCR_1 are set to B'100, TCNT_1 counts compare match A for channel 0. Channels 0 and 1 are controlled independently. Conditions such as setting of the CMF flag, generation of interrupts, output from the TMO pin, and counter clear are in accordance with the settings for each channel.

13.7 Interrupt Sources

13.7.1 Interrupt Sources

There are three interrupt sources for the 8-bit timer (TMR_0 or TMR_1): CMIA, CMIB, and OVI. Their interrupt sources and priorities are shown in table 13.3. Each interrupt source is enabled or disabled by the corresponding interrupt enable bit in TCR or TCSR, and independent interrupt requests are sent for each to the interrupt controller.

13.7.2 A/D Converter Activation

The A/D converter can be activated only by TMR_0 compare match A.

If the ADTE bit in TCSR_0 is set to 1 when the CMFA flag in TCSR_0 is set to 1 by the occurrence of TMR_0 compare match A, a request to start A/D conversion is sent to the A/D converter. If the 8-bit timer conversion start trigger has been selected on the A/D converter side at this time, A/D conversion is started.

13.8 Usage Notes

13.8.1 Notes on Setting Cycle

If the compare match is selected for counter clear, TCNT is cleared at the last state in the cycle in which the values of TCNT and TCOR match. TCNT updates the counter value at this last state. Therefore, the counter frequency is obtained by the following formula.

 $f = P\phi/(N + 1)$ f: Counter frequency Pφ: Operating frequency N: TCOR value

13.8.2 Conflict between TCNT Write and Clear

If a counter clear signal is generated during the T_2 state of a TCNT write cycle, the clear takes priority and the write is not performed as shown in figure 13.12.

Figure 13.12 Conflict between TCNT Write and Clear

13.8.3 Conflict between TCNT Write and Increment

If a TCNT input clock pulse is generated during the T_2 state of a TCNT write cycle, the write takes priority and the counter is not incremented as shown in figure 13.13.

Figure 13.13 Conflict between TCNT Write and Increment

13.8.4 Conflict between TCOR Write and Compare Match

If a compare match event occurs during the T_2 state of a TCOR write cycle, the TCOR write takes priority and the compare match signal is inhibited as shown in figure 13.14.

Figure 13.14 Conflict between TCOR Write and Compare Match

13.8.5 Conflict between Compare Matches A and B

If compare match events A and B occur at the same time, the 8-bit timer operates in accordance with the priorities for the output statuses set for compare match A and compare match B, as shown in table 13.4.

Table 13.4 Timer Output Priorities

13.8.6 Switching of Internal Clocks and TCNT Operation

TCNT may be incremented erroneously depending on when the internal clock is switched. Table 13.5 shows the relationship between the timing at which the internal clock is switched (by writing to bits CKS1 and CKS0) and the TCNT operation.

When the TCNT clock is generated from an internal clock, the rising or falling edge of the internal clock pulse are always monitored. Table 13.5 assumes that the falling edge is selected. If the signal levels of the clocks before and after switching change from high to low as shown in item 3, the change is considered as the falling edge. Therefore, a TCNT clock pulse is generated and TCNT is incremented. This is similar to when the rising edge is selected.

The erroneous incrementation of TCNT can also happen when switching between rising and falling edges of the internal clock, and when switching between internal and external clocks.

Table 13.5 Switching of Internal Clock and TCNT Operation

Notes: 1. Includes switching from low to stop, and from stop to low.

- 2. Includes switching from stop to high.
- 3. Includes switching from high to stop.
- 4. Generated because the change of the signal levels is considered as a falling edge; TCNT is incremented.

13.8.7 Mode Setting with Cascaded Connection

If 16-bit counter mode and compare match count mode are specified at the same time, input clocks for TCNT 0 and TCNT 1 are not generated, and the counter stops. Do not specify 16-bit counter mode and compare match count mode simultaneously.

13.8.8 Module Standby Setting

Operation of the TMR can be disabled or enabled using the standby control register. The initial setting is for operation of the TMR to be halted. Register access is enabled by clearing module standby mode. For details, see section 27, Power-Down Modes.

13.8.9 Interrupts in Module Standby Mode

If module standby mode is entered when an interrupt has been requested, it will not be possible to clear the CPU interrupt source. Interrupts should therefore be disabled before entering module standby mode.

Section 14 Watchdog Timer (WDT)

This LSI includes the watchdog timer (WDT), which externally outputs an overflow signal (WDTOVF) on overflow of the counter when the value of the counter has not been updated because of a system malfunction. The WDT can simultaneously generate an internal reset signal for the entire LSI.

The WDT is a single channel timer that counts up the clock oscillation settling period when the system leaves software standby mode or the temporary standby periods that occur when the clock frequency is changed. It can also be used as a general watchdog timer or interval timer.

14.1 Features

- Can be used to ensure the clock oscillation settling time The WDT is used in leaving software standby mode or the temporary standby periods that occur when the clock frequency is changed.
- Can switch between watchdog timer mode and interval timer mode.
- Outputs WDTOVF signal in watchdog timer mode

When the counter overflows in watchdog timer mode, the WDTOVF signal is output externally. It is possible to select whether to reset the LSI internally when this happens. Either the power-on reset or manual reset signal can be selected as the internal reset type.

- Interrupt generation in interval timer mode An interval timer interrupt is generated when the counter overflows.
- Choice of eight counter input clocks

Eight clocks ($P\phi \times 1$ to $P\phi \times 1/16384$) that are obtained by dividing the peripheral clock can be selected.

Figure 14.1 shows a block diagram of the WDT.

Figure 14.1 Block Diagram of WDT

14.2 Input/Output Pin

Table 14.1 shows the pin configuration of the WDT.

Table 14.1 Pin Configuration

14.3 Register Descriptions

The WDT has the following registers.

Table 14.2 Register Configuration

Note: $*$ For the access size, see section 14.3.4, Notes on Register Access.

14.3.1 Watchdog Timer Counter (WTCNT)

WTCNT is an 8-bit readable/writable register that is incremented by cycles of the selected clock signal. When an overflow occurs, it generates a watchdog timer overflow signal (WDTOVF) in watchdog timer mode and an interrupt in interval timer mode. WTCNT is initialized to H'00 by a power-on reset caused by the RES pin or in deep standby mode or software standby mode.

Use word access to write to WTCNT, writing H'5A in the upper byte. Use byte access to read from WTCNT.

Note: The method for writing to WTCNT differs from that for other registers to prevent erroneous writes. See section 14.3.4, Notes on Register Access, for details.

14.3.2 Watchdog Timer Control/Status Register (WTCSR)

WTCSR is an 8-bit readable/writable register composed of bits to select the clock used for the count, overflow flags, and timer enable bit.

WTCSR is initialized to H'18 by a power-on reset caused by the \overline{RES} pin or in deep standby mode or software standby mode. When used to count the clock oscillation settling time for canceling software standby mode, it retains its value after counter overflow.

Use word access to write to WTCSR, writing H'A5 in the upper byte. Use byte access to read from WTCSR.

Note: The method for writing to WTCSR differs from that for other registers to prevent erroneous writes. See section 14.3.4, Notes on Register Access, for details.

14.3.3 Watchdog Reset Control/Status Register (WRCSR)

WRCSR is an 8-bit readable/writable register that controls output of the internal reset signal generated by watchdog timer counter (WTCNT) overflow.

WRCSR is initialized to H'1F by input of a reset signal from the \overline{RES} pin or in deep standby mode, but is not initialized by the internal reset signal generated by overflow of the WDT. WRCSR is initialized to H'1F in software standby mode.

Note: The method for writing to WRCSR differs from that for other registers to prevent erroneous writes. See section 14.3.4, Notes on Register Access, for details.

14.3.4 Notes on Register Access

The watchdog timer counter (WTCNT), watchdog timer control/status register (WTCSR), and watchdog reset control/status register (WRCSR) are more difficult to write to than other registers. The procedures for reading or writing to these registers are given below.

(1) Writing to WTCNT and WTCSR

These registers must be written by a word transfer instruction. They cannot be written by a byte or longword transfer instruction.

When writing to WTCNT, set the upper byte to H'5A and transfer the lower byte as the write data, as shown in figure 14.2. When writing to WTCSR, set the upper byte to H'A5 and transfer the lower byte as the write data. This transfer procedure writes the lower byte data to WTCNT or WTCSR.

Figure 14.2 Writing to WTCNT and WTCSR

(2) Writing to WRCSR

WRCSR must be written by a word access to address H'FFFE0004. It cannot be written by byte transfer or longword transfer instructions.

Procedures for writing 0 to WOVF (bit 7) and for writing to RSTE (bit 6) and RSTS (bit 5) are different, as shown in figure 14.3.

To write 0 to the WOVF bit, the write data must be H'A5 in the upper byte and H'00 in the lower byte. This clears the WOVF bit to 0. The RSTE and RSTS bits are not affected. To write to the RSTE and RSTS bits, the upper byte must be H'5A and the lower byte must be the write data. The values of bits 6 and 5 of the lower byte are transferred to the RSTE and RSTS bits, respectively. The WOVF bit is not affected.

(3) Reading from WTCNT, WTCSR, and WRCSR

WTCNT, WTCSR, and WRCSR are read in a method similar to other registers. WTCSR is allocated to address H'FFFE0000, WTCNT to address H'FFFE0002, and WRCSR to address H'FFFE0004. Byte transfer instructions must be used for reading from these registers.

14.4 WDT Usage

14.4.1 Canceling Software Standby Mode

The WDT can be used to cancel software standby mode with an interrupt such as an NMI interrupt. The procedure is described below. (The WDT does not operate when resets are used for canceling, so keep the RES or MRES pin low until clock oscillation settles.)

- 1. Before making a transition to software standby mode, always clear the TME bit in WTCSR to 0. When the TME bit is 1, an erroneous reset or interval timer interrupt may be generated when the count overflows.
- 2. Set the type of count clock used in the CKS[2:0] bits in WTCSR and the initial value of the counter in WTCNT. These values should ensure that the time till count overflow is longer than the clock oscillation settling time.
- 3. After setting the STBY bit to 1 and the DEEP bit to 0 in the standby control register (STBCR: see section 27, Power-Down Modes), the execution of a SLEEP instruction places the system in software standby mode and clock operation then stops.
- 4. The WDT starts counting by detecting the edge change of the NMI signal.
- 5. When the WDT count overflows, the CPG starts supplying the clock and this LSI resumes operation. The WOVF flag in WRCSR is not set when this happens.

14.4.2 Changing the Frequency

To change the frequency used by the PLL, use the WDT. When changing the frequency only by switching the divider, do not use the WDT.

- 1. Before changing the frequency, always clear the TME bit in WTCSR to 0. When the TME bit is 1, an erroneous reset or interval timer interrupt may be generated when the count overflows.
- 2. Set the type of count clock used in the CKS[2:0] bits in WTCSR and the initial value of the counter in WTCNT. These values should ensure that the time till count overflow is longer than the clock oscillation settling time. Note that, the WDT counts up by the clock to be set.
- 3. When the frequency control register (FRQCR) is written to, this LSI stops temporarily. The WDT starts counting.
- 4. When the WDT count overflows, the CPG resumes supplying the clock and this LSI resumes operation. The WOVF flag in WRCSR is not set when this happens.
- 5. The counter stops at the value of H'00.
- 6. Before changing WTCNT after execution of the frequency change instruction, always confirm that the value of WTCNT is H'00 by reading from WTCNT.

14.4.3 Using Watchdog Timer Mode

- 1. Set the WT/ \overline{IT} bit in WTCSR to 1, the type of count clock in the CKS[2:0] bits in WTCSR, whether this LSI is to be reset internally or not in the RSTE bit in WRCSR, the reset type if it is generated in the RSTS bit in WRCSR, and the initial value of the counter in WTCNT.
- 2. Set the TME bit in WTCSR to 1 to start the count in watchdog timer mode.
- 3. While operating in watchdog timer mode, rewrite the counter periodically to H'00 to prevent the counter from overflowing.
- 4. When the counter overflows, the WDT sets the WOVF flag in WRCSR to 1, and the WDTOVF signal is output externally (figure 14.4). The WDTOVF signal can be used to reset the system. The WDTOVF signal is output for $64 \times$ P ϕ clock cycles.
- 5. If the RSTE bit in WRCSR is set to 1, a signal to reset the inside of this LSI can be generated simultaneously with the WDTOVF signal. Either power-on reset or manual reset can be selected for this interrupt by the RSTS bit in WRCSR. The internal reset signal is output for $128 \times$ P ϕ clock cycles.
- 6. When a WDT overflow reset is generated simultaneously with a reset input on the RES pin, the RES pin reset takes priority, and the WOVF bit in WRCSR is cleared to 0.

Figure 14.4 Operation in Watchdog Timer Mode

14.4.4 Using Interval Timer Mode

When operating in interval timer mode, interval timer interrupts are generated at every overflow of the counter. This enables interrupts to be generated at set periods.

- 1. Clear the WT/ \overline{IT} bit in WTCSR to 0, set the type of count clock in the CKS[2:0] bits in WTCSR, and set the initial value of the counter in WTCNT.
- 2. Set the TME bit in WTCSR to 1 to start the count in interval timer mode.
- 3. When the counter overflows, the WDT sets the IOVF bit in WTCSR to 1 and an interval timer interrupt request is sent to the INTC. The counter then resumes counting.

Figure 14.5 Operation in Interval Timer Mode

14.5 Usage Notes

Pay attention to the following points when using the WDT in either the interval timer or watchdog timer mode.

14.5.1 Timer Variation

After timer operation has started, the period from the power-on reset point to the first count up timing of WTCNT varies depending on the time period that is set by the TME bit of WTCSR. The shortest such time period is thus one cycle of the peripheral clock, Pφ, while the longest is the result of frequency division according to the value in the CKS[2:0] bits. The timing of subsequent incrementation is in accord with the selected frequency division ratio. Accordingly, this time difference is referred to as timer variation.

This also applies to the timing of the first incrementation after WTCNT has been written to during timer operation.

14.5.2 Prohibition against Setting H'FF to WTCNT

When the value in WTCNT reaches H'FF, the WDT assumes that an overflow has occurred. Accordingly, when H'FF is set in WTCNT, an interval timer interrupt or WDT reset will occur immediately, regardless of the current clock selection by the CKS[2:0] bits.

14.5.3 Interval Timer Overflow Flag

When the value in WTCNT is H'FF, the IOVF flag in WTCSR cannot be cleared.

Only clear the IOVF flag when the value in WTCNT has either become H'00 or been changed to a value other than H'FF.

14.5.4 System Reset by WDTOVF **Signal**

If the WDTOVF signal is input to the RES pin of this LSI, this LSI cannot be initialized correctly.

Avoid input of the WDTOVF signal to the RES pin of this LSI through glue logic circuits. To reset the entire system with the WDTOVF signal, use the circuit shown in figure 14.6.

Figure 14.6 Example of System Reset Circuit Using WDTOVF **Signal**

14.5.5 Manual Reset in Watchdog Timer Mode

When a manual reset occurs in watchdog timer mode, the bus cycle is continued. If a manual reset occurs during DMAC burst transfer, manual reset exception handling will be pended until the CPU acquires the bus mastership.

Section 15 Realtime Clock (RTC)

This LSI has a realtime clock (RTC) with its own 32.768-kHz crystal oscillator.

15.1 Features

- Clock and calendar functions (BCD format): Seconds, minutes, hours, date, day of the week, month, and year
- 1-Hz to 64-Hz timer (binary format) 64-Hz counter indicates the state of the RTC divider circuit between 64 Hz and 1 Hz
- Start/stop function
- 30-second adjust function
- Alarm interrupt: Frame comparison of seconds, minutes, hours, date, day of the week, month, and year can be used as conditions for the alarm interrupt
- Periodic interrupts: the interrupt cycle may be $1/256$ second, $1/64$ second, $1/16$ second, $1/4$ second, 1/2 second, 1 second, or 2 seconds
- Carry interrupt: a carry interrupt indicates when a carry occurs during a counter read
- Automatic leap year adjustment

Figure 15.1 shows the block diagram of RTC.

Figure 15.1 RTC Block Diagram

15.2 Input/Output Pin

Table 15.1 shows the RTC pin configuration.

Table 15.1 Pin Configuration

15.3 Register Descriptions

The RTC has the following registers.

Table 15.2 Register Configuration

15.3.1 64-Hz Counter (R64CNT)

R64CNT indicates the state of the divider circuit between 64 Hz and 1 Hz.

Reading this register, when carry from 128-Hz divider stage is generated, sets the CF bit in the RTC control register 1 (RCR1) to 1 so that the carrying and reading 64 Hz counter are performed at the same time is indicated. In this case, the R64CNT should be read again after writing 0 to the CF bit in RCR1 since the read value is not valid.

After the RESET bit or ADJ bit in the RTC control register 2 (RCR2) is set to 1, the RTC divider circuit is initialized and R64CNT is initialized to H'00.

R64CNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

01234567 0 RRRRRRRR Bit: Initial value: R/W: — — — — — — — — 1Hz 2Hz 4Hz 8Hz 16Hz 32Hz 64Hz

15.3.2 Second Counter (RSECCNT)

RSECCNT is used for setting/counting in the BCD-coded second section. The count operation is performed by a carry for each second of the 64-Hz counter.

The assignable range is from 00 through 59 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RSECCNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.3 Minute Counter (RMINCNT)

RMINCNT is used for setting/counting in the BCD-coded minute section. The count operation is performed by a carry for each minute of the second counter.

The assignable range is from 00 through 59 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RMINCNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.4 Hour Counter (RHRCNT)

RHRCNT is used for setting/counting in the BCD-coded hour section. The count operation is performed by a carry for each 1 hour of the minute counter.

The assignable range is from 00 through 23 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RHRCNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.5 Day of Week Counter (RWKCNT)

RWKCNT is used for setting/counting day of week section. The count operation is performed by a carry for each day of the date counter.

The assignable range is from 0 through 6 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RWKCNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.6 Date Counter (RDAYCNT)

RDAYCNT is used for setting/counting in the BCD-coded date section. The count operation is performed by a carry for each day of the hour counter.

The assignable range is from 01 through 31 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RDAYCNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

The range of date changes with each month and in leap years. Please confirm the correct setting. Leap years are recognized by dividing the year counter values by 400, 100, and 4 and obtaining a fractional result of 0. The year counter value of 0000 is included in the leap year.

15.3.7 Month Counter (RMONCNT)

RMONCNT is used for setting/counting in the BCD-coded month section. The count operation is performed by a carry for each month of the date counter.

The assignable range is from 01 through 12 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RMONCNT is not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.8 Year Counter (RYRCNT)

RYRCNT is used for setting/counting in the BCD-coded year section. The count operation is performed by a carry for each year of the month counter.

The assignable range is from 0000 through 9999 (practically in BCD), otherwise operation errors occur. Carry out write processing after stopping the count operation through the setting of the START bit in RCR2.

RYRCNT is not initialized by a power-on reset or manual reset, in deep standby mode or software standby mode.

15.3.9 Second Alarm Register (RSECAR)

RSECAR is an alarm register corresponding to the BCD coded second counter RSECCNT of the RTC. When the ENB bit is set to 1, a comparison with the RSECCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 00 through $59 + ENB$ bits (practically in BCD), otherwise operation errors occur.

The ENB bit in RSECAR is initialized to 0 by a power-on reset or in deep standby mode. The other bits are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.10 Minute Alarm Register (RMINAR)

RMINAR is an alarm register corresponding to the minute counter RMINCNT. When the ENB bit is set to 1, a comparison with the RMINCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 00 through 59 + ENB bits (practically in BCD), otherwise operation errors occur.

The ENB bit in RMINAR is initialized by a power-on reset or in deep standby mode. The other bits are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.11 Hour Alarm Register (RHRAR)

RHRAR is an alarm register corresponding to the BCD coded hour counter RHRCNT of the RTC. When the ENB bit is set to 1, a comparison with the RHRCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 00 through $23 + ENB$ bits (practically in BCD), otherwise operation errors occur.

The ENB bit in RHRAR is initialized by a power-on reset or in deep standby mode. The other bits are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.12 Day of Week Alarm Register (RWKAR)

RWKAR is an alarm register corresponding to the BCD coded day of week counter RWKCNT. When the ENB bit is set to 1, a comparison with the RWKCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 0 through $6 +$ ENB bits (practically in BCD), otherwise operation errors occur.

The ENB bit in RWKAR is initialized by a power-on reset or in deep standby mode. The other bits are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.13 Date Alarm Register (RDAYAR)

RDAYAR is an alarm register corresponding to the BCD coded date counter RDAYCNT. When the ENB bit is set to 1, a comparison with the RDAYCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 01 through 31 + ENB bits (practically in BCD), otherwise operation errors occur.

The ENB bit in RDAYAR is initialized by a power-on reset or in deep standby mode. The other bits are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.14 Month Alarm Register (RMONAR)

RMONAR is an alarm register corresponding to the BCD coded month counter RMONCNT. When the ENB bit is set to 1, a comparison with the RMONCNT value is performed. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The assignable range is from 01 through 12 + ENB bits (practically in BCD), otherwise operation errors occur.

The ENB bit in RMONAR is initialized by a power-on reset or in deep standby mode. The other bits are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.3.15 Year Alarm Register (RYRAR)

RYRAR is an alarm register corresponding to the year counter RYRCNT. The assignable range is from 0000 through 9999 (practically in BCD), otherwise operation errors occur. RYRAR is not initialized by a power-on reset, a manual reset, or in deep standby mode or software standby mode.

15.3.16 RTC Control Register 1 (RCR1)

RCR1 is a register that affects carry flags and alarm flags. It also selects whether to generate interrupts for each flag.

RCR1 is initialized to H'00 by a power-on reset, a manual reset, or in deep standby mode. The CF flag is retained undefined until the division circuit is reset (the RESET and ADJ bits in RCR2 are set to 1). When using the CF flag, make sure to reset the divider circuit beforehand. This register is not initialized in software standby mode.

15.3.17 RTC Control Register 2 (RCR2)

RCR2 is a register for periodic interrupt control, 30-second adjustment ADJ, divider circuit RESET, and RTC count control.

RCR2 is initialized to H'09 by a power-on reset or in deep standby mode. Bits other than the RTCEN and START bits are initialized by a manual reset. It is not initialized in software standby mode, and retains its contents.

15.3.18 RTC Control Register 3 (RCR3)

When the ENB bit in RCR3 is set to 1, RCR3 compares the value of RYRCNT and that of RYRAR. From among RSECAR/RMINAR/RHRAR/RWKAR/RDAYAR/RMONAR/RCR3, the counter and alarm register comparison is performed only on those with ENB bits set to 1, and if each of those coincides, an alarm flag of RCR1 is set to 1.

The ENB bit in RCR3 is initialized by a power-on reset or in deep standby mode. Remaining fields of RCR3 are not initialized by a power-on reset or manual reset, or in deep standby and software standby modes.

15.4 Operation

RTC usage is shown below.

15.4.1 Initial Settings of Registers after Power-On

All the registers should be set after the power is turned on.

15.4.2 Setting Time

Figure 15.2 shows how to set the time when the clock is stopped.

Figure 15.2 Setting Time

15.4.3 Reading Time

Figure 15.3 shows how to read the time.

Figure 15.3 Reading Time

If a carry occurs while reading the time, the correct time will not be obtained, so it must be read again. Part (a) in figure 15.3 shows the method of reading the time without using interrupts; part (b) in figure 15.3 shows the method using carry interrupts. To keep programming simple, method (a) should normally be used.

15.4.4 Alarm Function

Figure 15.4 shows how to use the alarm function.

Figure 15.4 Using Alarm Function

Alarms can be generated using seconds, minutes, hours, day of the week, date, month, year, or any combination of these. Set the ENB bit in the register on which the alarm is placed to 1, and then set the alarm time in the lower bits. Clear the ENB bit in the register on which the alarm is not placed to 0.

When the clock and alarm times match, 1 is set in the AF bit in RCR1. Alarm detection can be checked by reading this bit, but normally it is done by interrupt. If 1 is set in the AIE bit in RCR1, an interrupt is generated when an alarm occurs.

The alarm flag is set when the clock and alarm times match. However, the alarm flag can be cleared by writing 0.

15.5 Usage Notes

15.5.1 Register Writing during RTC Count

Do not write to the count registers (RSECCNT, RMINCNT, RHRCNT, RDAYCNT, RWKCNT, RMONCNT, and RYRCNT) during the RTC counting (while the START bit in RCR2 is 1). If any of the count registers is written to during the RTC counting, the count register may not be read correctly immediately after the execution of a write instruction. The RTC counting must be stopped before writing to any of the count registers.

15.5.2 Use of Realtime Clock (RTC) Periodic Interrupts

The method of using the periodic interrupt function is shown in figure 15.5.

A periodic interrupt can be generated periodically at the interval set by the flags PES0 to PES2 in RCR2. When the time set by the PES0 to PES2 has elapsed, the PEF is set to 1.

The PEF is cleared to 0 upon periodic interrupt generation or when the flags PES0 to PES2 are set. Periodic interrupt generation can be confirmed by reading this bit, but normally the interrupt function is used.

Figure 15.5 Using Periodic Interrupt Function

15.5.3 Transition to Standby Mode after Setting Register

When a transition to standby mode is made after registers in the RTC are set, sometimes counting is not performed correctly. In case the registers are set, be sure to make a transition to standby mode after waiting for two count clocks or more.

15.5.4 Crystal Oscillator Circuit for RTC

Crystal oscillator circuit constants (recommended values) for the RTC are shown in table 15.3, and the RTC crystal oscillator circuit in figure 15.6.

Figure 15.6 Example of Connecting Crystal Oscillator Circuit for RTC

15.5.5 Procedure for Setting the 30-Second Adjustment Function

Figure 15.7 shows the procedure for setting the 30-second adjustment function.

Figure 15.7 Procedure for Setting the 30-Second Adjustment Function

To use the 30-second adjustment function, the minutes, hours, date, day of the week, month, and year counters need to be written to. Thus, after clearing the START bit in RCR2 and reading out the minutes, hours, date, day of the week, month, and year counters and then writing the read values back, set the ADJ bit in RCR2 to 1. After the 30-second adjustment, set the START bit in RCR2 to 1 to start the clock operation.

Section 16 Serial Communication Interface with FIFO (SCIF)

This LSI has an eight-channel serial communication interface with FIFO (SCIF) that supports both asynchronous and clocked synchronous serial communication. It also has 16-stage FIFO registers for both transmission and reception independently for each channel that enable this LSI to perform efficient high-speed continuous communication.

16.1 Features

- Asynchronous serial communication:
	- ⎯ Serial data communication is performed by start-stop in character units. The SCIF can communicate with a universal asynchronous receiver/transmitter (UART), an asynchronous communication interface adapter (ACIA), or any other communications chip that employs a standard asynchronous serial system. There are eight selectable serial data communication formats.
	- Data length: 7 or 8 bits
	- Stop bit length: 1 or 2 bits
	- Parity: Even, odd, or none
	- Receive error detection: Parity, framing, and overrun errors
	- ⎯ Break detection: Break is detected when a framing error is followed by at least one frame at the space 0 level (low level). It is also detected by reading the RxD level directly from the serial port register when a framing error occurs.
- Clocked synchronous serial communication:
	- ⎯ Serial data communication is synchronized with a clock signal. The SCIF can communicate with other chips having a clocked synchronous communication function. There is one serial data communication format.
	- Data length: 8 bits
	- ⎯ Receive error detection: Overrun errors
- Full duplex communication: The transmitting and receiving sections are independent, so the SCIF can transmit and receive simultaneously. Both sections use 16-stage FIFO buffering, so high-speed continuous data transfer is possible in both the transmit and receive directions.
- On-chip baud rate generator with selectable bit rates
- Internal or external transmit/receive clock source: From either baud rate generator (internal) or SCK pin (external)
- Four types of interrupts: Transmit-FIFO-data-empty, break, receive-FIFO-data-full, and receive-error interrupts are requested independently.
- When the SCIF is not in use, it can be stopped by halting the clock supplied to it, saving power.
- The quantity of data in the transmit and receive FIFO registers and the number of receive errors of the receive data in the receive FIFO register can be ascertained.
- A time-out error (DR) can be detected when receiving in asynchronous mode.

Figure 16.1 shows a block diagram of the SCIF.

Figure 16.1 Block Diagram of SCIF

16.2 Input/Output Pins

Table 16.1 shows the pin configuration of the SCIF.

Table 16.1 Pin Configuration

16.3 Register Descriptions

The SCIF has the following registers.

Table 16.2 Register Configuration

Notes: 1. Only 0 can be written to clear the flag. Bits 15 to 8, 3, and 2 are read-only bits that cannot be modified.

 2. Only 0 can be written to clear the flag. Bits 15 to 1 are read-only bits that cannot be modified.

16.3.1 Receive Shift Register (SCRSR)

SCRSR receives serial data. Data input at the RxD pin is loaded into SCRSR in the order received, LSB (bit 0) first, converting the data to parallel form. When one byte has been received, it is automatically transferred to the receive FIFO data register (SCFRDR).

The CPU cannot read or write to SCRSR directly.

16.3.2 Receive FIFO Data Register (SCFRDR)

SCFRDR is a 16-byte FIFO register that stores serial receive data. The SCIF completes the reception of one byte of serial data by moving the received data from the receive shift register (SCRSR) into SCFRDR for storage. Continuous reception is possible until 16 bytes are stored. The CPU can read but not write to SCFRDR. If data is read when there is no receive data in the SCFRDR, the value is undefined.

When SCFRDR is full of receive data, subsequent serial data is lost.

SCFRDR is initialized to an undefined value by a power-on reset or in deep standby mode.

16.3.3 Transmit Shift Register (SCTSR)

SCTSR transmits serial data. The SCIF loads transmit data from the transmit FIFO data register (SCFTDR) into SCTSR, then transmits the data serially from the TxD pin, LSB (bit 0) first. After transmitting one data byte, the SCIF automatically loads the next transmit data from SCFTDR into SCTSR and starts transmitting again.

The CPU cannot read or write to SCTSR directly.

16.3.4 Transmit FIFO Data Register (SCFTDR)

SCFTDR is a 16-byte FIFO register that stores data for serial transmission. When the SCIF detects that the transmit shift register (SCTSR) is empty, it moves transmit data written in the SCFTDR into SCTSR and starts serial transmission. Continuous serial transmission is performed until there is no transmit data left in SCFTDR. The CPU can write to SCFTDR at all times.

When SCFTDR is full of transmit data (16 bytes), no more data can be written. If writing of new data is attempted, the data is ignored.

SCFTDR is initialized to an undefined value by a power-on reset or in deep standby mode.

16.3.5 Serial Mode Register (SCSMR)

SCSMR specifies the SCIF serial communication format and selects the clock source for the baud rate generator.

The CPU can always read and write to SCSMR. SCSMR is initialized to H'0000 by a power-on reset or in deep standby mode.

16.3.6 Serial Control Register (SCSCR)

SCSCR operates the SCIF transmitter/receiver, enables/disables interrupt requests, and selects the transmit/receive clock source. The CPU can always read and write to SCSCR. SCSCR is initialized to H'0000 by a power-on reset or in deep standby mode.

16.3.7 Serial Status Register (SCFSR)

SCFSR is a 16-bit register. The upper 8 bits indicate the number of receives errors in the receive FIFO data register, and the lower 8 bits indicate the status flag indicating SCIF operating state.

The CPU can always read and write to SCFSR, but cannot write 1 to the status flags (ER, TEND, TDFE, BRK, RDF, and DR). These flags can be cleared to 0 only if they have first been read (after being set to 1). Bits 3 (FER) and 2 (PER) are read-only bits that cannot be written. SCFSR is initialized by a power-on reset or in deep standby mode.

Note: $*$ Only 0 can be written to clear the flag after 1 is read.

Note: * Only 0 can be written to clear the flag after 1 is read.

16.3.8 Bit Rate Register (SCBRR)

SCBRR is an 8-bit register that, together with the baud rate generator clock source selected by the CKS1 and CKS0 bits in the serial mode register (SCSMR), determines the serial transmit/receive bit rate.

The CPU can always read and write to SCBRR. SCBRR is initialized to H'FF by a power-on reset or in deep standby mode. Each channel has independent baud rate generator control, so different values can be set in eight channels.

The SCBRR setting is calculated as follows:

• Asynchronous mode:

$$
N = \frac{P\phi}{64 \times 2^{2n-1} \times B} \times 10^6 - 1
$$

• Clocked synchronous mode:

$$
N = \frac{P\phi}{8 \times 2^{2n-1} \times B} \times 10^6 - 1
$$

- B: Bit rate (bits/s)
- N: SCBRR setting for baud rate generator $(0 \le N \le 255)$ (The setting must satisfy the electrical characteristics.)
- Pφ: Operating frequency for peripheral modules (MHz)
- n: Baud rate generator clock source $(n = 0, 1, 2, 3)$ (for the clock sources and values of n, see table 16.3.)

Table 16.3 SCSMR Settings

The bit rate error in asynchronous is given by the following formula:

$$
\text{Error } (\%) = \left\{\!\!\left.\frac{\mathsf{P} \varphi \, \times 10^6}{(N+1) \times \mathsf{B} \times 64 \times \ 2^{2n\text{-}1}} - 1\!\!\right\} \, \times 100\!\!\right.
$$

Table 16.4 lists examples of SCBRR settings in asynchronous mode, and table 16.5 lists examples of SCBRR settings in clocked synchronous mode.

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode) (1)

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode) (2)

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode) (3)

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode) (4)

Table 16.4 Bit Rates and SCBRR Settings (Asynchronous Mode) (5)

Note: Settings with an error of 1% or less are recommended.

Table 16.5 Bit Rates and SCBRR Settings (Clocked Synchronous Mode) (1)

Table 16.5 Bit Rates and SCBRR Settings (Clocked Synchronous Mode) (2)

[Legend]

Blank: No setting possible, or it is not possible to satisfy the electrical characteristics of the MCU regardless of the communication partner device.

—: Setting possible, but error occurs

Table 16.6 indicates the maximum bit rates in asynchronous mode when the baud rate generator is used. Tables 16.7 and 16.8 list the maximum rates when the external clock input is used (when tscyc = 12 tpcyc^{*}).

Note: * Make sure that the electrical characteristics of this MCU and that of a connected MCU are satisfied.

Table 16.6 Maximum Bit Rates for Various Frequencies with Baud Rate Generator (Asynchronous Mode)

Table 16.7 Maximum Bit Rates with External Clock Input (Asynchronous Mode)

Table 16.8 Maximum Bit Rates with External Clock Input (Clocked Synchronous Mode, $t_{\text{scyc}} = 12 t_{\text{pcyc}}$)

16.3.9 FIFO Control Register (SCFCR)

SCFCR resets the quantity of data in the transmit and receive data FIFO registers, sets the trigger data quantity, and contains an enable bit for loop-back testing. SCFCR can always be read and written to by the CPU. It is initialized to H'0000 by a power-on reset or in deep standby mode.

16.3.10 FIFO Data Count Register (SCFDR)

SCFDR is a 16-bit register which indicates the quantity of data stored in the transmit FIFO data register (SCFTDR) and the receive FIFO data register (SCFRDR).

It indicates the quantity of transmit data in SCFTDR with the upper 8 bits, and the quantity of receive data in SCFRDR with the lower 8 bits. SCFDR can always be read by the CPU. SCFDR is initialized to H'0000 by a power on reset or in deep standby mode.

16.3.11 Serial Port Register (SCSPTR)

SCSPTR controls input/output and data of pins multiplexed to SCIF function. Bits 3 and 2 can control input/output data of SCK pin. Bits 1 and 0 can input data from RxD pin and output data to TxD pin, so they control break of serial transmitting/receiving.

The CPU can always read and write to SCSPTR. SCSPTR is initialized to H'0050 by a power-on reset or in deep standby mode.

16.3.12 Line Status Register (SCLSR)

The CPU can always read or write to SCLSR, but cannot write 1 to the ORER flag. This flag can be cleared to 0 only if it has first been read (after being set to 1).

SCLSR is initialized to H'0000 by a power-on reset or in deep standby mode.

Note: $*$ Only 0 can be written to clear the flag after 1 is read.

16.4 Operation

16.4.1 Overview

For serial communication, the SCIF has an asynchronous mode in which characters are synchronized individually, and a clocked synchronous mode in which communication is synchronized with clock pulses.

The SCIF has a 16-stage FIFO buffer for both transmission and receptions, reducing the overhead of the CPU, and enabling continuous high-speed communication. The transmission format is selected in the serial mode register (SCSMR), as shown in table 16.9. The SCIF clock source is selected by the combination of the CKE1 and CKE0 bits in the serial control register (SCSCR), as shown in table 16.10.

(1) Asynchronous Mode

- Data length is selectable: 7 or 8 bits
- Parity bit is selectable. So is the stop bit length (1 or 2 bits). The combination of the preceding selections constitutes the communication format and character length.
- In receiving, it is possible to detect framing errors, parity errors, receive FIFO data full, overrun errors, receive data ready, and breaks.
- The number of stored data bytes is indicated for both the transmit and receive FIFO registers.
- An internal or external clock can be selected as the SCIF clock source.
	- ⎯ When an internal clock is selected, the SCIF operates using the on-chip baud rate generator.
	- When an external clock is selected, the external clock input must have a frequency 16 times the bit rate. (The on-chip baud rate generator is not used.)

(2) Clocked Synchronous Mode

- The transmission/reception format has a fixed 8-bit data length.
- In receiving, it is possible to detect overrun errors (ORER).
- An internal or external clock can be selected as the SCIF clock source.
	- When an internal clock is selected, the SCIF operates using the on-chip baud rate generator, and outputs a serial clock signal to external devices.
	- ⎯ When an external clock is selected, the SCIF operates on the input serial clock. The onchip baud rate generator is not used.

Table 16.9 SCSMR Settings and SCIF Communication Formats

[Legend]

x: Don't care

Table 16.10 SCSMR and SCSCR Settings and SCIF Clock Source Selection

[Legend]

x: Don't care

16.4.2 Operation in Asynchronous Mode

In asynchronous mode, each transmitted or received character begins with a start bit and ends with a stop bit. Serial communication is synchronized one character at a time.

The transmitting and receiving sections of the SCIF are independent, so full duplex communication is possible. The transmitter and receiver are 16-byte FIFO buffered, so data can be written and read while transmitting and receiving are in progress, enabling continuous transmitting and receiving.

Figure 16.2 shows the general format of asynchronous serial communication.

In asynchronous serial communication, the communication line is normally held in the mark (high) state. The SCIF monitors the line and starts serial communication when the line goes to the space (low) state, indicating a start bit. One serial character consists of a start bit (low), data (LSB first), parity bit (high or low), and stop bit (high), in that order.

When receiving in asynchronous mode, the SCIF synchronizes at the falling edge of the start bit. The SCIF samples each data bit on the eighth pulse of a clock with a frequency 16 times the bit rate. Receive data is latched at the center of each bit.

Figure 16.2 Example of Data Format in Asynchronous Communication (8-Bit Data with Parity and Two Stop Bits)

(1) Transmit/Receive Formats

Table 16.11 lists the eight communication formats that can be selected in asynchronous mode. The format is selected by settings in the serial mode register (SCSMR).

[Legend]

START: Start bit

STOP: Stop bit

P: Parity bit

(2) Clock

An internal clock generated by the on-chip baud rate generator or an external clock input from the SCK pin can be selected as the SCIF transmit/receive clock. The clock source is selected by the C/A bit in the serial mode register (SCSMR) and bits CKE1 and CKE0 in the serial control register (SCSCR). For clock source selection, refer to table 16.10.

When an external clock is input at the SCK pin, it must have a frequency equal to 16 times the desired bit rate.

When the SCIF operates on an internal clock, it can output a clock signal on the SCK pin. The frequency of this output clock is 16 times the desired bit rate.

(3) Transmitting and Receiving Data

• **SCIF Initialization (Asynchronous Mode)**

Before transmitting or receiving, clear the TE and RE bits to 0 in the serial control register (SCSCR), then initialize the SCIF as follows.

When changing the operation mode or the communication format, always clear the TE and RE bits to 0 before following the procedure given below. Clearing TE to 0 initializes the transmit shift register (SCTSR). Clearing TE and RE to 0, however, does not initialize the serial status register (SCFSR), transmit FIFO data register (SCFTDR), or receive FIFO data register (SCFRDR), which retain their previous contents. Clear TE to 0 after all transmit data has been transmitted and the TEND flag in the SCFSR is set. The TE bit can be cleared to 0 during transmission, but the transmit data goes to the Mark state after the bit is cleared to 0. Set the TFRST bit in SCFCR to 1 and reset SCFTDR before TE is set again to start transmission.

When an external clock is used, the clock should not be stopped during initialization or subsequent operation. SCIF operation becomes unreliable if the clock is stopped.

Figure 16.3 shows a sample flowchart for initializing the SCIF.

Figure 16.3 Sample Flowchart for SCIF Initialization

• **Transmitting Serial Data (Asynchronous Mode)**

Figure 16.4 shows a sample flowchart for serial transmission.

Use the following procedure for serial data transmission after enabling the SCIF for transmission.

In serial transmission, the SCIF operates as described below.

- 1. When data is written into the transmit FIFO data register (SCFTDR), the SCIF transfers the data from SCFTDR to the transmit shift register (SCTSR) and starts transmitting. Confirm that the TDFE flag in the serial status register (SCFSR) is set to 1 before writing transmit data to SCFTDR. The number of data bytes that can be written is (16 – transmit trigger setting).
- 2. When data is transferred from SCFTDR to SCTSR and transmission is started, consecutive transmit operations are performed until there is no transmit data left in SCFTDR. When the number of transmit data bytes in SCFTDR falls below the transmit trigger number set in the FIFO control register (SCFCR), the TDFE flag is set. If the TIE bit in the serial control register (SCSR) is set to 1 at this time, a transmit-FIFO-data-empty interrupt (TXI) request is generated.

The serial transmit data is sent from the TxD pin in the following order.

- A. Start bit: One-bit 0 is output.
- B. Transmit data: 8-bit or 7-bit data is output in LSB-first order.
- C. Parity bit: One parity bit (even or odd parity) is output. (A format in which a parity bit is not output can also be selected.)
- D. Stop bit(s): One or two 1 bits (stop bits) are output.
- E. Mark state: 1 is output continuously until the start bit that starts the next transmission is sent.
- 3. The SCIF checks the SCFTDR transmit data at the timing for sending the stop bit. If data is present, the data is transferred from SCFTDR to SCTSR, the stop bit is sent, and then serial transmission of the next frame is started.

Figure 16.5 shows an example of the operation for transmission.

Figure 16.5 Example of Transmit Operation (8-Bit Data, Parity, 1 Stop Bit)

• **Receiving Serial Data (Asynchronous Mode)**

Figures 16.6 and 16.7 show sample flowcharts for serial reception.

Use the following procedure for serial data reception after enabling the SCIF for reception.

Figure 16.6 Sample Flowchart for Receiving Serial Data

Figure 16.7 Sample Flowchart for Receiving Serial Data (cont)

In serial reception, the SCIF operates as described below.

- 1. The SCIF monitors the transmission line, and if a 0 start bit is detected, performs internal synchronization and starts reception.
- 2. The received data is stored in SCRSR in LSB-to-MSB order.
- 3. The parity bit and stop bit are received. After receiving these bits, the SCIF carries out the following checks.
	- A. Stop bit check: The SCIF checks whether the stop bit is 1. If there are two stop bits, only the first is checked.
	- B. The SCIF checks whether receive data can be transferred from the receive shift register (SCRSR) to SCFRDR.
	- C. Overrun check: The SCIF checks that the ORER flag is 0, indicating that the overrun error has not occurred.
	- D. Break check: The SCIF checks that the BRK flag is 0, indicating that the break state is not set.

If all the above checks are passed, the receive data is stored in SCFRDR.

Note: When a parity error or a framing error occurs, reception is not suspended.

4. If the RIE bit in SCSCR is set to 1 when the RDF or DR flag changes to 1, a receive-FIFOdata-full interrupt (RXI) request is generated. If the RIE bit or the REIE bit in SCSCR is set to 1 when the ER flag changes to 1, a receive-error interrupt (ERI) request is generated. If the RIE bit or the REIE bit in SCSCR is set to 1 when the BRK or ORER flag changes to 1, a break reception interrupt (BRI) request is generated.

Figure 16.8 shows an example of the operation for reception.

Figure 16.8 Example of SCIF Receive Operation (8-Bit Data, Parity, 1 Stop Bit)

16.4.3 Operation in Clocked Synchronous Mode

In clocked synchronous mode, the SCIF transmits and receives data in synchronization with clock pulses. This mode is suitable for high-speed serial communication.

The SCIF transmitter and receiver are independent, so full-duplex communication is possible while sharing the same clock. The transmitter and receiver are also 16-byte FIFO buffered, so continuous transmitting or receiving is possible by reading or writing data while transmitting or receiving is in progress.

Figure 16.9 shows the general format in clocked synchronous serial communication.

In clocked synchronous serial communication, each data bit is output on the communication line from one falling edge of the serial clock to the next. Data is guaranteed valid at the rising edge of the serial clock.

In each character, the serial data bits are transmitted in order from the LSB (first) to the MSB (last). After output of the MSB, the communication line remains in the state of the MSB.

In clocked synchronous mode, the SCIF receives data by synchronizing with the rising edge of the serial clock.

(1) Transmit/Receive Formats

The data length is fixed at eight bits. No parity bit can be added.

(2) Clock

An internal clock generated by the on-chip baud rate generator by the setting of the C/\overline{A} bit in SCSMR and CKE[1:0] in SCSCR, or an external clock input from the SCK pin can be selected as the SCIF transmit/receive clock.

When the SCIF operates on an internal clock, it outputs the clock signal at the SCK pin. Eight clock pulses are output per transmitted or received character. When the SCIF is not transmitting or receiving, the clock signal remains in the high state. When only receiving, the clock signal outputs while the RE bit of SCSCR is 1 and the number of data in receive FIFO is more than the receive FIFO data trigger number.

(3) Transmitting and Receiving Data

• **SCIF Initialization (Clocked Synchronous Mode)**

Before transmitting, receiving, or changing the mode or communication format, the software must clear the TE and RE bits to 0 in the serial control register (SCSCR), then initialize the SCIF. Clearing TE to 0 initializes the transmit shift register (SCTSR). Clearing RE to 0, however, does not initialize the RDF, PER, FER, and ORER flags and receive data register (SCRDR), which retain their previous contents.

Figure 16.10 Sample Flowchart for SCIF Initialization

• **Transmitting Serial Data (Clocked Synchronous Mode)**

Figure 16.11 shows a sample flowchart for transmitting serial data.

Use the following procedure for serial data transmission after enabling the SCIF for transmission.

Figure 16.11 Sample Flowchart for Transmitting Serial Data
In serial transmission, the SCIF operates as described below.

- 1. When data is written into the transmit FIFO data register (SCFTDR), the SCIF transfers the data from SCFTDR to the transmit shift register (SCTSR) and starts transmitting. Confirm that the TDFE flag in the serial status register (SCFSR) is set to 1 before writing transmit data to SCFTDR. The number of data bytes that can be written is (16 – transmit trigger setting).
- 2. When data is transferred from SCFTDR to SCTSR and transmission is started, consecutive transmit operations are performed until there is no transmit data left in SCFTDR. When the number of transmit data bytes in SCFTDR falls below the transmit trigger number set in the FIFO control register (SCFCR), the TDFE flag is set. If the TIE bit in the serial control register (SCSR) is set to 1 at this time, a transmit-FIFO-data-empty interrupt (TXI) request is generated.

If clock output mode is selected, the SCIF outputs eight synchronous clock pulses. If an external clock source is selected, the SCIF outputs data in synchronization with the input clock. Data is output from the TxD pin in order from the LSB (bit 0) to the MSB (bit 7).

- 3. The SCIF checks the SCFTDR transmit data at the timing for sending the MSB (bit 7). If data is present, the data is transferred from SCFTDR to SCTSR, and then serial transmission of the next frame is started. If there is no data, the TxD pin holds the state after the TEND flag in SCFSR is set to 1 and the MSB (bit 7) is sent.
- 4. After the end of serial transmission, the SCK pin is held in the high state.

Figure 16.12 shows an example of SCIF transmit operation.

Figure 16.12 Example of SCIF Transmit Operation

• **Receiving Serial Data (Clocked Synchronous Mode)**

Figures 16.13 and 16.14 show sample flowcharts for receiving serial data. When switching from asynchronous mode to clocked synchronous mode without SCIF initialization, make sure that ORER, PER, and FER are cleared to 0.

Figure 16.13 Sample Flowchart for Receiving Serial Data (1)

Figure 16.14 Sample Flowchart for Receiving Serial Data (2)

In serial reception, the SCIF operates as described below.

- 1. The SCIF synchronizes with serial clock input or output and starts the reception.
- 2. Receive data is shifted into SCRSR in order from the LSB to the MSB. After receiving the data, the SCIF checks the receive data can be loaded from SCRSR into SCFRDR or not. If this check is passed, the RDF flag is set to 1 and the SCIF stores the received data in SCFRDR. If the check is not passed (overrun error is detected), further reception is prevented.
- 3. After setting RDF to 1, if the receive-data-full interrupt enable bit (RIE) is set to 1 in SCSCR, the SCIF requests a receive-data-full interrupt (RXI). If the ORER bit is set to 1 and the receive-data-full interrupt enable bit (RIE) or the receive error interrupt enable bit (REIE) in SCSCR is also set to 1, the SCIF requests a break interrupt (BRI).

Figure 16.15 shows an example of SCIF receive operation.

Figure 16.15 Example of SCIF Receive Operation

• **Transmitting and Receiving Serial Data Simultaneously (Clocked Synchronous Mode)**

Figure 16.16 shows a sample flowchart for transmitting and receiving serial data simultaneously.

Use the following procedure for the simultaneous transmission/reception of serial data, after enabling the SCIF for transmission/reception.

16.5 SCIF Interrupts

The SCIF has four interrupt sources: transmit FIFO data empty (TXI), receive error (ERI), receive FIFO data full (RXI), and break (BRI).

Table 16.12 shows the interrupt sources and their order of priority. The interrupt sources are enabled or disabled by means of the TIE, RIE, and REIE bits in SCSCR. A separate interrupt request is sent to the interrupt controller for each of these interrupt sources.

When a TXI request is enabled by the TIE bit and the TDFE flag in the serial status register (SCFSR) is set to 1, a TXI interrupt request is generated. The DMAC can be activated and data transfer performed by this TXI interrupt request. At this time, an interrupt request is not sent to the CPU.

When an RXI request is enabled by the RIE bit, and the RDF flag or the DR flag in SCFSR is set to 1, an RXI interrupt request is generated. The DMAC can be activated and data transfer performed by this RXI interrupt request. At this time, an interrupt request is not sent to the CPU. The RXI interrupt request caused by the DR flag is generated only in asynchronous mode.

When the RIE bit is set to 0 and the REIE bit is set to 1, the SCIF requests only an ERI or a BRI interrupt without requesting an RXI interrupt.

The TXI interrupt indicates that transmit data can be written, and the RXI interrupt indicates that there is receive data in SCFRDR.

Table 16.12 SCIF Interrupt Sources

16.6 Usage Notes

Note the following when using the SCIF.

16.6.1 SCFTDR Writing and TDFE Flag

The TDFE flag in the serial status register (SCFSR) is set when the number of transmit data bytes written in the transmit FIFO data register (SCFTDR) has fallen below the transmit trigger number set by bits TTRG1 and TTRG0 in the FIFO control register (SCFCR). After TDFE is set, transmit data up to the number of empty bytes in SCFTDR can be written, allowing efficient continuous transmission.

However, if the number of data bytes written in SCFTDR is equal to or less than the transmit trigger number, the TDFE flag will be set to 1 again after being read as 1 and cleared to 0. TDFE clearing should therefore be carried out when SCFTDR contains more than the transmit trigger number of transmit data bytes.

The number of transmit data bytes in SCFTDR can be found from the upper 8 bits of the FIFO data count register (SCFDR).

16.6.2 SCFRDR Reading and RDF Flag

The RDF flag in the serial status register (SCFSR) is set when the number of receive data bytes in the receive FIFO data register (SCFRDR) has become equal to or greater than the receive trigger number set by bits RTRG1 and RTRG0 in the FIFO control register (SCFCR). After RDF is set, receive data equivalent to the trigger number can be read from SCFRDR, allowing efficient continuous reception.

However, if the number of data bytes in SCFRDR exceeds the trigger number, the RDF flag will be set to 1 again if it is cleared to 0. RDF should therefore be cleared to 0 after being read as 1 after reading the number of the received data in the receive FIFO data register (SCFRDR) which is less than the trigger number.

The number of receive data bytes in SCFRDR can be found from the lower 8 bits of the FIFO data count register (SCFDR).

16.6.3 Restriction on DMAC Usage

- 1. When the DMAC writes data to SCFTDR with a TXI interrupt request, the state of the TEND flag becomes undefined. Therefore, the TEND flag should not be used as the transfer end flag in such a case.
- 2. When one channel is used in full duplex communication with the DMAC used for transmission and the CPU used for reception, if the receive data are read from the receive FIFO data register (SCFRDR) after the RDF or DR flag in the serial status register (SCFSR) has been set, the RDF or DR flag may be cleared.
- 3. When one channel is used in full duplex communication with the DMAC used for reception and the CPU used for transmission, if the transmit data is written to the transmit FIFO data register (SCFTDR) after the TDFE or TEND flag in the serial status register (SCFSR) has been set, the TDFE or TEND flags may be cleared.

16.6.4 Break Detection and Processing

Break signals can be detected by reading the RxD pin directly when a framing error (FER) is detected. In the break state the input from the RxD pin consists of all 0s, so the FER flag is set and the parity error flag (PER) may also be set.

Note that, although transfer of receive data to SCFRDR is halted in the break state, the SCIF receiver continues to operate.

16.6.5 Sending a Break Signal

The I/O condition and level of the TxD pin are determined by the SPB2IO and SPB2DT bits in the serial port register (SCSPTR). This feature can be used to send a break signal.

Until TE bit is set to 1 (enabling transmission) after initializing, the TxD pin does not work. During the period, mark status is performed by the SPB2DT bit. Therefore, the SPB2IO and SPB2DT bits should be set to 1 (high level output).

To send a break signal during serial transmission, clear the SPB2DT bit to 0 (designating low level), then clear the TE bit to 0 (halting transmission). When the TE bit is cleared to 0, the transmitter is initialized regardless of the current transmission state, and 0 is output from the TxD pin.

16.6.6 Receive Data Sampling Timing and Receive Margin (Asynchronous Mode)

The SCIF operates on a base clock with a frequency of 16 times the transfer rate. In reception, the SCIF synchronizes internally with the fall of the start bit, which it samples on the base clock. Receive data is latched at the rising edge of the eighth base clock pulse. The timing is shown in figure 16.17.

Figure 16.17 Receive Data Sampling Timing in Asynchronous Mode

The receive margin in asynchronous mode can therefore be expressed as shown in equation 1.

Equation 1:

$$
M = \left(0.5 - \frac{1}{2N}\right) - (L - 0.5) F - \frac{|D - 0.5|}{N} (1 + F) \right) \times 100\%
$$

Where: M: Receive margin $(\%)$

N: Ratio of clock frequency to bit rate $(N = 16)$

D: Clock duty ($D = 0$ to 1.0)

L: Frame length $(L = 9$ to 12)

F: Absolute deviation of clock frequency

From equation 1, if $F = 0$ and $D = 0.5$, the receive margin is 46.875%, as given by equation 2.

Equation 2:

When $D = 0.5$ and $F = 0$. $M = (0.5 - 1/(2 \times 16)) \times 100\%$ $= 46.875%$

This is a theoretical value. A reasonable margin to allow in system designs is 20% to 30%.

Section 17 I^2C Bus Interface 3 (IIC3)

The I^2C bus interface 3 conforms to and provides a subset of the Philips I^2C (Inter-IC) bus interface functions. However, the configuration of the registers that control the $I²C$ bus differs partly from the Philips register configuration.

17.1 Features

- Selection of I^2C format or clocked synchronous serial format
- Continuous transmission/reception

Since the shift register, transmit data register, and receive data register are independent from each other, the continuous transmission/reception can be performed.

I ²C bus format:

- Start and stop conditions generated automatically in master mode
- Selection of acknowledge output levels when receiving
- Automatic loading of acknowledge bit when transmitting
- Bit synchronization/wait function

In master mode, the state of SCL is monitored per bit, and the timing is synchronized automatically. If transmission/reception is not yet possible, set the SCL to low until preparations are completed.

• Six interrupt sources

Transmit data empty (including slave-address match), transmit end, receive data full (including slave-address match), arbitration lost, NACK detection, and stop condition detection

- The direct memory access controller (DMAC) can be activated by a transmit-data-empty request or receive-data-full request to transfer data.
- Direct bus drive

Two pins, SCL0 to SCL2 and SDA0 to SDA2, function as NMOS open-drain outputs when the bus drive function is selected.

Clocked synchronous serial format:

• Four interrupt sources

Transmit-data-empty, transmit-end, receive-data-full, and overrun error

• The direct memory access controller (DMAC) can be activated by a transmit-data-empty request or receive-data-full request to transfer data.

Figure 17.1 Block Diagram of I²C Bus Interface 3

17.2 Input/Output Pins

Table 17.1 shows the pin configuration of the I^2C bus interface 3. Specifications for the voltage applied to I/O pins for the I^2C bus interface are different from others because of the pin configuration difference. For details, see section 31, Electrical Characteristics.

Table 17.1 Pin Configuration

Channel	Pin Name	Symbol	I/O	Function
0 to 2	Serial clock	SCL0 to SCL2	I/O	I'C serial clock input/output
	Serial data	SDA0 to SDA2	1/O	I ² C serial data input/output

Figure 17.2 shows an example of I/O pin connections to external circuits.

Specifications for the voltage applied to I/O pins for the $I²C$ bus interface are different from others because of the pin configuration difference. For details, see section 31, Electrical Characteristics.

Figure 17.2 External Circuit Connections of I/O Pins

17.3 Register Descriptions

The I^2C bus interface 3 has the following registers.

Table 17.2 Register Configuration

17.3.1 I ²C Bus Control Register 1 (ICCR1)

ICCR1 is an 8-bit readable/writable register that enables or disables the I^2C bus interface 3, controls transmission or reception, and selects master or slave mode, transmission or reception, and transfer clock frequency in master mode.

ICCR1 is initialized to H'00 by a power-on reset or deep standby mode.

Table 17.3 Transfer Rate

Note: The settings should satisfy external specifications.

SH7261 Group

17.3.2 I ²C Bus Control Register 2 (ICCR2)

ICCR2 is an 8-bit readable/writable register that issues start/stop conditions, manipulates the SDA pin, monitors the SCL pin, and controls reset in the control part of the $I²C$ bus.

ICCR2 is initialized to H'7D by a power-on reset or deep standby mode.

17.3.3 I ²C Bus Mode Register (ICMR)

ICMR is an 8-bit readable/writable register that selects whether the MSB or LSB is transferred first, performs master mode wait control, and selects the transfer bit count.

ICMR is initialized to H'38 by a power-on reset or deep standby mode. Bits BC[2:0] are initialized to H'0 by the IICRST bit in ICCR2.

17.3.4 I ²C Bus Interrupt Enable Register (ICIER)

ICIER is an 8-bit readable/writable register that enables or disables interrupt sources and acknowledge bits, sets acknowledge bits to be transferred, and confirms acknowledge bits received.

ICIER is initialized to H'00 by a power-on reset or deep standby mode.

17.3.5 I ²C Bus Status Register (ICSR)

ICSR is an 8-bit readable/writable register that confirms interrupt request flags and their status.

ICSR is initialized to H'00 by a power-on reset or deep standby mode.

17.3.6 Slave Address Register (SAR)

SAR is an 8-bit readable/writable register that selects the communications format and sets the slave address. In slave mode with the I^2C bus format, if the upper seven bits of SAR match the upper seven bits of the first frame received after a start condition, this module operates as the slave device.

SAR is initialized to H'00 by a power-on reset or deep standby mode.

17.3.7 I ²C Bus Transmit Data Register (ICDRT)

ICDRT is an 8-bit readable/writable register that stores the transmit data. When ICDRT detects the empty space in the shift register (ICDRS), it transfers the transmit data which is written in ICDRT to ICDRS and starts transferring data. If the next transfer data is written to ICDRT during transferring data of ICDRS, continuous transfer is possible. ICDRT is initialized to H'FF.

ICDRT is initialized to H'FF by a power-on reset or deep standby mode.

17.3.8 I ²C Bus Receive Data Register (ICDRR)

ICDRR is an 8-bit register that stores the receive data. When data of one byte is received, ICDRR transfers the receive data from ICDRS to ICDRR and the next data can be received. ICDRR is a receive-only register, therefore the CPU cannot write to this register.

ICDRR is initialized to H'FF by a power-on reset or deep standby mode.

17.3.9 I ²C Bus Shift Register (ICDRS)

ICDRS is a register that is used to transfer/receive data. In transmission, data is transferred from ICDRT to ICDRS and the data is sent from the SDA pin. In reception, data is transferred from ICDRS to ICDRR after data of one byte is received. This register cannot be read directly from the CPU.

17.3.10 NF2CYC Register (NF2CYC)

NF2CYC is an 8-bit readable/writable register that selects the range of the noise filtering for the SCL and SDA pins. For details of the noise filter, see section 17.4.7, Noise Filter.

NF2CYC is initialized to H'02 by a power-on reset or in deep standby mode.

SH7261 Group

17.4 Operation

The I^2C bus interface 3 can communicate either in I^2C bus mode or clocked synchronous serial mode by setting FS in SAR.

17.4.1 I ²C Bus Format

Figure 17.3 shows the I^2C bus formats. Figure 17.4 shows the I^2C bus timing. The first frame following a start condition always consists of eight bits.

Figure 17.3 I²C Bus Formats

Figure 17.4 I²C Bus Timing

[Legend]

S: Start condition. The master device drives SDA from high to low while SCL is high.

SLA: Slave address

- R/\overline{W} : Indicates the direction of data transfer: from the slave device to the master device when R/W is 1, or from the master device to the slave device when R/W is 0.
- A: Acknowledge. The receive device drives SDA to low.

DATA: Transfer data

P: Stop condition. The master device drives SDA from low to high while SCL is high.

17.4.2 Master Transmit Operation

In master transmit mode, the master device outputs the transmit clock and transmit data, and the slave device returns an acknowledge signal. For master transmit mode operation timing, refer to figures 17.5 and 17.6. The transmission procedure and operations in master transmit mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set the WAIT bit in ICMR and bits CKS[3:0] in ICCR1. (Initial setting)
- 2. Read the BBSY flag in ICCR2 to confirm that the bus is released. Set the MST and TRS bits in ICCR1 to select master transmit mode. Then, write 1 to BBSY and 0 to SCP. (Start condition issued) This generates the start condition.
- 3. After confirming that TDRE in ICSR has been set, write the transmit data (the first byte data show the slave address and R/\overline{W}) to ICDRT. At this time, TDRE is automatically cleared to 0, and data is transferred from ICDRT to ICDRS. TDRE is set again.
- 4. When transmission of one byte data is completed while TDRE is 1, TEND in ICSR is set to 1 at the rise of the 9th transmit clock pulse. Read the ACKBR bit in ICIER, and confirm that the slave device has been selected. Then, write second byte data to ICDRT. When ACKBR is 1, the slave device has not been acknowledged, so issue the stop condition. To issue the stop condition, write 0 to BBSY and SCP. SCL is fixed low until the transmit data is prepared or the stop condition is issued.
- 5. The transmit data after the second byte is written to ICDRT every time TDRE is set.
- 6. Write the number of bytes to be transmitted to ICDRT. Wait until TEND is set (the end of last byte data transmission) while TDRE is 1, or wait for NACK (NACKF in $ICSR = 1$) from the receive device while ACKE in ICIER is 1. Then, issue the stop condition to clear TEND or NACKF.
- 7. When the STOP bit in ICSR is set to 1, the operation returns to the slave receive mode.

Figure 17.5 Master Transmit Mode Operation Timing (1)

Figure 17.6 Master Transmit Mode Operation Timing (2)

17.4.3 Master Receive Operation

In master receive mode, the master device outputs the receive clock, receives data from the slave device, and returns an acknowledge signal. For master receive mode operation timing, refer to figures 17.7 and 17.8. The reception procedure and operations in master receive mode are shown below.

- 1. Clear the TEND bit in ICSR to 0, then clear the TRS bit in ICCR1 to 0 to switch from master transmit mode to master receive mode. Then, clear the TDRE bit to 0.
- 2. When ICDRR is read (dummy data read), reception is started*, and the receive clock is output, and data received, in synchronization with the internal clock. The master device outputs the level specified by ACKBT in ICIER to SDA, at the 9th receive clock pulse.
- 3. After the reception of first frame data is completed, the RDRF bit in ICSR is set to 1 at the rise of 9th receive clock pulse. At this time, the receive data is read by reading ICDRR, and RDRF is cleared to 0.
- 4. The continuous reception is performed by reading ICDRR every time RDRF is set. If 8th receive clock pulse falls after reading ICDRR by the other processing while RDRF is 1, SCL is fixed low until ICDRR is read.
- 5. If next frame is the last receive data, set the RCVD bit in ICCR1 to 1 before reading ICDRR. This enables the issuance of the stop condition after the next reception.
- 6. When the RDRF bit is set to 1 at rise of the 9th receive clock pulse, issue the stage condition.
- 7. When the STOP bit in ICSR is set to 1, read ICDRR. Then clear the RCVD bit to 0.
- 8. The operation returns to the slave receive mode.
- Note: * If only one byte is received, read ICDRR (dummy-read) after the RCVD bit in ICCR1 is set.

Figure 17.7 Master Receive Mode Operation Timing (1)

Figure 17.8 Master Receive Mode Operation Timing (2)

17.4.4 Slave Transmit Operation

In slave transmit mode, the slave device outputs the transmit data, while the master device outputs the receive clock and returns an acknowledge signal. For slave transmit mode operation timing, refer to figures 17.9 and 17.10.

The transmission procedure and operations in slave transmit mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting) Set the MST and TRS bits in ICCR1 to select slave receive mode, and wait until the slave address matches.
- 2. When the slave address matches in the first frame following detection of the start condition, the slave device outputs the level specified by ACKBT in ICIER to SDA, at the rise of the 9th clock pulse. At this time, if the 8th bit data (R/W) is 1, the TRS bit in ICCR1 and the TDRE bit in ICSR are set to 1, and the mode changes to slave transmit mode automatically. The continuous transmission is performed by writing transmit data to ICDRT every time TDRE is set.
- 3. If TDRE is set after writing last transmit data to ICDRT, wait until TEND in ICSR is set to 1, with TDRE = 1. When TEND is set, clear TEND.
- 4. Clear TRS for the end processing, and read ICDRR (dummy read). SCL is opened.
- 5. Clear TDRE.

Figure 17.9 Slave Transmit Mode Operation Timing (1)

Figure 17.10 Slave Transmit Mode Operation Timing (2)

17.4.5 Slave Receive Operation

In slave receive mode, the master device outputs the transmit clock and transmit data, and the slave device returns an acknowledge signal. For slave receive mode operation timing, refer to figures 17.11 and 17.12. The reception procedure and operations in slave receive mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting) Set the MST and TRS bits in ICCR1 to select slave receive mode, and wait until the slave address matches.
- 2. When the slave address matches in the first frame following detection of the start condition, the slave device outputs the level specified by ACKBT in ICIER to SDA, at the rise of the 9th clock pulse. At the same time, RDRF in ICSR is set to read ICDRR (dummy read). (Since the read data show the slave address and R/W, it is not used.)
- 3. Read ICDRR every time RDRF is set. If 8th receive clock pulse falls while RDRF is 1, SCL is fixed low until ICDRR is read. The change of the acknowledge before reading ICDRR, to be returned to the master device, is reflected to the next transmit frame.
- 4. The last byte data is read by reading ICDRR.

Figure 17.11 Slave Receive Mode Operation Timing (1)

Figure 17.12 Slave Receive Mode Operation Timing (2)

17.4.6 Clocked Synchronous Serial Format

This module can be operated with the clocked synchronous serial format, by setting the FS bit in SAR to 1. When the MST bit in ICCR1 is 1, the transfer clock output from SCL is selected. When MST is 0, the external clock input is selected.

(1) Data Transfer Format

Figure 17.13 shows the clocked synchronous serial transfer format.

The transfer data is output from the fall to the fall of the SCL clock, and the data at the rising edge of the SCL clock is guaranteed. The MLS bit in ICMR sets the order of data transfer, in either the MSB first or LSB first. The output level of SDA can be changed during the transfer wait, by the SDAO bit in ICCR2.

Figure 17.13 Clocked Synchronous Serial Transfer Format

(2) Transmit Operation

In transmit mode, transmit data is output from SDA, in synchronization with the fall of the transfer clock. The transfer clock is output when MST in ICCR1 is 1, and is input when MST is 0. For transmit mode operation timing, refer to figure 17.14. The transmission procedure and operations in transmit mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set the MST and CKS[3:0] bits in ICCR1. (Initial setting)
- 2. Set the TRS bit in ICCR1 to select the transmit mode. Then, TDRE in ICSR is set.
- 3. Confirm that TDRE has been set. Then, write the transmit data to ICDRT. The data is transferred from ICDRT to ICDRS, and TDRE is set automatically. The continuous transmission is performed by writing data to ICDRT every time TDRE is set. When changing from transmit mode to receive mode, clear TRS while TDRE is 1.

Figure 17.14 Transmit Mode Operation Timing

(3) Receive Operation

In receive mode, data is latched at the rise of the transfer clock. The transfer clock is output when MST in ICCR1 is 1, and is input when MST is 0. For receive mode operation timing, refer to figure 17.15. The reception procedure and operations in receive mode are described below.

- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting)
- 2. When the transfer clock is output, set MST to 1 to start outputting the receive clock.
- 3. When the receive operation is completed, data is transferred from ICDRS to ICDRR and RDRF in ICSR is set. When $MST = 1$, the next byte can be received, so the clock is continually output. The continuous reception is performed by reading ICDRR every time RDRF is set. When the 8th clock is risen while RDRF is 1, the overrun is detected and AL/OVE in ICSR is set. At this time, the previous reception data is retained in ICDRR.
- 4. To stop receiving when MST = 1, set RCVD in ICCR1 to 1, then read ICDRR. Then, SCL is fixed high after receiving the next byte data.
- Notes: Follow the steps below to receive only one byte with MST = 1 specified. See figure 17.16 for the operation timing.
	- 1. Set the ICE bit in ICCR1 to 1. Set bits CKS[3:0] in ICCR1. (Initial setting)
	- 2. Set MST = 1 while the RCVD bit in ICCR1 is 0. This causes the receive clock to be output.
	- 3. Check if the BC2 bit in ICMR is set to 1 and then set the RCVD bit in ICCR1 to 1. This causes the SCL to be fixed to the high level after outputting one byte of the receive clock.

Figure 17.15 Receive Mode Operation Timing

Figure 17.16 Operation Timing For Receiving One Byte (MST = 1)

17.4.7 Noise Filter

The logic levels at the SCL and SDA pins are routed through noise filters before being latched internally. Figure 17.17 shows a block diagram of the noise filter circuit.

The noise filter consists of three cascaded latches and a match detector. The SCL (or SDA) input signal is sampled on the peripheral clock. When NF2CYC is set to 0, this signal is not passed forward to the next circuit unless the outputs of both latches agree. When NF2CYC is set to 1, this signal is not passed forward to the next circuit unless the outputs of three latches agree. If they do not agree, the previous value is held.

Figure 17.17 Block Diagram of Noise Filter

17.4.8 Example of Use

Flowcharts in respective modes that use the I^2C bus interface 3 are shown in figures 17.18 to 17.21.

Figure 17.18 Sample Flowchart for Master Transmit Mode

Figure 17.19 Sample Flowchart for Master Receive Mode

Figure 17.21 Sample Flowchart for Slave Receive Mode

17.5 Interrupt Requests

There are six interrupt requests in this module; transmit data empty, transmit end, receive data full, NACK detection, STOP recognition, and arbitration lost/overrun error. Table 17.4 shows the contents of each interrupt request.

Table 17.4 Interrupt Requests

When the interrupt condition described in table 17.4 is 1, the CPU executes an interrupt exception handling. Note that a TXI or RXI interrupt can activate the DMAC if the setting for DMAC activation has been made. In such a case, an interrupt request is not sent to the CPU. Interrupt sources should be cleared in the exception handling. The TDRE and TEND bits are automatically cleared to 0 by writing the transmit data to ICDRT. The RDRF bit is automatically cleared to 0 by reading ICDRR. The TDRE bit is set to 1 again at the same time when the transmit data is written to ICDRT. Therefore, when the TDRE bit is cleared to 0, then an excessive data of one byte may be transmitted.

17.6 Bit Synchronous Circuit

In master mode, this module has a possibility that high level period may be short in the two states described below.

- When SCL is driven to low by the slave device
- When the rising speed of SCL is lowered by the load of the SCL line (load capacitance or pullup resistance)

Therefore, it monitors SCL and communicates by bit with synchronization.

Figure 17.22 shows the timing of the bit synchronous circuit and table 17.5 shows the time when the SCL output changes from low to Hi-Z then SCL is monitored.

Table 17.5 Time for Monitoring SCL

Notes: 1. Monitors the (on-board) SCL level after the time (pcyc) for monitoring SCL has passed since the rising edge of the SCL monitor timing reference clock.

2. $pcyc = P\phi \times cyc$

17.7 Usage Note

17.7.1 Issuance of Stop Condition and Start Condition (Retransmission)

Issue a start (retransmission) or stop condition after the falling edge of the 9th clock has been recognized. The falling edge of the 9th clock can be recognized by checking the SCLO bit in the I 2 C bus control register 2 (ICCR2). When a start (retransmission) or stop condition is issued with a certain timing under the following conditions (1 or 2), the start (retransmission) or stop condition may not be output correctly.

- 1. SCL takes longer to rise than the period defined in section 17.6, Bit Synchronous Circuit, due to the load of the SCL bus (load capacitance or pull-up resistance).
- 2. The low-level period between the 8th and 9th clock is prolonged by the slave device, which activates the bit synchronous circuit.

17.7.2 Note on Setting for Multi-Master Operation

In multi-master operation, when the transfer rate setting for this module (ICCR1.CKS[3:0]) makes this LSI slower than the other masters, pulse cycles with an unexpected length will infrequently be output on SCL.

Be sure to specify a transfer rate that is at least 1/1.8 of the fastest transfer rate among the other masters.

17.7.3 Note on Master Receive Mode

Reading ICDRR around the falling edge of the 8th clock might fail to fetch the receive data.

In addition, when RCVD is set to 1 around the falling edge of the 8th clock and the receive buffer full, a stop condition may not be issued.

Use either 1 or 2 below as a measure against the situations above.

- 1. In master receive mode, read ICDRR before the rising edge of the 8th clock.
- 2. In master receive mode, set the RCVD bit to 1 so that transfer proceeds in byte units.

17.7.4 Note on Setting ACKBT in Master Receive Mode

In master receive mode operation, set ACKBT before the falling edge of the 8th SCL cycle of the last data being continuously transferred. Not doing so can lead to an overrun for the slave transmission device.

17.7.5 Note on the States of Bits MST and TRN when Arbitration is Lost

When sequential bit-manipulation instructions are used to set the MST and TRS bits to select master transmission in multi-master operation, a conflicting situation where AL in $ICSR = 1$ but the mode is master transmit mode ($MST = 1$ and $TRS = 1$) may arise; this depends on the timing of the loss of arbitration when the bit manipulation instruction for TRS is executed.

This can be avoided in either of the following ways.

- In multi-master operation, use the MOV instruction to set the MST and TRS bits.
- When arbitration is lost, check whether the MST and TRS bits are 0. If the MST and TRS bits have been set to a value other than 0, clear the bits to 0.

17.7.6 Note on IICRST and BBSY bits

When 1 is written to IICRST in ICCR2, this LSI release SCL and SDA pins. Then, if the SDA level changes from low to high under the condition of SCL = high, BBSY in ICCR2 is cleared to 0 assuming that the stop condition has been issued.

Section 18 Serial Sound Interface (SSI)

The serial sound interface (hereinafter referred to as the "SSI") is a transceiver module designed to send or receive audio data interface with a variety of devices compatible with I^2S bus. It also provides additional modes for other common formats as well as multi-channel mode.

18.1 Features

- Number of channels: Two channels
- Operating mode: Non-compressed mode
	- The non-compressed mode supports serial audio streams divided by channels.
- Serves as both a transmitter and a receiver
- Capable of using serial bus format
- Asynchronous transfer takes place between the data buffer and the shift register.
- It is possible to select a value as the dividing ratio for the clock used by the serial but interface.
- It is possible to control data transmission or reception with DMAC and interrupt requests.
- Selects the oversample clock from among the pins AUDIO CLK, or AUDIO X1 and AUDIO_X2.
	- External clock frequency input through the pins AUDIO CLK, or AUDIO X1 and AUDIO_X2: 1 to 40 MHz
	- Crystal oscillator frequency for the pins AUDIO X1 and AUDIO X2: 10 to 25 MHz

Figure 18.1 shows a schematic diagram of the four channels in the SSI module.

Figure 18.1 Schematic Diagram of SSI Module

Figure 18.2 shows a block diagram of the SSI module when it is used alone.

18.2 Input/Output Pins

Table 18.1 shows the pin assignments relating to the SSI module.

Table 18.1 Pin Assignments

18.3 Register Description

The SSI has the following registers. Note that explanation in the text does not refer to the channels.

Table 18.2 Register Description

Note: * For this register, bits 26 and 27 are capable of reading and writing, although the others are read-only bits. For details, refer to section 18.3.2, Status Register (SSISR).

18.3.1 Control Register (SSICR)

SSICR is a readable/writable 32-bit register that controls the IRQ, selects the polarity status, and sets operating mode.

SSICR is initialized to H'00000000 by a power-on reset or in deep standby mode.

18.3.2 Status Register (SSISR)

SSISR consists of status flags indicating the operational status of the SSI module and bits indicating the current channel numbers and word numbers.

SSISR is initialized to H'02000003 by a power-on reset or in deep standby mode.

Notes: 1. This bit can be read from or written to. Writing 0 initializes the bit, but writing 1 is ignored. 2. The SSI clock must be kept supplied until the SSI is in the idle state.

ignored.

2. The SSI clock must be kept supplied until the SSI is in the idle state.

18.3.3 Transmit Data Register (SSITDR)

SSITDR is a 32-bit register that stores data to be transmitted.

Data written to this register is transferred to the shift register upon transmission request. If the data word length is less than 32 bits, the alignment is determined by the setting of the PDTA control bit in SSICR. The data in the buffer can be accessed by reading this register.

SSITDR is initialized to H'00000000 by a power-on reset or in deep standby mode.

18.3.4 Receive Data Register (SSIRDR)

SSIRDR is a 32-bit register that stores receive messages.

Data in this register is transferred from the shift register each time data word is received. If the data word length is less than 32 bits, the alignment is determined by the setting of the PDTA control bit in SSICR.

SSIRDR is initialized to H'00000000 by a power-on reset or in deep standby mode.

18.4 Operation Description

18.4.1 Bus Format

The SSI module can operate as a transmitter or a receiver and can be configured into many serial bus formats in either mode.

The bus format can be selected from one of the four major modes shown in table 18.3.

18.4.2 Non-Compressed Modes

The non-compressed modes support all serial audio streams split into channels. It supports I^2S compatible format as well as many more variants on these modes.

(1) Slave Receiver

This mode allows the module to receive serial data from another device. The clock and word select signal used for the serial data stream is also supplied from an external device. If these signals do not conform to the format specified in the configuration fields of the SSI module, operation is not guaranteed.

(2) Slave Transmitter

This mode allows the module to transmit serial data to another device. The clock and word select signal used for the serial data stream is also supplied from an external device. If these signals do not conform to the format specified in the configuration fields of the SSI module, operation is not guaranteed.

(3) Master Receiver

This mode allows the module to receive serial data from another device. The clock and word select signals are internally derived from the oversampling clock. The format of these signals is defined in the configuration fields of the SSI module. If the incoming data does not follow the configured format, operation is not guaranteed.

(4) Master Transmitter

This mode allows the module to transmit serial data to another device. The clock and word select signals are internally derived from the oversampling clock. The format of these signals is defined in the configuration fields of the SSI module.

(5) Operating Setting Related to Word Length

All bits related to the SSICR's word length are valid in non-compressed modes. The SSI module supports many configurations, but the formats described below are I^2S compatible, MSB-first leftaligned, and MSB-first right-aligned.

1. I^2S Compatible Format

Figures 18.3 and 18.4 demonstrate the supported I^2S compatible format both with and without padding. Padding occurs when the data word length is smaller than the system word length.

Figure 18.3 I² S Compatible Format (without Padding)

Figure 18.4 I² S Compatible Format (with Padding)

Figure 18.5 shows MSB-first left-aligned format, and figure 18.6 shows MSB-first right-aligned format.

2. MSB-First Left-Aligned Format

Figure 18.5 MSB-First Left-Aligned Format (Transmitted and Received in the order of Serial Data and Padding Bits)

3. MSB-First Right-Aligned Format

Figure 18.6 MSB-First Right-Aligned Format (Transmitted and Received in the order of Padding Bits and Serial Data)

(6) Multi-channel Formats

Some devices extend the definition of the specification by I^2S bus and allow more than 2 channels to be transferred within two system words.

The SSI module supports the transfer of 4, 6 and 8 channels by using the CHNL, SWL and DWL bits only when the system word length (SWL) is greater than or equal to the data word length (DWL) multiplied by channels (CHNL).

Table 18.4 shows the number of padding bits for each of the valid setting. If setting is not valid, "—" is indicated instead of a number.

Padding Bits	Per System Word		DWL[2:0] 000		001	010	011	100	101	110
CHNL [1:0]	Decoded Channels per System Word	SWL [2:0]	Decoded Word Length	8	16	18	20	22	24	32
00	$\mathbf{1}$	000	8	0						
		001	16	8	0					
		010	24	16	8	6	$\overline{4}$	2	$\mathbf 0$	
		011	32	24	16	14	12	10	8	0
		100	48	40	32	30	28	26	24	16
		101	64	56	48	46	44	42	40	32
		110	128	120	112	110	108	106	104	96
		111	256	248	240	238	236	234	232	224
01	\overline{c}	000	8							
		001	16	0						
		010	24	8						
		011	32	16	0					
		100	48	32	16	12	8	4	0	
		101	64	48	32	28	24	20	16	0
		110	128	112	96	92	88	84	80	64
		111	256	240	224	220	216	212	208	192

Table 18.4 The Number of Padding Bits for Each Valid Setting

When the SSI module acts as a transmitter, each word written to SSITDR is transmitted to the serial audio bus in the order they are written. When the SSI module acts as a receiver, each word received by the serial audio bus is read in the order received from the SSIRDR register.

Figures 18.7 to 18.9 show how 4, 6 and 8 channels are transferred to the serial audio bus. Note that there are no padding bits in the first example, the second example is left-aligned and the third is right-aligned. This selection is arbitrary and is just for demonstration purposes only.

Figure 18.7 Multichannel Format (4 Channels Without Padding)

Figure 18.8 Multichannel Format (6 Channels with High Padding)

Figure 18.9 Multichannel Format (8 Channels; Transmitting and Receiving in the order of Padding Bits and Serial Data; with Padding)

(7) Bit Setting Configuration Format

Several more configuration bits in non-compressed mode are shown below. These bits are not mutually exclusive, but some combinations may not be useful for any other device.

These configuration bits are described below with reference to figure 18.10, Basic Sample Format.

Figure 18.10 Basic Sample Format (Transmit Mode with Example System/Data Word Length)

Figure 18.10 uses a system word length of 6 bits and a data word length of 4 bits. These settings are not possible with the SSI module but are used only for clarification of the other configuration bits.

1. Inverted Clock

Figure 18.11 Inverted Clock

2. Inverted Word Select

Figure 18.12 Inverted Word Select

3. Inverted Padding Polarity

Figure 18.13 Inverted Padding Polarity

4. Transmitting and Receiving in the Order of Padding Bits and Serial Data; with Delay

Figure 18.14 Transmitting and Receiving in the Order of Padding Bits and Serial Data; with Delay

5. Transmitting and Receiving in the Order of Padding Bits and Serial Data; without Delay

Figure 18.15 Transmitting and Receiving in the Order of Padding Bits and Serial Data; without Delay

6. Transmitting and Receiving in the Order of Serial Data and Padding Bits; without Delay

Figure 18.16 Transmitting and Receiving in the Order of Serial Data and Padding Bits; without Delay

7. Parallel Right-Aligned with Delay

Figure 18.17 Parallel Right-Aligned with Delay

8. Mute Enabled

Figure 18.18 Mute Enabled

18.4.3 Operation Modes

There are three modes of operation: configuration, enabled and disabled. Figure 18.19 shows how the module enters each of these modes.

Figure 18.19 Operation Modes

(1) Configuration Mode

This mode is entered after the module is released from reset. All required configuration fields in the control register should be defined in this mode, before the SSI module is enabled by setting the EN bit.

Setting the EN bit causes the module to enter the module enabled mode.

(2) Module Enabled Mode

Operation of the module in this mode is dependent on the operation mode selected. For details, refer to section 18.4.4, Transmit Operation and section 18.4.5, Receive Operation, below.

18.4.4 Transmit Operation

Transmission can be controlled either by DMA or interrupt.

DMA control is preferred to reduce the processor load. In DMA control mode the processor will only receive interrupts if there is an underflow or overflow of data or the DMAC has finished its transfer.

The alternative method is using the interrupts that the SSI module generates to supply data as required. This mode has a higher interrupt load as the module is only double buffered and will require data to be written at least every system word period.

When disabling the module, the SSI clock* must remain present until the SSI module is in idle state, indicated by the IIRQ bit.

Figure 18.20 shows the transmit operation in DMA control mode, and figure 18.21 shows the transmit operation in interrupt control mode.

Note: $*$ Input clock from the SSISCK pin when SCKD = 0. Oversampling clock when $SCKD = 1$.

(1) Transmission Using DMA Controller

(2) Transmission using Interrupt Data Flow Control

Figure 18.21 Transmission Using Interrupt Data Flow Control

18.4.5 Receive Operation

Like transmission, reception can be controlled either by DMA or interrupt.

Figures 18.22 and 18.23 show the flow of operation.

When disabling the SSI module, the SSI clock* must be kept supplied until the IIRQ bit is in idle state.

Note: $*$ Input clock from the SSISCK pin when SCKD = 0. Oversampling clock when $SCKD = 1$.

(1) Reception Using DMA Controller

Figure 18.22 Reception Using DMA Controller

(2) Reception Using Interrupt Data Flow Control

Figure 18.23 Reception Using Interrupt Data Flow Control

When an underflow or overflow error condition has matched, the CHNO [1:0] bit and the SWNO bit can be used to recover the SSI module to a known status. When an underflow or overflow occurs, the host can read the channel number and system word number to determine what point the serial audio stream has reached. In the transmitter case, the host can skip forward through the data it wants to transmit until it finds the sample data that matches what the SSI module is expecting to transmit next, and so resynchronize with the audio data stream. In the receiver case the host CPU can store null data to make the number of receive data items consistent until it is ready to store the sample data that the SSI module is indicating will be received next, and so resynchronize with the audio data stream.

18.4.6 Temporary Stop and Restart Procedures in Transmit Mode

The following procedures can be used for implementation.

(1) Procedure for the Repeated Transfer and Stop without having to Reconfigure the DMAC

- 1. Set SSICR.DMEN = 0 (disabling a DMA request) to stop the DMA transfer.
- 2. Wait for SSISR.DIRQ = 1 (transmit mode: the transmit buffer is empty) using a polling, interrupt, or the like.
- 3. With SSICR.EN = 0 (disabling an SSI module operation), stop the transfer.
- 4. Before attempting another transfer, make sure that $SSISR.IDST = 1$ is reached.
- 5. Set SSICR.EN = 1 (enabling an SSI module operation).
- 6. Wait for SSISR.DIRQ = 1, using a polling, interrupt, or the like.
- 7. Setting SSICR.DMEN = 1 (enabling a DMA request) will restart the DMA transfer.

(2) Procedure for Reconfiguring the DMAC after an SSI stop

- 1. Set SSICR.DMEN = 0 (disabling a DMA request) to stop the DMA transfer.
- 2. Wait for SSISR.DIRQ = 1 (transmit mode: the transmit buffer is empty), using a polling, interrupt, or the like.
- 3. With $SSCR. EN = 0$ (disabling an SSI module operation), stop the transfer.
- 4. Stop the DMAC with DMSCNT of the DMAC.
- 5. Before attempting another transfer, make sure that $SSISR.IDST = 1$ is reached.
- 6. Set SSICR.EN = 1 (enabling an SSI module operation).
- 7. Set the DMAC registers and start the transfer.
- 8. Setting SSICR.DMEN = 1 (enabling a DMA request) will restart the DMA transfer.

18.4.7 Serial Bit Clock Control

This function is used to control and select which clock is used for the serial bus interface.

If the serial clock direction is set to input $(SCKD = 0)$, the SSI module is in clock slave mode and the shift register uses the bit clock that was input to the SSISCK pin.

If the serial clock direction is set to output $(SCKD = 1)$, this module is in clock master mode, and the shift register uses the oversampling clock or a divided oversampling clock as the bit clock. The oversampling clock is divided by the ratio specified by the serial oversampling clock division ratio bits (CKDV) in SSICR for use as the bit clock by the shift register.

In either case the module pin, SSISCK, is the same as the bit clock.

18.5 Usage Notes

18.5.1 Limitations from Overflow during Receive DMA Operation

If an overflow occurs while the receive DMA is in operation, the module should be restarted. The receive buffer in the SSI consists of 32-bit registers that share the L and R channels. Therefore, data to be received at the L channel may sometimes be received at the R channel if an overflow occurs, for example, under the following condition: the control register (SSICR) has a 32-bit setting for both data word length (DWL2 to DWL0) and system word length (SWL2 to SWL0).

If an overflow is confirmed with the overflow error interrupt or overflow error status flag (the OIRQ bit in SSISR), write 0 to the EN bit in SSICR and DMEN bit to disable DMA in the SSI module, thus stopping the operation. (In this case, the controller setting should also be stopped.) After this, write 0 to the OIRQ bit to clear the overflow status, set DMA again and restart the transfer.

18.5.2 Note on Using Oversample Clock

To use the externally input clock as the oversample clock, refer to the section 4.6.1, Note on Inputting External Clock, in which the terms EXTAL and XTAL pins should be replaced by the AUDIO_X1 and AUDIO_X2 pins respectively.

To use the crystal resonator, refer to the section 4.6.2, Note on Using Crystal Resonator, in which the terms EXTAL and XTAL pins should be replaced by the AUDIO_X1 and AUDIO_X2 pins respectively.

Also, see section 4.6.3, Note on Resonator.

18.5.3 Restriction on Stopping Clock Supply

Once the bits MSTP53 and MSTP52 in the standby control register 5 (STBCR5) are cleared to 0 and the SSI operation is started, do not set these bits to 1 (stops clock supply to the SSI).

Section 19 Controller Area Network (RCAN-ET) [R5S72611] [R5S72613]

19.1 Summary

19.1.1 Overview

This document primarily describes the programming interface for the RCAN-ET module. It serves to facilitate the hardware/software interface so that engineers involved in the RCAN-ET implementation can ensure the design is successful.

19.1.2 Scope

The CAN Data Link Controller function is not described in this document. It is the responsibility of the reader to investigate the CAN Specification Document (see references). The interfaces from the CAN Controller are described, in so far as they pertain to the connection with the User Interface.

The programming model is described in some detail. It is not the intention of this document to describe the implementation of the programming interface, but to simply present the interface to the underlying CAN functionality.

The document places no constraints upon the implementation of the RCAN-ET module in terms of process, packaging or power supply criteria. These issues are resolved where appropriate in implementation specifications.

19.1.3 Audience

In particular this document provides the design reference for software authors who are responsible for creating a CAN application using this module.

In the creation of the RCAN-ET user interface LSI engineers must use this document to understand the hardware requirements.

19.1.4 References

- 1. CAN License Specification, Robert Bosch GmbH, 1992
- 2. CAN Specification Version 2.0 part A, Robert Bosch GmbH, 1991
- 3. CAN Specification Version 2.0 part B, Robert Bosch GmbH, 1991
- 4. Implementation Guide for the CAN Protocol, CAN Specification 2.0 Addendum, CAN In Automation, Erlangen, Germany, 1997
- 5. Road vehicles Controller area network (CAN): Part 1: Data link layer and physical signalling (ISO-11898-1, 2002)

19.1.5 Features

- Supports CAN specification 2.0B
- Bit timing compliant with ISO-11898-1
- 16 Mailbox version
- Clock 16 to 40 MHz
- 15 programmable Mailboxes for transmit/receive $+1$ receive-only mailbox
- Sleep mode for low power consumption and automatic recovery from sleep mode by detecting CAN bus activity
- Programmable receive filter mask (standard and extended identifier) supported by all Mailboxes
- Programmable CAN data rate up to 1MBit/s
- Transmit message queuing with internal priority sorting mechanism against the problem of priority inversion for real-time applications
- Data buffer access without SW handshake requirement in reception
- Flexible micro-controller interface
- Flexible interrupt structure

19.2 Architecture

19.2.1 Block Diagram

The RCAN-ET device offers a flexible and sophisticated way to organise and control CAN frames, providing the compliance to CAN2.0B Active and ISO-11898-1. The module is formed from 5 different functional entities. These are the Micro Processor Interface (MPI), Mailbox, Mailbox Control and CAN Interface. The figure below shows the block diagram of the RCAN-ET Module. The bus interface timing is designed according to the peripheral bus I/F required for each product.

Figure 19.1 RCAN-ET Architecture

Important: Although core of RCAN-ET is designed based on a 32-bit bus system, the whole RCAN-ET including MPI for the CPU has 16-bit bus interface to CPU. LongWord (32-bit) accesses are converted into two consecutive word accesses by the bus interface.

19.2.2 Functions of Each Block

(1) Micro Processor Interface (MPI)

The MPI allows communication between the Renesas CPU and RCAN-ET's registers/mailboxes to control the memory interface. It also contains the Wakeup Control logic that detects the CAN bus activities and notifies the MPI and the other parts of RCAN-ET so that the RCAN-ET can automatically exit the Sleep mode.

It contains registers such as MCR, IRR, GSR and IMR.

(2) Mailbox

The Mailboxes consists of RAM configured as message buffers and registers. There are 16 Mailboxes, and each mailbox has the following information.

\angle RAM \angle

- CAN message control (identifier, rtr, ide, etc)
- CAN message data (for CAN Data frames)
- Local Acceptance Filter Mask for reception
- <Registers>
- CAN message control (dlc)
- 3-bit wide Mailbox Configuration, Disable Automatic Re-Transmission bit, Auto-Transmission for Remote Request bit, New Message Control bit

(3) Mailbox Control

The Mailbox Control handles the following functions:

- For received messages, compare the IDs and generate appropriate RAM addresses/data to store messages from the CAN Interface into the Mailbox and set/clear appropriate registers accordingly.
- To transmit messages, RCAN-ET will run the internal arbitration to pick the correct priority message, and load the message from the Mailbox into the Tx-buffer of the CAN Interface and set/clear appropriate registers accordingly.
- Arbitrates Mailbox accesses between the CPU and the Mailbox Control.
- Contains registers such as TXPR, TXCR, TXACK, ABACK, RXPR, RFPR, UMSR and **MBIMR**

(4) CAN Interface

This block conforms to the requirements for a CAN Bus Data Link Controller which is specified in Ref. [3, 5]. It fulfils all the functions of a standard Data Link Controller as specified by the OSI 7 Layer Reference model. This functional entity also provides the registers and the logic which are specific to a given CAN bus, which includes the Receive Error Counter, Transmit Error Counter, the Bit Configuration Registers and various useful Test Modes. This block also contains functional entities to hold the data received and the data to be transmitted for the CAN Data Link Controller.

19.2.3 Input/Output Pins

Table 19.1 shows the pin configuration of the RCAN-ET.

Table 19.1 Pin Configuration

19.2.4 Memory Map

The diagram of the memory map is shown below.

Figure 19.2 RCAN-ET Memory Map

19.3 Mailbox

19.3.1 Mailbox Structure

Mailboxes play a role as message buffers to transmit/receive CAN frames. Each Mailbox is comprised of 3 identical storage fields that are 1): Message Control, 2): Local Acceptance Filter Mask, 3): Message Data. The following table shows the address map for the control, LAFM, data and addresses for each mailbox.

Table 19.2 Address Map for Each Mailbox

Mailbox-0 is a receive-only box, and all the other Mailboxes can operate as both receive and transmit boxes, dependant upon the MBC (Mailbox Configuration) bits in the Message Control. The following diagram shows the structure of a Mailbox in detail.

Table 19.3 Roles of Mailboxes

Notes: 1. All bits shadowed in grey are reserved and the write value should be 0. The value returned by a read may not always be 0 and should not be relied upon. 2. MBC1 bit in mailbox is fixed to 1.

3. ATX and DART are not supported by mailbox-0, and the MBC setting of mailbox-0 is limited.

4. When the MCR15 bit is 1, the order of STDID, RTR, IDE and EXTID of both message control and LAFM differs from HCAN2.

 $5. n = 0 to 15 (mailbox number)$

Figure 19.3 Mailbox-n Structure

19.3.2 Message Control Field

STDID[10:0]: These bits set the identifier (standard identifier) of data frames and remote frames.

EXTID[17:0]: These bits set the identifier (extended identifier) of data frames and remote frames.

RTR (Remote Transmission Request bit): Used to distinguish between data frames and remote frames. This bit is overwritten by received CAN Frames depending on Data Frames or Remote Frames.

Important: Please note that, when ATX bit is set with the setting MBC = B'001, the RTR bit will never be set. When a Remote Frame is received, the CPU can be notified by the corresponding RFPR set or IRR[2] (Remote Frame Request Interrupt), however, as RCAN-ET needs to transmit the current message as a Data Frame, the RTR bit remains unchanged. In case of overrun condition, the message received is discarded. Consequently, when a remote frame is causing overrun (UMSR is set) into a Mailbox configured with $ATX = 1/NMC = 0$, the transmission of the corresponding data frame is not carried out.

Important: In order to support automatic answer to remote frame when MBC = B'001 is used and $ATX = 1$ the RTR flag must be programmed to zero to allow data frame to be transmitted.

Note: when a Mailbox is configured to send a remote frame request the DLC used for transmission is the one stored into the Mailbox.

IDE (Identifier Extension bit): Used to distinguish between the standard format and extended format of CAN data frames and remote frames.

Mailbox-0

Note: MBC[1] of MB0 is always "1".

Mailbox-15 to 1

NMC (New Message Control): When this bit is set to '0', the Mailbox of which the RXPR or RFPR bit is already set does not store the new message but maintains the old one and sets the UMSR correspondent bit. When this bit is set to '1', the Mailbox of which the RXPR or RFPR bit is already set overwrites with the new message and sets the UMSR correspondent bit.

Important: Please note that if a remote frame is overwritten with a data frame or vice versa could be that both RXPR and RFPR flags (together with UMSR) are set for the same Mailbox. In this case the RTR bit within the Mailbox Control Field should be relied upon.

ATX (Automatic Transmission of Data Frame): When this bit is set to '1' and a Remote Frame is received into the Mailbox DLC is stored. Then, a Data Frame is transmitted from the same Mailbox using the current contents of the message data and updated DLC by setting the corresponding TXPR automatically. The scheduling of transmission is still governed by ID priority or Mailbox priority as configured with the Message Transmission Priority control bit (MCR.2). In order to use this function, MBC[2:0] needs to be programmed to be B'001. When a transmission is performed by this function, the DLC (Data Length Code) to be used is the one that has been received. Application needs to guarantee that the DLC of the remote frame correspond to the DLC of the data frame requested.

Important: When ATX is used and MBC = B'001 the filter for the IDE bit cannot be used since ID of remote frame has to be exactly the same as that of data frame as the reply message.

Important: Please note that, when this function is used, the RTR bit will never be set despite receiving a Remote Frame. When a Remote Frame is received, the CPU will be notified by the corresponding RFPR set, however, as RCAN-ET needs to transmit the current message as a Data Frame, the RTR bit remains unchanged.

DART (Disable Automatic Re-Transmission): When this bit is set, it disables the automatic retransmission of a message in the event of an error on the CAN bus or an arbitration lost on the CAN bus. In effect, when this function is used, the corresponding TXCR bit is automatically set at the start of transmission. When this bit is set to '0', RCAN-ET tries to transmit the message as many times as required until it is successfully transmitted or it is cancelled by the TXCR.

MBC[2:0] (Mailbox Configuration): These bits configure the nature of each Mailbox as follows. When $MBC = B'111$, the Mailbox is inactive, i.e., it does not receive or transmit a message regardless of TXPR or other settings. The MBC = B'110, B'101 and B'100 settings are prohibited. When the MBC is set to any other value, the LAFM field becomes available. Please don't set TXPR when MBC is set as reception. Similarly, please don't set TXPR, when MBC is set as remote frame transmission and RTR in Mailbox is cleared. There is no hardware protection, and TXPR remains set. MBC[1] of Mailbox-0 is fixed to "1" by hardware. This is to ensure that MB0 cannot be configured to transmit Messages.

Table 19.4 Mailbox Function Setting

When $ATX = 1$ is used the filter for IDE must not be used

DLC[3:0] (Data Length Code): These bits encode the number of data bytes from 0,1, 2, ... 8 that will be transmitted in a data frame. Please note that when a remote frame request is transmitted the DLC value to be used must be the same as the DLC of the data frame that is requested.

19.3.3 Local Acceptance Filter Mask (LAFM)

This area is used as Local Acceptance Filter Mask (LAFM) for receive boxes**.**

LAFM: When MBC is set to B'001, B'010, B'011, this field is used as LAFM Field. The LAFM is comprised of two 16-bit read/write areas as follows. It allows a Mailbox to accept more than one identifier.

Figure 19.4 Acceptance Filter

If a bit is set in the LAFM, then the corresponding bit of a received CAN identifier is ignored when the RCAN-ET searches a Mailbox with the matching CAN identifier. If the bit is cleared, then the corresponding bit of a received CAN identifier must match to the STDID/IDE/EXTID set in the mailbox to be stored. The structure of the LAFM is same as the message control in a Mailbox. If this function is not required, it must be filled with '0'.

Important: RCAN-ET starts to find a matching identifier from Mailbox-15 down to Mailbox-0. As soon as RCAN-ET finds one matching, it stops the search. The message will be stored or not depending on the NMC and RXPR/RFPR flags. This means that, even using LAFM, a received message can only be stored into 1 Mailbox.

Important: When a message is received and a matching Mailbox is found, the whole message is stored into the Mailbox. This means that, if the LAFM is used, the STDID, RTR, IDE and EXTID may differ to the ones originally set as they are updated with the STDID, RTR, IDE and EXTID of the received message.

STD_LAFM[10:0] — Filter mask bits for the CAN base identifier [10:0] bits.

EXT_LAFM[17:0] — Filter mask bits for the CAN Extended identifier [17:0] bits.

IDE_LAFM — Filter mask bit for the CAN IDE bit.

19.3.4 Message Data Fields

Storage for the CAN message data that is transmitted or received. MSG_DATA[0] corresponds to the first data byte that is transmitted or received. The bit order on the CAN bus is bit 7 through to bit 0.

19.4 RCAN-ET Control Registers

The following sections describe RCAN-ET control registers. The address is mapped as follow.

Important: These registers can only be accessed in Word size (16-bit).

Table 19.5 RCAN-ET Control Registers Configuration

19.4.1 Master Control Register (MCR)

The Master Control Register (MCR) is a 16-bit read/write register that controls RCAN-ET.

• MCR (Address $=$ H'000)

Bit 15 — ID Reorder (MCR15): This bit changes the order of STDID, RTR, IDE and EXTID of both message control and LAFM.

Figure 19.5 ID Reorder

This bit can be modified only in reset mode.

Bit 14 — Auto Halt Bus Off (MCR14): If both this bit and MCR6 are set, MCR1 is automatically set as soon as RCAN-ET enters BusOff.

This bit can be modified only in reset mode.

Bit 13 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 12 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 11 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 10 - 8 — Test Mode (TST[2:0]): This bit enables/disables the test modes. Please note that before activating the Test Mode it is requested to move RCAN-ET into Halt mode or Reset mode. This is to avoid that the transition to Test Mode could affect a transmission/reception in progress. For details, please refer to section 19.6.2, Test Mode Settings.

Please note that the test modes are allowed only for diagnosis and tests and not when RCAN-ET is used in normal operation.

Bit 7 — Auto-wake Mode (MCR7): MCR7 enables or disables the Auto-wake mode. If this bit is set, the RCAN-ET automatically cancels the sleep mode (MCR5) by detecting CAN bus activity (dominant bit). If MCR7 is cleared the RCAN-ET does not automatically cancel the sleep mode.

RCAN-ET cannot store the message that wakes it up.

Note: MCR7 cannot be modified while in sleep mode.

Bit 6 — Halt during Bus Off (MCR6): MCR6 enables or disables entering Halt mode immediately when MCR1 is set during Bus Off. This bit can be modified only in Reset or Halt mode. Please note that when Halt is entered in Bus Off the CAN engine is also recovering immediately to Error Active mode.

Bit 5 — Sleep Mode (MCR5): Enables or disables Sleep mode transition. If this bit is set, while RCAN-ET is in halt mode, the transition to sleep mode is enabled. Setting MCR5 is allowed after entering Halt mode. The two Error Counters (REC, TEC) will remain the same during Sleep mode. This mode will be exited in two ways:

- 1. by writing a '0' to this bit position,
- 2. or, if MCR[7] is enabled, after detecting a dominant bit on the CAN bus.

If Auto wake up mode is disabled, RCAN-ET will ignore all CAN bus activities until the sleep mode is terminated. When leaving this mode the RCAN-ET will synchronise to the CAN bus (by checking for 11 recessive bits) before joining CAN Bus activity. This means that, when the No.2 method is used, RCAN-ET will miss the first message to receive. CAN transceivers stand-by mode will also be unable to cope with the first message when exiting stand by mode, and the S/W needs to be designed in this manner.

In sleep mode only the following registers can be accessed: MCR, GSR, IRR and IMR.

Important: RCAN-ET is required to be in Halt mode before requesting to enter in Sleep mode. That allows the CPU to clear all pending interrupts before entering sleep mode. Once all interrupts are cleared RCAN-ET must leave the Halt mode and enter Sleep mode simultaneously (by writing $MCR[5] = 1$ and $MCR[1] = 0$ at the same time).

Bit 4 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 3 — Reserved. The written value should always be '0' and the returned value is '0'.

Bit 2 — Message Transmission Priority (MCR2): MCR2 selects the order of transmission for pending transmit data. If this bit is set, pending transmit data are sent in order of the bit position in the Transmission Pending Register (TXPR). The order of transmission starts from Mailbox-15 as the highest priority, and then down to Mailbox-1 (if those mailboxes are configured for transmission).

If MCR2 is cleared, all messages for transmission are queued with respect to their priority (by running internal arbitration). The highest priority message has the Arbitration Field (STDID + IDE bit + EXTID (if IDE = 1) + RTR bit) with the lowest digital value and is transmitted first. The internal arbitration includes the RTR bit and the IDE bit (internal arbitration works in the same way as the arbitration on the CAN Bus between two CAN nodes starting transmission at the same time).

This bit can be modified only in Reset or Halt mode.

Bit 1—Halt Request (MCR1): Setting the MCR1 bit causes the CAN controller to complete its current operation and then enter Halt mode (where it is cut off from the CAN bus). The RCAN-ET remains in Halt Mode until the MCR1 is cleared. During the Halt mode, the CAN Interface does not join the CAN bus activity and does not store messages or transmit messages. All the user registers (including Mailbox contents and TEC/REC) remain unchanged with the exception of IRR0 and GSR4 which are used to notify the halt status itself. If the CAN bus is in idle or intermission state regardless of MCR6, RCAN-ET will enter Halt Mode within one Bit Time. If MCR6 is set, a halt request during Bus Off will be also processed within one Bit Time. Otherwise the full Bus Off recovery sequence will be performed beforehand. Entering the Halt Mode can be notified by IRR0 and GSR4.

If both MCR14 and MCR6 are set, MCR1 is automatically set as soon as RCAN-ET enters BusOff.

In the Halt mode, the RCAN-ET configuration can be modified with the exception of the Bit Timing setting, as it does not join the bus activity. MCR[1] has to be cleared by writing a '0' in order to re-join the CAN bus. After this bit has been cleared, RCAN-ET waits until it detects 11 recessive bits, and then joins the CAN bus.

- Note: After issuing a Halt request the CPU is not allowed to set TXPR or TXCR or clear MCR1 until the transition to Halt mode is completed (notified by IRR0 and GSR4). After MCR1 is set this can be cleared only after entering Halt mode or through a reset operation (SW or HW).
- Note: Transition into or recovery from Halt mode, is only possible if the BCR1 and BCR0 registers are configured to a proper Baud Rate.

Bit 0 — Reset Request (MCR0): Controls resetting of the RCAN-ET module. When this bit is changed from '0' to '1' the RCAN-ET controller enters its reset routine, re-initialising the internal logic, which then sets GSR3 and IRR0 to notify the reset mode. During a re-initialisation, all user registers are initialised.

RCAN-ET can be re-configured while this bit is set. This bit has to be cleared by writing a '0' to join the CAN bus. After this bit is cleared, the RCAN-ET module waits until it detects 11 recessive bits, and then joins the CAN bus. The Baud Rate needs to be set up to a proper value in order to sample the value on the CAN Bus.

After Power On Reset, this bit and GSR3 are always set. This means that a reset request has been made and RCAN-ET needs to be configured.

The Reset Request is equivalent to a Power On Reset but controlled by Software.

19.4.2 General Status Register (GSR)

The General Status Register (GSR) is a 16-bit read-only register that indicates the status of RCAN-ET.

 GSR (Address = $H'002$)

Bits 15 to 6: Reserved. The written value should always be '0' and the returned value is '0'.

Bit 5 — Error Passive Status Bit (GSR5): Indicates whether the CAN Interface is in Error Passive or not. This bit will be set high as soon as the RCAN-ET enters the Error Passive state and is cleared when the module enters again the Error Active state (this means the GSR5 will stay high during Error Passive and during Bus Off). Consequently to find out the correct state both GSR5 and GSR0 must be considered.

Bit 4 — Halt/Sleep Status Bit (GSR4): Indicates whether the CAN engine is in the halt/sleep state or not. Please note that the clearing time of this flag is not the same as the setting time of IRR12.

Please note that this flag reflects the status of the CAN engine and not of the full RCAN-ET IP. RCAN-ET exits sleep mode and can be accessed once MCR5 is cleared. The CAN engine exits sleep mode only after two additional transmission clocks on the CAN Bus.

Bit 3— Reset Status Bit (GSR3): Indicates whether the RCAN-ET is in the reset state or not.

Bit 2 — Message Transmission in progress Flag (GSR2): Flag that indicates to the CPU if the RCAN-ET is in Bus Off or transmitting a message or an error/overload flag due to error detected during transmission. The timing to set TXACK is different from the time to clear GSR2. TXACK is set at the $7th$ bit of End Of Frame. GSR2 is set at the $3rd$ bit of intermission if there are no more messages ready to be transmitted. It is also set by arbitration lost, bus idle, reception, reset or halt transition.

Bit 1—Transmit/Receive Warning Flag (GSR1): Flag that indicates an error warning.

Note: REC is incremented during Bus Off to count the recurrences of 11 recessive bits as requested by the Bus Off recovery sequence. However the flag GSR1 is not set in Bus Off.

Bit 0—Bus Off Flag (GSR0): Flag that indicates that RCAN-ET is in the bus off state.

Note: Only the lower 8 bits of TEC are accessible from the user interface. The $9th$ bit is equivalent to GSR0.

19.4.3 Bit Configuration Register (BCR0, BCR1)

The bit configuration registers (BCR0 and BCR1) are 2×16 -bit read/write register that are used to set CAN bit timing parameters and the baud rate pre-scaler for the CAN Interface.

The Time quanta is defined as:

$$
Timequanta = \frac{2 \times BRP}{f_{clk}}
$$

Where: BRP (Baud Rate Pre-scaler) is the value stored in BCR0 incremented by 1 and fclk is the used peripheral clock frequency.

 $BCR1$ (Address = $H'004$)

Please refer to the table below for TSG1 and TSG2 setting.

Bits 15 to 12 — Time Segment 1 (TSG1[3:0] = BCR1[15:12]): These bits are used to set the segment TSEG1 (= PRSEG + PHSEG1) to compensate for edges on the CAN Bus with a positive phase error. A value from 4 to 16 time quanta can be set.

Bit 15: Bit 14: Bit 13: TSG1[3] TSG1[2] TSG1[1] Bit 12:

Bit 11: Reserved. The written value should always be '0' and the returned value is '0'.

Bits 10 to 8 — Time Segment 2 (TSG2[2:0] = BCR1[10:8]): These bits are used to set the segment $TSEG2$ (= PHSEG2) to compensate for edges on the CAN Bus with a negative phase error. A value from 2 to 8 time quanta can be set as shown below.

Bits 7 and 6: Reserved. The written value should always be '0' and the returned value is '0'.

Bits 5 and 4 - ReSynchronisation Jump Width (SJW[1:0] = BCR0[5:4]): These bits set the synchronisation jump width.

Bits 3 to 1: Reserved. The written value should always be '0' and the returned value is '0'.

\bullet BCR0 (Address = H'006)

Bits 8 to 15 : Reserved. The written value should always be '0' and the returned value is '0'.

Bits 7 to 0—Baud Rate Pre-scale (BRP[7:0] = BCR0 [7:0]): These bits are used to define the peripheral clock periods contained in a Time Quantum.

• Requirements of Bit Configuration Register

SYNC SEG: Segment for establishing synchronisation of nodes on the CAN bus. (Normal bit edge transitions occur in this segment.)

- PRSEG: Segment for compensating for physical delay between networks.
- PHSEG1: Buffer segment for correcting phase drift (positive). (This segment is extended when synchronisation (resynchronisation) is established.)
- PHSEG2: Buffer segment for correcting phase drift (negative). (This segment is shortened when synchronisation (resynchronisation) is established)
- $TSEG1: TSG1 + 1$
- $TSEG2 \cdot TSG2 + 1$

The RCAN-ET Bit Rate Calculation is:

Bit Rate = $\frac{\text{fclk}}{2 \times (\text{BRP} + 1) \times (\text{TSEG1} + \text{TSEG2} + 1)}$

where BRP is given by the register value and TSEG1 and TSEG2 are derived values from TSG1 and TSG2 register values.

 $f_{_{\mathrm{CLK}}}$ = Peripheral Clock

BCR Setting Constraints

TSEG1min > TSEG2 \geq SJWmax (SJW = 1 to 4) $8 \leq$ TSEG1 + TSEG2 + 1 \leq 25 time quanta (TSEG1 + TSEG2 + 1 = 7 is not allowed) $TSFG2 > 2$

These constraints allow the setting range shown in the table below for TSEG1 and TSEG2 in the Bit Configuration Register. The number in the table shows possible setting of SJW. "No" shows that there is no allowed combination of TSEG1 and TSEG2.

Example 1: To have a Bit rate of 500 Kbps with a frequency of fclk = 40 MHz it is possible to set: $BRP = 3$, $TSEG1 = 6$, $TSEG2 = 3$.

Then the configuration to write is $BCR1 = H'5200$ and $BCR0 = H'0003$.

Example 2: To have a Bit rate of 250 Kbps with a frequency of fclk = 35 MHz it is possible to set: $BRP = 4$, $TSEG1 = 8$, $TSEG2 = 5$.

Then the configuration to write is $BCR1 = H'7400$ and $BCR0 = H'0004$.

19.4.4 Interrupt Request Register (IRR)

The interrupt request register (IRR) is a 16-bit read/write-clearable register containing status flags for the various interrupt sources.

• IRR (Address = $H'008$)

Bits 15 to 14: Reserved.

Bit 13 - Message Error Interrupt (IRR13): This interrupt indicates that:

- A message error has occurred when in test mode.
- Note: If a Message Overload condition occurs when in Test Mode, then this bit will not be set. When not in test mode this interrupt is inactive.

Bit 12 – Bus Activity while in Sleep Mode (IRR12): IRR12 indicates that a CAN bus activity is present. While the RCAN-ET is in sleep mode and a dominant bit is detected on the CAN bus, this bit is set. This interrupt is cleared by writing a '1' to this bit position. Writing a '0' has no effect. If auto wakeup is not used and this interrupt is not requested it needs to be disabled by the related interrupt mask register. If auto wake up is not used and this interrupt is requested it should be cleared only after recovering from sleep mode. This is to avoid that a new falling edge of the reception line causes the interrupt to get set again.

Please note that the setting time of this interrupt is different from the clearing time of GSR4.

Bits 11 to 10: Reserved

Bit 9 – Message Overrun/Overwrite Interrupt Flag (IRR9): Flag indicating that a message has been received but the existing message in the matching Mailbox has not been read as the corresponding RXPR or RFPR is already set to '1' and not yet cleared by the CPU. The received message is either abandoned (overrun) or overwritten dependant upon the NMC (New Message Control) bit. This bit is cleared when all bit in UMSR (Unread Message Status Register) are cleared (by writing '1') or by setting MBIMR (MailBox interrupt Mast Register) for all UMSR flag set . It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit position has no effect.

Bit 8 - Mailbox Empty Interrupt Flag (IRR8): This bit is set when one of the messages set for transmission has been successfully sent (corresponding TXACK flag is set) or has been successfully aborted (corresponding ABACK flag is set). The related TXPR is also cleared and this mailbox is now ready to accept a new message data for the next transmission. In effect, this bit is set by an OR'ed signal of the TXACK and ABACK bits not masked by the corresponding MBIMR flag. Therefore, this bit is automatically cleared when all the TXACK and ABACK bits are cleared. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit position has no effect.

Bit 7 - Overload Frame (IRR7): Flag indicating that the RCAN-ET has detected a condition that should initiate the transmission of an overload frame. Note that on the condition of transmission being prevented, such as listen only mode, an Overload Frame will NOT be transmitted, but IRR7 will still be set. IRR7 remains asserted until reset by writing a '1' to this bit position - writing a '0' has no effect.

Bit 6 - Bus Off Interrupt Flag (IRR6): This bit is set when RCAN-ET enters the Bus-off state or when RCAN-ET leaves Bus-off and returns to Error-Active. The cause therefore is the existing condition TEC \geq 256 at the node or the end of the Bus-off recovery sequence (128 \times 11 consecutive recessive bits) or the transition from Bus Off to Halt (automatic or manual). This bit remains set even if the RCAN-ET node leaves the bus-off condition, and needs to be explicitly cleared by S/W. The S/W is expected to read the GSR0 to judge whether RCAN-ET is in the busoff or error active status. It is cleared by writing a '1' to this bit position even if the node is still bus-off. Writing a '0' has no effect.

Bit 5 - Error Passive Interrupt Flag (IRR5): Interrupt flag indicating the error passive state caused by the transmit or receive error counter or by Error Passive forced by test mode. This bit is reset by writing a '1' to this bit position, writing a '0' has no effect. If this bit is cleared the node may still be error passive. Please note that the SW needs to check GSR0 and GSR5 to judge whether RCAN-ET is in Error Passive or Bus Off status.

Bit 4 - Receive Error Counter Warning Interrupt Flag (IRR4): This bit becomes set if the receive error counter (REC) reaches a value greater than 95 when RCAN-ET is not in the Bus Off status. The interrupt is reset by writing a '1' to this bit position, writing '0' has no effect.

Bit 3 - Transmit Error Counter Warning Interrupt Flag (IRR3): This bit becomes set if the transmit error counter (TEC) reaches a value greater than 95. The interrupt is reset by writing a '1' to this bit position, writing '0' has no effect.

Bit 2 - Remote Frame Request Interrupt Flag (IRR2): Flag indicating that a remote frame has been received in a mailbox. This bit is set if at least one receive mailbox, with related MBIMR not set, contains a remote frame transmission request. This bit is automatically cleared when all bits in the Remote Frame Receive Pending Register (RFPR), are cleared. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit has no effect.

Bit 1 – Data Frame Received Interrupt Flag (IRR1): IRR1 indicates that there are pending Data Frames received. If this bit is set at least one receive mailbox contains a pending message. This bit is cleared when all bits in the Data Frame Receive Pending Register (RXPR) are cleared, i.e. there is no pending message in any receiving mailbox. It is in effect a logical OR of the RXPR flags from each configured receive mailbox with related MBIMR not set. It is also cleared by writing a '1' to all the correspondent bit position in MBIMR. Writing to this bit has no effect.

Bit 0 – Reset/Halt/Sleep Interrupt Flag (IRR0): This flag can get set for three different reasons. It can indicate that:

- 1. Reset mode has been entered after a SW (MCR0) or HW reset
- 2. Halt mode has been entered after a Halt request (MCR1)
- 3. Sleep mode has been entered after a sleep request (MCR5) has been made while in Halt mode.

The GSR may be read after this bit is set to determine which state RCAN-ET is in.

Important: When a Sleep mode request needs to be made, the Halt mode must be used beforehand. Please refer to the MCR5 description and figure 19.8.

IRR0 is set by the transition from "0" to "1" of GSR3 or GSR4 or by transition from Halt mode to Sleep mode. So, IRR0 is not set if RCAN-ET enters Halt mode again right after exiting from Halt mode, without GSR4 being cleared. Similarly, IRR0 is not set by direct transition from Sleep mode to Halt Request. At the transition from Halt/Sleep mode to Transition/Reception, clearing GSR4 needs (one-bit time - TSEG2) to (one-bit time * 2 - TSEG2).

In the case of Reset mode, IRR0 is set, however, the interrupt to the CPU is not asserted since IMR0 is automatically set by initialisation.

19.4.5 Interrupt Mask Register (IMR)

The interrupt mask register is a 16-bit register that protects all corresponding interrupts in the Interrupt Request Register (IRR) from generating an output signal on the IRQ. An interrupt request is masked if the corresponding bit position is set to '1'. This register can be read or written at any time. The IMR directly controls the generation of IRQ, but does not prevent the setting of the corresponding bit in the IRR.

• IMR (Address $=$ H'00A)

Bit 15 to 0: Maskable interrupt sources corresponding to IRR[15:0] respectively. When a bit is set, the interrupt signal is not generated, although setting the corresponding IRR bit is still performed.

Bit[15:0]: IMRn Description

19.4.6 Transmit Error Counter (TEC) and Receive Error Counter (REC)

The Transmit Error Counter (TEC) and Receive Error Counter (REC) is a 16-bit read/(write) register that functions as a counter indicating the number of transmit/receive message errors on the CAN Interface. The count value is stipulated in the CAN protocol specification Refs. [2], [3], [4] and [5]. When not in (Write Error Counter) test mode this register is read only, and can only be modified by the CAN Interface. This register can be cleared by a Reset request (MCR0) or entering to bus off.

In Write Error Counter test mode (i.e. $TST[2:0] = B'100$), it is possible to write to this register. The same value can only be written to TEC/REC, and the value written into TEC is set to TEC and REC. When writing to this register, RCAN-ET needs to be put into Halt Mode. This feature is only intended for test purposes.

• TEC/REC (Address = H'00C)

Note: * It is only possible to write the value in test mode when TST[2:0] in MCR is B'100. REC is incremented during Bus Off to count the recurrences of 11 recessive bits as requested by the Bus Off recovery sequence.

19.5 RCAN-ET Mailbox Registers

The following sections describe RCAN-ET Mailbox registers that control/flag individual Mailboxes. The address is mapped as follows.

Important: LongWord access is carried out as two consecutive Word accesses.

Table 19.7 RCAN-ET Mailbox Registers

Description	Address	Name	Access Size (bits)
Transmit Pending 1	H'020	TXPR1	LW
Transmit Pending 0	H'022	TXPR0	
	H'024		
	H'026		
	H'028		
Transmit Cancel 0	H'02A	TXCR ₀	
	H'02C		
	H'02E		
	H'030		
Transmit Acknowledge 0	H'032	TXACK0	Word
	H'034		
	H'036		
	H'038		
Abort Acknowledge 0	H'03A	ABACK0	Word
	H'03C		
	H'03E		
	H'040		
Data Frame Receive Pending 0	H'042	RXPR ₀	Word
	H'044		
	H'046		
	H'048		
Remote Frame Receive Pending 0	H'04A	RFPR ₀	Word
	H'04C		
	H'04E		
	H'050		

19.5.1 Transmit Pending Register (TXPR0, TXPR1)

The concatenation of TXPR0 and TXPR1 is a 32-bit register that contains any transmit pending flags for the CAN module. In the case of 16-bit bus interface, Long Word access is carried out as two consecutive word accesses.

<Longword Write Operation>

<Longword Read Operation>

The TXPR1 register cannot be modified and it is always fixed to '0'. The TXPR0 controls Mailbox-15 to Mailbox-1. The CPU may set the TXPR bits to affect any message being considered for transmission by writing a '1' to the corresponding bit location. Writing a '0' has no effect, and TXPR cannot be cleared by writing a '0' and must be cleared by setting the corresponding TXCR bits. TXPR may be read by the CPU to determine which, if any, transmissions are pending or in progress. In effect there is a transmit pending bit for all Mailboxes except for the Mailbox-0. Writing a '1' to a bit location when the mailbox is not configured to transmit is not allowed.

The RCAN-ET will clear a transmit pending flag after successful transmission of its corresponding message or when a transmission abort is requested successfully from the TXCR. The TXPR flag is not cleared if the message is not transmitted due to the CAN node losing the arbitration process or due to errors on the CAN bus, and RCAN-ET automatically tries to transmit it again unless its DART bit (Disable Automatic Re-Transmission) is set in the Message-Control of the corresponding Mailbox. In such case (DART set), the transmission is cleared and notified through Mailbox Empty Interrupt Flag (IRR8) and the correspondent bit within the Abort Acknowledgement Register (ABACK).

If the status of the TXPR changes, the RCAN-ET shall ensure that in the identifier priority scheme $(MCR2 = 0)$, the highest priority message is always presented for transmission in an intelligent way even under circumstances such as bus arbitration losses or errors on the CAN bus. Please refer to section 19.6, Application Note, for details.

When the RCAN-ET changes the state of any TXPR bit position to a '0', an empty slot interrupt (IRR8) may be generated. This indicates that either a successful or an aborted mailbox transmission has just been made. If a message transmission is successful it is signalled in the TXACK register, and if a message transmission abortion is successful it is signalled in the ABACK register. By checking these registers, the contents of the Message of the corresponding Mailbox may be modified to prepare for the next transmission.

• TXPR1

Read value is always H'0000. Long word access is mandatory when reading or writing TXPR1/TXPR0. Writing any value to TXPR1 is allowed, however, write operation to TXPR1 has no effect. Writing to the bit 0 in TXPR0 has no effect.

• TXPR0

Note: * it is possible only to write a '1' for a Mailbox configured as transmitter.

Bit 15 to 1 — Indicates that the corresponding Mailbox is requested to transmit a CAN Frame. The bit 15 to 1 corresponds to Mailbox-15 to 1 respectively. When multiple bits are set, the order of the transmissions is governed by the MCR2 – CAN-ID or Mailbox number.

Bit 0— Reserved: This bit is always '0' as this is a receive-only Mailbox. Writing a '1' to this bit position has no effect. The returned value is '0'.

19.5.2 Transmit Cancel Register 0 (TXCR0)

TXCR0 is a 16-bit read/conditionally-write registers. The TXCR0 controls Mailbox-15 to Mailbox-1.This register is used by the CPU to request the pending transmission requests in the TXPR to be cancelled. To clear the corresponding bit in the TXPR the CPU must write a '1' to the bit position in the TXCR. Writing a '0' has no effect.

When an abort has succeeded the CAN controller clears the corresponding TXPR + TXCR bits, and sets the corresponding ABACK bit. However, once a Mailbox has started a transmission, it cannot be cancelled by this bit. In such a case, if the transmission finishes in success, the CAN controller clears the corresponding TXPR + TXCR bit, and sets the corresponding TXACK bit, however, if the transmission fails due to a bus arbitration loss or an error on the bus, the CAN controller clears the corresponding TXPR + TXCR bit, and sets the corresponding ABACK bit. If an attempt is made by the CPU to clear a mailbox transmission that is not transmit-pending it has no effect. In this case the CPU will be not able at all to set the TXCR flag.

• TXCR0

Note: * Only writing a '1' to a Mailbox that is requested for transmission and is configured as transmit.

Bit 15 to 1 — Requests the corresponding Mailbox, that is in the queue for transmission, to cancel its transmission. The bit 15 to 1 corresponds to Mailbox-15 to 1 (and TXPR0[15:1]) respectively.

Bit 0 — This bit is always '0' as this is a receive-only mailbox. Writing a '1' to this bit position has no effect and always read back as a '0'.

19.5.3 Transmit Acknowledge Register 0 (TXACK0)

The TXACK0 is a 16-bit read/conditionally-write registers. This register is used to signal to the CPU that a mailbox transmission has been successfully made. When a transmission has succeeded the RCAN-ET sets the corresponding bit in the TXACK register. The CPU may clear a TXACK bit by writing a '1' to the corresponding bit location. Writing a '0' has no effect.

• TXACK0

Note: $*$ Only when writing a '1' to clear.

Bit 15 to 1 — Notifies that the requested transmission of the corresponding Mailbox has been finished successfully. The bit 15 to 1 corresponds to Mailbox-15 to 1 respectively.

Bit[15:1]:TXACK0 Description

Bit 0 — This bit is always '0' as this is a receive-only mailbox. Writing a '1' to this bit position has no effect and always read back as a '0'.

19.5.4 Abort Acknowledge Register 0 (ABACK0)

The ABACK0 is a 16-bit read/conditionally-write registers. This register is used to signal to the CPU that a mailbox transmission has been aborted as per its request. When an abort has succeeded the RCAN-ET sets the corresponding bit in the ABACK register. The CPU may clear the Abort Acknowledge bit by writing a '1' to the corresponding bit location. Writing a '0' has no effect. An ABACK bit position is set by the RCAN-ET to acknowledge that a TXPR bit has been cleared by the corresponding TXCR bit.

• ABACK0

Note: $*$ Only when writing a '1' to clear.

Bit 15 to 1 — Notifies that the requested transmission cancellation of the corresponding Mailbox has been performed successfully. The bit 15 to 1 corresponds to Mailbox-15 to 1 respectively.

Bit 0 — This bit is always '0' as this is a receive-only mailbox. Writing a '1' to this bit position has no effect and always read back as a '0'.

19.5.5 Data Frame Receive Pending Register 0 (RXPR0)

The RXPR0 is a 16-bit read/conditionally-write registers. The RXPR is a register that contains the received Data Frames pending flags associated with the configured Receive Mailboxes. When a CAN Data Frame is successfully stored in a receive mailbox the corresponding bit is set in the RXPR. The bit may be cleared by writing a '1' to the corresponding bit position. Writing a '0' has no effect. However, the bit may only be set if the mailbox is configured by its MBC (Mailbox Configuration) to receive Data Frames. When a RXPR bit is set, it also sets IRR1 (Data Frame Received Interrupt Flag) if its MBIMR (Mailbox Interrupt Mask Register) is not set, and the interrupt signal is generated if IMR1 is not set. Please note that these bits are only set by receiving Data Frames and not by receiving Remote frames.

• RXPR0

Note: * Only when writing a '1' to clear.

Bit 15 to 0 — Configurable receive mailbox locations corresponding to each mailbox position from 15 to 0 respectively.

19.5.6 Remote Frame Receive Pending Register 0 (RFPR0)

The RFPR0 is a 16-bit read/conditionally-write registers. The RFPR is a register that contains the received Remote Frame pending flags associated with the configured Receive Mailboxes. When a CAN Remote Frame is successfully stored in a receive mailbox the corresponding bit is set in the RFPR. The bit may be cleared by writing a '1' to the corresponding bit position. Writing a '0' has no effect. In effect there is a bit position for all mailboxes. However, the bit may only be set if the mailbox is configured by its MBC (Mailbox Configuration) to receive Remote Frames. When a RFPR bit is set, it also sets IRR2 (Remote Frame Request Interrupt Flag) if its MBIMR (Mailbox Interrupt Mask Register) is not set, and the interrupt signal is generated if IMR2 is not set. Please note that these bits are only set by receiving Remote Frames and not by receiving Data frames.

• RFPR0

Note: * Only when writing a '1' to clear.

Bit 15 to 0 — Remote Request pending flags for mailboxes 15 to 0 respectively.

19.5.7 Mailbox Interrupt Mask Register 0 (MBIMR0)

The MBIMR1 and MBIMR0 are 16-bit read/write registers. The MBIMR only prevents the setting of IRR related to the Mailbox activities, that are IRR[1] – Data Frame Received Interrupt, IRR[2] – Remote Frame Request Interrupt, IRR[8] – Mailbox Empty Interrupt, and IRR[9] – Message OverRun/OverWrite Interrupt. If a mailbox is configured as receive, a mask at the corresponding bit position prevents the generation of a receive interrupt (IRR[1] and IRR[2] and IRR[9]) but does not prevent the setting of the corresponding bit in the RXPR or RFPR or UMSR. Similarly when a mailbox has been configured for transmission, a mask prevents the generation of an Interrupt signal and setting of an Mailbox Empty Interrupt due to successful transmission or abortion of transmission (IRR[8]), however, it does not prevent the RCAN-ET from clearing the corresponding TXPR/TXCR bit + setting the TXACK bit for successful transmission, and it does not prevent the RCAN-ET from clearing the corresponding TXPR/TXCR bit + setting the ABACK bit for abortion of the transmission.

A mask is set by writing a '1' to the corresponding bit position for the mailbox activity to be masked. At reset all mailbox interrupts are masked.

• MBIMR0

Bit 15 to 0 — Enable or disable interrupt requests from individual Mailbox-15 to Mailbox-0 respectively.

Bit[15:0]: MBIMR0 Description

19.5.8 Unread Message Status Register 0 (UMSR0)

This register is a 16-bit read/conditionally write register and it records the mailboxes whose contents have not been accessed by the CPU prior to a new message being received. If the CPU has not cleared the corresponding bit in the RXPR or RFPR when a new message for that mailbox is received, the corresponding UMSR bit is set to '1'. This bit may be cleared by writing a '1' to the corresponding bit location in the UMSR. Writing a '0' has no effect.

If a mailbox is configured as transmit box, the corresponding UMSR will not be set.

• UMSR0

Bit 15 to 0 — Indicate that an unread received message has been overwritten or overrun condition has occurred for Mailboxes 15 to 0.

19.6 Application Note

19.6.1 Configuration of RCAN-ET

RCAN-ET is considered in configuration mode or after a H/W (Power On Reset)/ S/W (MCR[0]) reset or when in Halt mode. In both conditions RCAN-ET cannot join the CAN Bus activity and configuration changes have no impact on the traffic on the CAN Bus.

(1) After a reset request

The following sequence must be implemented to configure the RCAN-ET after (S/W or H/W) reset. After reset, all the registers are initialized, therefore, RCAN-ET needs to be configured before joining the CAN bus activity. Please read the notes carefully.

If there is a TXPR(s) set, RCAN-ET will start transmission of the message and will be arbitrated by the CAN bus. If it loses the arbitration, it will become a receiver.

(2) Halt mode

When RCAN-ET is in Halt mode, it cannot take part to the CAN bus activity. Consequently the user can modify all the requested registers without influencing existing traffic on the CAN Bus. It is important for this that the user waits for the RCAN-ET to be in halt mode before to modify the requested registers - note that the transition to Halt Mode is not always immediate (transition will occurs when the CAN Bus is idle or in intermission). After RCAN-ET transit to Halt Mode, GSR4 is set.

Once the configuration is completed the Halt request needs to be released. RCAN-ET will join CAN Bus activity after the detection of 11 recessive bits on the CAN Bus.

(3) Sleep mode

When RCAN-ET is in sleep mode the clock for the main blocks of the IP is stopped in order to reduce power consumption. Only the following user registers are clocked and can be accessed: MCR, GSR, IRR and IMR. Interrupt related to transmission (TXACK and ABACK) and reception (RXPR and RFPR) cannot be cleared when in sleep mode (as TXACK, ABACK, RXPR and RFPR are not accessible) and must to be cleared beforehand.

The following diagram shows the flow to follow to move RCAN-ET into sleep mode.

(4) CAN sleep mode

Figure 19.7 Halt Mode/Sleep Mode

Figure 19.8 - Halt Mode/Sleep Mode shows allowed state transition.

- Don't set MCR5 (Sleep Mode) without entering Halt Mode.
- After setting MCR1, make sure that GSR4 is set and the RCAN-ET has entered Halt Mode before clearing MCR1.

Figure 19.8 Halt Mode/Sleep Mode

The following table shows conditions to access registers.

Table 19.8 Conditions to Access Registers

Notes: 1. No hardware protection

2. When TXPR is not set.

19.6.2 Test Mode Settings

The RCAN-ET has various test modes. The register TST[2:0] (MCR[10:8]) is used to select the RCAN-ET test mode. The default (initialized) settings allow RCAN-ET to operate in Normal mode. The following table is examples for test modes.

Test Mode can be selected only while in configuration mode. The user must then exit the configuration mode (ensuring BCR0/BCR1 is set) in order to run the selected test mode.

Table 19.9 Test Mode Settings

• Normal Mode

RCAN-ET operates in the normal mode.

Listen-Only Mode:

ISO-11898 requires this mode for baud rate detection. The Error Counters are cleared and disabled so that the TEC/REC does not increase the values, and the CTx Output is disabled so that RCAN-ET does not generate error frames or acknowledgment bits. IRR13 is set when a message error occurs.

• Self Test Mode 1

RCAN-ET generates its own Acknowledge bit, and can store its own messages into a reception mailbox (if required). The CRx/CTx pins must be connected to the CAN bus.

• Self Test Mode 2

RCAN-ET generates its own Acknowledge bit, and can store its own messages into a reception mailbox (if required). The CRx/CTx pins do not need to be connected to the CAN bus or any external devices, as the internal CTx is looped back to the internal CRx. CTx pin outputs only recessive bits and CRx pin is disabled.

• Write Error Counter

TEC/REC can be written in this mode. RCAN-ET can be forced to become an Error Passive mode by writing a value greater than 127 into the Error Counters. The value written into TEC is used to write into REC, so only the same value can be set to these registers. Similarly, RCAN-ET can be forced to become an Error Warning by writing a value greater than 95 into them.

• Error Passive mode

RCAN-ET needs to be in Halt Mode when writing into TEC/REC (MCR1 must be "1" when writing to the Error Counter). Furthermore this test mode needs to be exited prior to leaving Halt mode.Error Passive Mode: RCAN-ET can be forced to enter Error Passive mode.

Note: the REC will not be modified by implementing this Mode. However, once running in Error Passive Mode, the REC will increase normally should errors be received. In this Mode, RCAN-ET will enter BusOff if TEC reaches 256 (Dec). However when this mode is used RCAN-ET will not be able to become Error Active. Consequently, at the end of the Bus Off recovery sequence, RCAN-ET will move to Error Passive and not to Error Active

When message error occurs, IRR13 is set in all test modes.

19.6.3 Message Transmission Sequence

(1) Message Transmission Request

The following sequence is an example to transmit a CAN frame onto the bus. As described in the previous register section, please note that IRR8 is set when one of the TXACK or ABACK bits is set, meaning one of the Mailboxes has completed its transmission or transmission abortion and is now ready to be updated for the next transmission, whereas, the GSR2 means that there is currently no transmission request made (No TXPR flags set).

Figure 19.9 Transmission Request

(2) Internal Arbitration for Transmission

The following diagram explains how RCAN-ET manages to schedule transmission-requested messages in the correct order based on the CAN identifier. 'Internal arbitration' picks up the highest priority message amongst transmit-requested messages.

Figure 19.10 Internal Arbitration for transmission

The RCAN-ET has two state machines. One is for transmission, and the other is for reception.

- 1-1: When a TXPR bit(s) is set while the CAN bus is idle, the internal arbitration starts running immediately and the transmission is started.
- 1-2: Operations for both transmission and reception starts at SOF. Since there is no reception frame, RCAN-ET becomes transmitter.
- 2-1: At crc delimiter, internal arbitration to search next message transmitted starts.
- 2-2: Operations for both transmission and reception starts at SOF. Because of a reception frame with higher priority, RCAN-ET becomes receiver. Therefore, Reception is carried out instead of transmitting Frame-3.
- 3-1: At crc delimiter, internal arbitration to search next message transmitted starts.
- 3-2: Operations for both transmission and reception starts at SOF. Since a transmission frame has higher priority than reception one, RCAN-ET becomes transmitter.

Internal arbitration for the next transmission is also performed at the beginning of each error delimiter in case of an error is detected on the CAN Bus. It is also performed at the beginning of error delimiters following overload frame.

As the arbitration for transmission is performed at CRC delimiter, in case a remote frame request is received into a Mailbox with $ATX = 1$ the answer can join the arbitration for transmission only at the following Bus Idle, CRC delimiter or Error Delimiter.

Depending on the status of the CAN bus, following the assertion of the TXCR, the corresponding Message abortion can be handled with a delay of maximum 1 CAN Frame.
19.6.4 Message Receive Sequence

The diagram below shows the message receive sequence.

Figure 19.11 Message receive sequence

When RCAN-ET recognises the end of the Arbitration field while receiving a message, it starts comparing the received identifier to the identifiers set in the Mailboxes, starting from Mailbox-15 down to Mailbox-0. It first checks the MBC if it is configured as a receive box, and reads LAFM, and reads the CAN-ID of Mailbox-15 (if configured as receive) to finally compare them to the received ID. If it does not match, the same check takes place at Mailbox-14 (if configured as receive). Once RCAN-ET finds a matching identifier, it stores the number of Mailbox-[n] into an internal buffer, stops the search, and goes back to idle state, waiting for the EndOfFrame (EOF) to come. When the $6th$ bit of EOF is notified by the CAN Interface logic, the received message is written or abandoned, depending on the NMC bit. No modification of configuration during communication is allowed. Entering Halt Mode is one of ways to modify configuration. If it is written into the corresponding Mailbox, including the CAN-ID, i.e., there is a possibility that the CAN-ID is overwritten by a different CAN-ID of the received message due to the LAFM used. This also implies that, if the identifier of a received message matches to $ID + LAFM$ of 2 or more Mailboxes, the higher numbered Mailbox will always store the relevant messages and the lower numbered Mailbox will never receive messages. Therefore, the settings of the identifiers and LAFMs need to be carefully selected.

With regards to the reception of data and remote frames described in the above flow diagram the clearing of the UMSR flag after the reading of IRR is to detect situations where a message is overwritten by a new incoming message stored in the same mailbox while the interrupt service routine is running. If during the final check of UMSR a overwrite condition is detected the message needs to be discarded and read again.

Please note that in the case a received remote frame is overwritten by a data frame, both the remote frame request interrupt (IRR2) and data frame received interrupt (IRR1) and also the Receive Flags (RXPR and RFPR) are set. In an analogous way, the overwriting of a data frame by a remote frame, leads to setting both IRR2 and IRR1.

In the Overrun Mode (NMC = $\dot{0}$), only the first Mailbox will cause the flags to be asserted. So, if a Data Frame is initially received, then RXPR and IRR1 are both asserted. If a Remote Frame is then received before the Data Frame has been read, then RFPR and IRR2 are NOT set. In this case UMSR of the corresponding Mailbox will still be set.

19.6.5 Reconfiguration of Mailbox

When re-configuration of Mailboxes is required, the following procedures should be taken.

(1) Change configuration of transmit box

Two cases are possible.

- Change of ID, RTR, IDE, LAFM, Data, DLC, NMC, ATX, DART This change is possible only when $MBC = B'000$. Confirm that the corresponding TXPR is not set. The configuration (except MBC bit) can be changed at any time.
- Change from transmit to receive configuration (MBC) Confirm that the corresponding TXPR is not set. The configuration can be changed only in Halt or reset state. Please note that it might take longer for RCAN-ET to transit to halt state if it is receiving or transmitting a message (as the transition to the halt state is delayed until the end of the reception/transmission), and also RCAN-ET will not be able to receive/transmit messages during the Halt state.

In case RCAN-ET is in the Bus Off state the transition to halt state depends on the configuration of the bit 6 of MCR and also bit and 14 of MCR.

(2) Change configuration (ID, RTR, IDE, LAFM, Data, DLC, NMC, ATX, DART) of receive box or Change receive box to transmit box

The configuration can be changed only in Halt Mode.

RCAN-ET will not lose a message if the message is currently on the CAN bus and RCAN-ET is a receiver. RCAN-ET will be moving into Halt Mode after completing the current reception. Please note that it might take longer if RCAN-ET is receiving or transmitting a message (as the transition to the halt state is delayed until the end of the reception/transmission), and also RCAN-ET will not be able to receive/transmit messages during the Halt Mode.

In case RCAN-ET is in the Bus Off state the transition to halt mode depends on the configuration of the bit 6 and 14 of MCR.

Figure 19.12 Change ID of Receive Box or Change Receive Box to Transmit Box

19.7 Interrupt Sources

Table 19.10 lists the RCAN-ET interrupt sources. With the exception of the reset processing interrupt (IRR0) by a power-on reset, these sources can be masked. Masking is implemented using the mailbox interrupt mask register 0 (MBIMR0) and interrupt mask register (IMR). For details on the interrupt vector of each interrupt source, see section 6, Interrupt Controller (INTC).

Table 19.10 RCAN-ET Interrupt Sources

Notes: 1. Available only in Test Mode.

 2. RM0 is an interrupt generated by the remote request pending flag for mailbox 0 (RFPR0[0]) or the data frame receive flag for mailbox 0 (RXPR0[0]). RM1 is an interrupt generated by the remote request pending flag for mailbox n (RFPR0[n]) or the data frame receive flag for mailbox n (RXPR0[n]) ($n = 1$ to 15).

3. IRR1 is a data frame received interrupt flag for mailboxes 0 to 15, and IRR2 is a remote frame request interrupt flag for mailboxes 0 to 15.

19.8 CAN Bus Interface

A bus transceiver IC is necessary to connect this LSI to a CAN bus. A Renesas HA13721 transceiver IC and its compatible products are recommended. As the CRx and CTx pins use 3 V, an external level shifter is necessary. Figure 19.13 shows a sample connection diagram.

Figure 19.13 High-Speed Interface Using HA13721

19.9 Usage Notes

19.9.1 Module Standby Mode

The standby control register 2 (STBCR2) controls the supply of clocks to RCAN-ET. As an initial value, the clock to RCAN-ET is halted. Registers should be accessed after the module stop mode is released.

19.9.2 Reset

Two types of resets are supported for RCAN-ET.

• Hardware reset

RCAN-ET is initialized by a power-on reset, deep standby mode, or software standby mode.

• Software reset

The MCR0 bit in the master control register (MCR) initializes registers other than MCR and CAN communication functions.

As the IRR0 bit in the interrupt request register (IRR) is initialized and set to 1 at a reset, it should be cleared to 0 in the configuration mode shown in the reset sequence diagram.

The area except for the message control field 1 (CONTROL1) of Mailbox is consisted of RAM, and not initialized at a reset. After a power-on reset, all the Mailboxes should be initialized in the configuration mode shown in the reset sequence diagram.

19.9.3 CAN Sleep Mode

The supply of main clocks in the modules is stopped in CAN sleep mode. Therefore, registers other than MCR, GSR, IRR, and IMR should not be accessed in CAN sleep mode.

19.9.4 Register Access

When the CAN bus receive frame is being stored in the Mailbox with the CAN communication functions of RCAN-ET, accessing the Mailbox area generates 0 to 5 peripheral bus cycles as a wait.

19.9.5 Interrupts

As shown in table 19.2, the Mailbox 0 receive interrupt enables the DMAC activation. When an interrupt is specified as to be activated by the Mailbox 0 receive interrupt and cleared by the interrupt source at the DMA transfer, up to the message control field 1 (CONTROL1) of Mailbox 0 should be read using the block transfer mode.

Section 20 IEBus[™] Controller (IEB) [R5S72612] [R5S72613]

This LSI has an on-chip one-channel IEBus controller (IEB). The Inter Equipment Bus™ $(IEBus^TTM)^*$ is a small-scale digital data transfer system for inter-equipment data transfer.

This LSI does not have an on-chip IEBus driver/receiver, so it is necessary to mount a dedicated driver/receiver externally. In addition, as the IERxD and IETxD pins need 3V to operate, a dedicated external level shifter is necessary.

Note: $*$ The Inter Equipment Bus™ (IEBus™) is a trademark of Renesas Electronics Corporation.

20.1 Features

- IEBus protocol control (layer 2) supported
	- Half-duplex asynchronous communications
	- Multi-master system
	- ⎯ Broadcast communications function
	- Selectable mode (three types) with different transfer speeds
- On-chip buffers for data transmission and reception
	- Transmission and reception buffers: 128 bytes each
	- ⎯ Up to 128 bytes of consecutive transmit/reception (maximum number of transfer bytes in mode 2)
- Operating frequency
	- $-$ 6 MHz, 6.29 MHz (IEB uses clocks of P ϕ or AUDIO_X1*/AUDIO_X2*.)
	- -12 MHz, 12.58 MHz (IEB uses 1/2 divided clocks of P ϕ or AUDIO X1*/AUDIO X2*.)
	- $-$ 18 MHz, 18.87 MHz (IEB uses 1/3 divided clocks of P ϕ or AUDIO_X1*/AUDIO_X2*.)
	- $-$ 24 MHz, 25.16 MHz (IEB uses 1/4 divided clocks of P ϕ or AUDIO_X1*/AUDIO_X2*.)
	- \sim 30 MHz, 31.45 MHz (IEB uses 1/5 divided clocks of P ϕ or AUDIO_X1*/AUDIO_X2*.)
	- \sim 36 MHz, 37.74 MHz (IEB uses 1/6 divided clocks of P ϕ or AUDIO_X1*/AUDIO_X2*.)

Note: * Available as the IEB clock input only when not used as the clock input for SSI audio

Module standby mode can be set.

20.1.1 IEBus Communications Protocol

An overview of the IEBus is provided below.

- Communications method: Half-duplex asynchronous communications
- Multi-master system

All units connected to the IEBus can transfer data to other units.

- Broadcast communications function (one-to-many communications)
	- ⎯ Group broadcast communications: Broadcast communications to group unit
	- ⎯ General broadcast communications: Broadcast communications to all units
- Mode is selectable (three modes with different transfer speeds)

Table 20.1 Mode Types

Notes: 1. Peripheral clock (Pφ), or clocks for AUDIO_X1 and AUDIO_X2

2. Oscillation frequency when this LSI is used

- Access control: CSMA/CD (Carrier Sense Multiple Access with Collision Detection) Priority of bus mastership is as follows.
	- ⎯ Broadcast communications (one-to-many communications) have priority over normal communications (one-to-one communications).
	- $-$ A smaller master address has priority.
- Communications scale
	- Number of units: Up to 50
	- Cable length: Up to 150 m (when using a twisted-pair cable)
- Note: The communications scale of the actual system depends on the characteristics of the externally mounted IEBus driver/receiver and the cable used.

(1) Determination of Bus Mastership (Arbitration)

A unit connected to the IEBus performs an operation to get the bus to control other units. This operation is called arbitration. In arbitration, when multiple units start transferring simultaneously, the bus mastership is given to one unit among them.

Only one unit can obtain bus mastership through arbitration, so the following priority for bus mastership is determined.

(a) Priority according to communications type

Broadcast communications (one-to-many communications) has priority over normal communications (one-to-one communications).

(b) Priority according to master address

The unit with the smallest master address has priority among units of the same communications type.

Example: The master address is configured with 12 bits. A unit with H'000 has the highest priority, while a unit with H'FFF has the lowest priority.

Note: When a unit loses in arbitration, the unit can automatically enter retransfer mode (0 to 7 retransfer times can be selected by the RN bit in IEMCR).

(2) Communications Mode

The IEBus has three communications modes with different transfer speeds. Table 20.2 shows the transfer speed in each communications mode and the maximum number of transfer bytes in one communications frame.

Table 20.2 Transfer Speed and Maximum Number of Transfer Bytes in Each Communications Mode

Notes: Each unit connected to the IEBus should select a communications mode prior to performing communications. Note that correct communications is not guaranteed if the master and slave units do not adopt the same communications mode. In the case of communications between a unit with $\phi = 6$ MHz and a unit with $\phi = 6.29$

MHz, correct communications are not possible even if the same communications mode is adopted. Communications must be done with the same oscillation frequency.

- 1. Effective transfer speed when the maximum number of transfer bytes is transmitted.
- 2. Peripheral clock (Pφ), or clocks for AUDIO_X1 and AUDIO_X2
- 3. Oscillation frequency when this LSI is used

(3) Communications Address

In the IEBus, a specific 12-bit communications address is allocated to each individual unit. A communications address is configured as follows.

- Upper four bits: group number (number identifying a group to which the unit belongs)
- Lower eight bits: unit number (number identifying individual units in a group)

(4) Broadcast Communications

In normal transfer, a single master unit communicates with a single slave unit, so one-to-one transfer or reception takes place. In broadcast communications, a single master unit communicates with multiple slave units. Since there are multiple slave units, no acknowledgements are returned from the slave units during communications.

A broadcast bit decides whether broadcast or normal communications is done. (For details of the broadcast bit, see section 20.1.2 (1) (b), Broadcast Bit.

There are two types of broadcast communications.

(a) Group broadcast communications

Broadcast communications is aimed at units with the same group number, meaning that those units have the same upper four bits of the communications address.

(b) General broadcast communications

Broadcast communications is aimed at all units regardless of group number.

Group broadcast and general broadcast communications are identified by a slave address. (For details on the slave address, see section 20.1.2 (3), Slave Address Field.)

20.1.2 Communications Protocol

Figure 20.1 shows an IEBus transfer signal format.

Communications data is transferred as a series of signals referred to as a communications frame. The number of data which can be transmitted in a single communications frame and the transfer speed differ according to the communications mode.

Figure 20.1 Transfer Signal Format

(1) Header

A header is comprised of a start bit and a broadcast bit.

(a) Start Bit

The start bit is a signal to inform other units of the start of data transfer. A unit attempting to start data transfer outputs a low-level signal (the start bit) for a specified period and then outputs the broadcast bit.

If another unit is already outputting a start bit when a unit attempts to output a start bit, the unit waits for completion of the start bit from the other unit without outputting its own start bit, and then outputs the broadcast bit synchronized with the completion timing.

Other units enter the receive state after detecting the start bit.

(b) Broadcast Bit

The broadcast bit is a bit to identify the type of communications: broadcast or normal.

When this bit is cleared to 0, it indicates broadcast communications. When it is set to 1, it indicates normal communications. Broadcast communications includes group broadcast and general broadcast, which are identified by a value of the slave address. (For details of the slave address, see section 20.1.2 (3), Slave Address Field.)

Since multiple slave units are communications destination units, in the case of broadcast communications, the acknowledge bit is not returned from each field described in (2) and below.

When more than one unit starts to transfer a communications frame with the same timing, broadcast communications has priority over normal communications, and arbitration occurs.

(2) Master Address Field

The master address field is a field for transmitting the unit address (master address) to other units. The master address field is comprised of master address bits and a parity bit.

The master address consists of 12 bits and the MSB is output first.

When more than one unit start to transfer broadcast bits having the same value with the same timing, arbitration is decided by the master address field.

In the master address field, self-output data and data on the bus are compared for every one-bit transfer. If the self-output master address and data on the bus are different, the unit that loses arbitration will stop its transfer and enter the receive state.

Since the IEBus is configured with wired AND, the unit having the smallest master address of the units in arbitration (arbitration master) wins in arbitration.

Finally, only a single unit remains in the transfer state as a master unit after outputting a 12-bit master address.

Next, this master unit outputs a parity bit*, defines the master address for other units, and then enters the slave address field output state.

Note: * Since even parity is used, when the number of one bit in the master address is odd, the parity bit is 1.

(3) Slave Address Field

The slave address field is a field to transmit an address (the slave address) of a unit (the slave unit) to be transmitted. The slave address field is comprised of slave address bits, a parity bit, and an acknowledge bit.

The slave address consists of 12 bits and the MSB is output first. The parity bit is output after the 12-bit slave address is transmitted to avoid receiving the slave address accidentally. The master unit then detects the acknowledgement from the slave unit to confirm that the slave unit exists on the bus. When the acknowledgement is detected, the master unit enters the control field output state. However, the master unit enters the control field output state without detecting the acknowledgement in broadcast communications.

The slave unit returns an acknowledgement when the slave addresses match and the parities of the master and slave addresses are correct. When the parity of either the master or slave address is incorrect, the slave unit decides that the master or slave address was not correctly received and does not return the acknowledgement. In this case, the master unit enters the waiting (monitor) state and communications ends.

In the case of broadcast communications, the slave address is used to identify the type of broadcast communications (group or general) as follows:

- When the slave address is H'FFF: General broadcast communications
- When the slave address is other than H'FFF: Group broadcast communications
- Note: The group number is the upper 4-bit value of the slave address in group broadcast communications.

(4) Control Field

The control field is a field for transmitting the type and direction of the following data field. The control field is comprised of control bits, a parity bit, and an acknowledge bit.

The control bits consist of four bits and the MSB is output first.

The parity bit is output following the control bits. When the parity is correct, and the slave unit can implement the function required from the master unit, the slave unit returns an acknowledgement and enters the message length field output state. However, if the slave unit cannot implement the requirements from the master unit even though the parity is correct, or if the parity is not correct, the slave unit does not return an acknowledgement and returns to the waiting (monitor) state.

The master unit enters the subsequent message length field output state after confirming the acknowledgement.

When the acknowledgement is not confirmed, the master unit enters the waiting (monitor) state, and communications ends. However, in the case of broadcast communications, the master unit enters the following message length field output state without confirming the acknowledgement. For details of the contents of the control bit, see table 20.4.

(5) Message Length Field

The message length field is a field for specifying the number of transfer bytes. The message length field is comprised of message length bits, a parity bit, and an acknowledge bit.

The message length has eight bits and the MSB is output first. Table 20.3 shows the number of transfer bytes.

Table 20.3 Contents of Message Length bits

Note: If a number greater than the maximum number of transfer bytes in one frame is specified, communications are done in multiple frames depending on the communications mode. In this case, the message length bits indicate the number of remaining communications data after the first transfer. In this LSI, the message length bits must be smaller than the maximum number of transfer bytes in one frame. Set these within the ranges shown below. Mode 0: 1 to 16 bytes Mode 1: 1 to 32 bytes Mode 2: 1 to 128 bytes

This field operation differs depending on the value of bit 3 in the control field: master transmission (the bit 3 of the control bits is 1) or master reception (the bit 3 of the control bits is 0).

(a) Master Transmission

The master unit outputs the message length bits and the parity bit. When the parity is even, the slave unit returns an acknowledgement and enters the following data field. Note that the slave unit does not return an acknowledgement in broadcast communications.

When the parity is odd, the slave unit decides that the message length field is not correctly received, does not return an acknowledgement, and returns to the waiting (monitor) state. In this case, the master unit also returns to the waiting state and communications end.

(b) Master Reception

The slave unit outputs the message length bits and parity bit. When even parity is confirmed, the master unit returns an acknowledgement.

When the parity is not correct, the master unit decides that the message length bits are not correctly received, does not return an acknowledgement, and returns to the waiting state. In this case, the slave unit also returns to the waiting state and communications end.

(6) Data Field

The data field is a field for data transmission/reception to and from the slave unit. The master unit transmits/receives data to and from the slave unit using the data field. The data field is comprised of data bits, a parity bit, and an acknowledge bit.

The data bits consist of eight bits and the MSB is output first.

The parity and acknowledge bits are output following the data bits from the master unit and slave unit, respectively.

Broadcast communications are performed only for the transmission of the master unit. In this case, the acknowledge bit is ignored. Operations in master transmission and master reception are described below.

(a) Master Transmission

The master unit transmits the data bits and parity bit to the slave unit to write data from the master unit to the slave unit. The slave unit receives the data bits and parity bit, and returns an acknowledgement if the parity bit is even and the receive buffer is empty. If the parity bit is odd or the receive buffer is not empty, the slave unit does not accept the corresponding data and does not return an acknowledgement.

When the slave unit does not return an acknowledgement, the master unit retransmits the data. This operation is repeated until either an acknowledgement from the slave unit is detected or the maximum number of data transfer bytes is reached.

When the parity is even and the acknowledgement is output from the slave unit, the master unit transmits the subsequent data if data remains and the maximum number of transfer bytes are not exceeded.

In the case of broadcast communications, the slave unit does not return the acknowledgement, and the master unit transfers data byte by byte.

(b) Master Reception

The master unit outputs synchronous signals corresponding to all data bits to be read from the slave unit.

The slave unit outputs the data bits and parity bit on the bus in accordance with the synchronous signals from the master unit.

The master unit reads the parity bit output from the slave unit, and checks the parity. If the parity is not even, or the receive buffer is not empty, the master unit rejects acceptance of the data, and does not return the acknowledgement. The master unit reads the same data repeatedly if the number of data does not exceed the maximum number of transfer bytes in one frame. If the parity is even and the receive buffer is empty, the master unit accepts data and returns an acknowledgement. The master unit reads in the subsequent data if the number of data does not exceed the maximum number of transfer bytes in one frame.

(7) Parity bit

The parity bit is used to confirm that transfer data occurs with no errors.

The parity bit is added to respective data of the master address, slave address, control, message length, and data bits.

Even parity is used. When the number of bits having the value 1 is odd, the parity bit is 1. When the number of bits having the value 1 is even, the parity bit is 0.

(8) Acknowledge bit

In normal communications (single unit to single unit communications), the acknowledge bit is added in the following positions to confirm that data is correctly accepted.

- At the end of the slave address field
- At the end of the control field
- At the end of the message length field
- At the end of the data field

The acknowledge bit is defined below.

- 0: indicates that the transfer data is acknowledged. (ACK)
- 1: indicates that the transfer data is not acknowledged. (NAK)

Note that the acknowledge bit is ignored in the case of broadcast communications.

(a) Acknowledge bit at the End of the Slave Address Field

The acknowledge bit at the end of the slave address field becomes NAK in the following cases and transfer is stopped.

- When the parity of the master address or slave address bits is incorrect
- When a timing error (an error in bit format) occurs
- When there is no slave unit

(b) Acknowledge bit at the End of the Control Field

The acknowledge bit at the end of the control field becomes NAK in the following cases and transfer is stopped.

- When the parity of the control bits is incorrect
- When the bit 3 of the control bits is 1 (data write) although the slave receive buffer* is not empty
- When the control bits are set to data read $(H³, H⁷)$ although the slave transmit buffer* is empty
- When another unit which locked the slave unit requests H'3, H'6, H'7, H'A, H'B, H'E, or H'F in the control bits although the slave unit has been locked
- When the control bits are the locked address read (H'4, H'5) although the unit is not locked
- When a timing error occurs
- When the control bits are undefined

Note: See section 20.1.3 (1), Slave Status Read (Control Bits: H'0, H'6).

(c) Acknowledge Bit at the End of the Message Length Field

The acknowledge bit at the end of the message length field becomes NAK in the following cases and transfer is stopped.

- When the parity of the message length bits is incorrect
- When a timing error occurs

(d) Acknowledge Bit at the End of the Data Field

The acknowledge bit at the end of the data field becomes NAK in the following cases and transfer is stopped.

- When the parity of the data bits is incorrect^{*}
- When a timing error occurs after the previous transfer of the acknowledge bit
- When the receive buffer becomes full and cannot accept further data*
- Note: * In this case, the data field is transferred repeatedly until the number of data reaches the maximum number of transfer bytes if the number of data does not exceed the maximum number of transfer bytes in one frame.

20.1.3 Transfer Data (Data Field Contents)

The data field contents are specified by the control bits.

Table 20.4 Control Bit Contents

Notes: 1. Depending on the value of bit 3 (MSB), the transfer directions of the message length bits in the following message length field and data in the data field vary. When bit 3 is 1: Data is transferred from the master unit to the slave unit. When bit 3 is 0: Data is transferred from the slave unit to the master unit.

 2. H'3, H'6, H'A, and H'B are control bits to specify lock setting and cancellation. When the undefined values of H'1, H'2, H'8, H'9, H'C, and H'D are transmitted, the acknowledge signal is not returned.

When the control bits received from another unit which locked are not included in table 20.5, the slave unit which has been locked by the master unit does not accept the control bits and does not return the acknowledge bit.

Table 20.5 Control Field for Locked Slave Unit

(1) Slave Status Read (Control Bits: H'0, H'6)

The master unit can decide the reason the slave unit does not return the acknowledgement (ACK) by reading the slave status (H'0, H'6). The slave status indicates the result of the last communications that the slave unit performed. All slave units can provide slave status information. Figure 20.2 shows the bit configuration of the slave status.

Figure 20.2 Bit Configuration of Slave Status (SSR)

(2) Data Command Transfer (Control Bits: Read (H'3, H'7), Write (H'A, H'B, H'E, H'F))

In the case of data read (H'3, H'7), data in the data buffer of the slave unit is read in the master unit. In the case of data write (H'B or H'F) or command write (H'A or H'E), data received in the slave unit is processed in accordance with the operation specification of the slave unit.

Notes: 1. The user can select data and commands freely in accordance with the system.

2. H'3, H'A, or H'B may lock depending on the communications condition and status.

(3) Locked Address Read (Control Bits: H′**4, H**′**5)**

In the case of the locked address read (H'4 or H'5), the address (12 bits) of the master unit which issues the lock instruction is configured in bytes as shown in figure 20.3.

Figure 20.3 Locked Address Configuration

(4) Locking/Unlocking (Control Bits: Setting (H'3, H'A, H'B), Cancellation: (H'6))

The lock function is used for message transfer over multiple communications frames. A locked unit receives data only from the unit which locked it.

Locking and unlocking are described below.

(a) Locking

When an acknowledge bit of 0 in the message length field is transmitted/received with the control bits (H'3, H'A, H'B) indicating the lock operation, and then the communications frame is completed before completion of data transmission/reception for the number of bytes specified by the message length bits, the slave unit is locked by the master unit. In this case, the bit (bit 2) relevant to locking in the byte data indicating the slave status is set to 1.

Lock is set only when the number of data exceeds the maximum number of transfer bytes in one frame. Lock is not set by other error terminations.

(b) Unlocking

When the control bits indicate the lock (H'3, H'A, or H'B) or unlock (H'6) operation and the byte data for the number of bytes specified by the message length bits are transmitted/received in a single communications frame, the slave unit is unlocked by the master unit. In this case, the bit (bit 2) relevant to locking in the byte indicating the slave status is cleared to 0.

Note that locking and unlocking are not done in broadcast communications.

Note: * There are three ways to cause a locked unit to unlock itself.

- Perform a power-on reset
- Put the unit in deep standby mode
- Issue an unlock command through the IEBus command register (IECMR)

 Note that the LCK flag in IEFLG can be used to check whether the unit is locked or unlocked.

20.1.4 Bit Format

Figure 20.4 shows the bit format (conceptual diagram) configuring the IEBus communications frame.

Each period of the bit format for use of active high signals is described below.

- Preparation period: first logic 1 period (high level)
- Synchronous period: subsequent logic 0 period (low level)
- Data period: period indicating bit value (logic 1: high level, logic 0: low level)
- Halt period: last logic 1 period (high level)

For use of active low signals, levels are reversed from the active high signals.

The synchronous and data periods have approximately the same length.

The IEBus is synchronized bit by bit. The specifications for the time of all bits and the periods allocated to the bits differ depending on the type of transfer bits and the unit (master or slave unit).

20.1.5 Configuration

Figure 20.5 shows the entire block configuration and table 20.6 lists the functions of each block.

Figure 20.5 IEB Block Diagram

Table 20.6 Functions of Each Block

20.2 Input/Output Pins

Table 20.7 shows the IEB pin configuration.

Table 20.7 Pin Configuration

20.3 Register Descriptions

The IEB has the following registers.

Each register, in principle, has 8-bit width and is accessed in 8 bits.

Table 20.8 Register Configuration

Note: * Only 1 can be written to clear the flag.

20.3.1 IEBus Control Register (IECTR)

IECTR is used to control the IEB operation.

IECTR is initialized by a power-on reset or in deep standby.

20.3.2 IEBus Command Register (IECMR)

IECMR issues commands to control IEB communications. Since this register is a write-only register, the read value is undefined.

IECMR is initialized by a power-on reset or in deep standby.

Notes: 1. Do not execute this command in slave communications.

- 2. This command is valid during master communications (MRQ $= 1$). In other states, this command issuance is ignored. If this command is issued in master communications, the communications controller immediately enters the wait state. At this time, the issued master transmission request ends $(MRQ = 0)$.
- 3. This command is valid during slave communications (SRQ = 1). In other states, this command issuance is ignored. Once this command is issued in slave transmission, the SRQ flag is 0 before slave transmission. Therefore, a transmit request from the master is not responded to. If a transmit request is issued during slave transmission, the transmission stops and the wait state is entered $(SRQ = 0)$.
- 4. Undefined bits. Issuing this command does not affect operation.

20.3.3 IEBus Master Control Register (IEMCR)

IEMCR sets the communication conditions for master communications.

IEMCR is initialized by a power-on reset or in deep standby.

Notes: 1. CTL3 decides the data transfer direction of the message length bits in the message length field and data bits in the data field:

CTL3 = 1: Transfer is from master unit to slave unit

CTL3 = 0: Transfer is from slave unit to master unit

- 2. Control bits to lock and unlock
- 3. Setting prohibited.

20.3.4 IEBus Master Unit Address Register 1 (IEAR1)

IEAR1 sets the lower four bits of the master unit address and communications mode. In master communications, the master unit address becomes the master address field value. In slave communications, the master unit address is compared with the received slave address field.

IEAR1 is initialized by a power-on reset or in deep standby.

20.3.5 IEBus Master Unit Address Register 2 (IEAR2)

IEAR2 sets the upper eight bits of the master unit address. In master communications, this register becomes the master address field value. In slave communications, this register is compared with the received slave address field.

IEAR2 is initialized by a power-on reset or in deep standby.

20.3.6 IEBus Slave Address Setting Register 1 (IESA1)

IESA1 sets the lower four bits of the communications destination slave unit address.

IESA1 is initialized by a power-on reset or in deep standby.

20.3.7 IEBus Slave Address Setting Register 2 (IESA2)

IESA2 sets the upper eight bits of the communications destination slave unit address.

IESA2 is initialized by a power-on reset or in deep standby.

20.3.8 IEBus Transmit Message Length Register (IETBFL)

IETBFL sets the message length for master or slave transmission.

IETBFL is initialized by a power-on reset or in deep standby.

Note: * Setting prohibited

20.3.9 IEBus Reception Master Address Register 1 (IEMA1)

IEMA1 indicates the lower four bits of the communication destination master unit address in slave/broadcast reception.

IEMA1 is initialized by a power-on reset or in deep standby.

20.3.10 IEBus Reception Master Address Register 2 (IEMA2)

IEMA2 indicates the upper eight bits of the communications destination master unit address in slave/broadcast reception. This register is enabled when slave/broadcast reception starts, and the contents are changed at the time of setting the RXS flag in IERSR.

If a broadcast receive error interrupt is selected with the DEE bit in IECTR and the receive buffer is not in the receive enabled state at control field reception, a receive error interrupt is generated and the upper eight bits of the master address are stored in IEMA2. This register cannot be modified by a write.

IEMA2 is initialized by a power-on reset or in deep standby.

20.3.11 IEBus Receive Control Field Register (IERCTL)

IERCTL indicates the control field value in slave/broadcast reception. This register is enabled when slave/broadcast receive starts, and the contents are changed at the time of setting the RXS flag in IERSR. This register cannot be modified.

IERCTL is initialized by a power-on reset or in deep standby.

20.3.12 IEBus Receive Message Length Register (IERBFL)

IERBFL indicates the message length field in slave/broadcast reception. This register is enabled when slave/broadcast receive starts, and the contents are changed at the time of setting the RXS flag in IERSR.

This register cannot be modified.

IERBFL is initialized by a power-on reset or in deep standby.

20.3.13 IEBus Lock Address Register 1 (IELA1)

IELA1 specifies the lower eight bits of a locked address when a unit is locked.

IELA1 is initialized by a power-on reset or in deep standby.

20.3.14 IEBus Lock Address Register 2 (IELA2)

IELA2 specifies the upper four bits of a locked address when a unit is locked.

IELA2 is initialized by a power-on reset or in deep standby.

20.3.15 IEBus General Flag Register (IEFLG)

IEFLG indicates the IEB command execution status, lock status and slave address match, and broadcast reception detection.

IEFLG is initialized by a power-on reset or in deep standby.

20.3.16 IEBus Transmit Status Register (IETSR)

IETSR detects events such as transmit start, transmit normal completion, and transmit error end. Each status flag in IETSR corresponds to a bit in the IEBus transmit interrupt enable register (IEIET) that enables or disables each interrupt. This register is cleared by writing 1 to each bit.

IETSR is initialized by a power-on reset or in deep standby.

Note: * only 1 can be written to clear the flag.

20.3.17 IEBus Transmit Interrupt Enable Register (IEIET)

IEIET enables/disables interrupts for sources such as transmit start, transmit normal completion, and transmit error completion in IETSR.

IEIET is initialized by a power-on reset or in deep standby.

20.3.18 IEBus Receive Status Register (IERSR)

IERSR detects receive busy, receive start, receive normal completion, or receive completion with an error. Each status flag in IERSR corresponds to a bit in the IEIER that enables/disables each interrupt. This register is cleared by writing 1 to each bit.

IERSR is initialized by a power-on reset or in deep standby.

Note: $*$ only 1 can be written to clear the flag.

20.3.19 IEBus Receive Interrupt Enable Register (IEIER)

IEIER enables/disables interrupts for sources such as IERSR receive busy, receive start, receive normal completion, and receive error completion.

IEIER is initialized by a power-on reset or in deep standby.

20.3.20 IEBus Clock Selection Register (IECKSR)

IECKSR is an 8-bit readable/writable register that specifies the clock used in the IEB. IECKSR can be read from or written to even when the IEB module is stopped. IECKSR is initialized by a power-on reset or in deep standby mode.

Notes: 1. Do not change the setting of CKS3 to CKS0 while IEBus is in transmit/receive operation

 2. After the MSTP4 bit in STBCR2 is cleared to 0 to start the operation of the IEB with the CKS3 bit set to 1 (the AUDIO_X1 and AUDIO_X2 clocks are used), do not set the MSTP4 bit to 1 (clock supply to the IEB except IECKSR is stopped). As for the setting of the STBCR2 register, see section 27, Power-Down Modes.

20.3.21 IEBus Transmit Data Buffer 001 to 128 (IETB001 to IETB128)

IETB001 to IETB128 are 128-byte (8×128) buffers to which data to be transmitted during master transmission is written.

IETB001 to IETB128 are initialized by a power-on reset or in deep standby. The initial values are undefined.

Note: * Writing to these bits during master transmission (MRQ in IEFLG is 1) is prohibited.

20.3.22 IEBus Receive Data Buffer 001 to 128 (IERB001 to IERB128)

IERB001 to IERB128 are 128-byte (8×128) buffers to which data to be transmitted during slave transmission is written.

IERB001 to IERB128 are initialized by a power-on reset or in deep standby. The initial values are undefined.

Note: * Reading these bits during slave reception (SRE in IEFLG is 1 and RXBSY in IERSR 0) is prohibited. (Read value is undefined.)

20.4 Data Format

20.4.1 Transmission Format

Figure 20.6 shows the relationship between the transfer format and each register during the IEBus data transmission.

Figure 20.6 Relationship between Transfer Format and Each Register during IEBus Data Transmission

20.4.2 Reception Format

Figure 20.7 shows the relationship between the transfer format and each register during the IEBus data reception.

Figure 20.7 Relationship between Transfer Format and Each Register during IEBus Data Reception

20.5 Software Control Flows

20.5.1 Initial Setting

Figure 20.8 shows the flowchart for the initial setting.

Figure 20.8 Flowchart for Initial Setting

20.5.2 Master Transmission

Figure 20.9 shows the flowchart for master transmission.

Figure 20.9 Flowchart for Master Transmission

20.5.3 Slave Reception

Figure 20.10 shows the flowchart for slave reception.

Figure 20.10 Flowchart for Slave Reception

20.5.4 Master Reception

Figure 20.11 shows the flowchart for master reception.

Figure 20.11 Flowchart for Master Reception

20.5.5 Slave Transmission

Figure 20.12 shows the flowchart for slave transmission.

Figure 20.12 Flowchart for Slave Transmission

20.6 Operation Timing

20.6.1 Master Transmit Operation

Figure 20.13 shows the timing for master transmit operation.

Figure 20.13 Master Transmit Operation Timing

20.6.2 Slave Receive Operation

Figure 20.14 shows the timing for slave receive operation.

Figure 20.14 Slave Receive Operation Timing

20.6.3 Master Receive Operation

Figure 20.15 shows the timing for master receive operation.

Figure 20.15 Master Receive Operation Timing

20.6.4 Slave Transmit Operation

Figure 20.16 shows the timing for slave transmit operation.

Figure 20.16 Slave Transmit Operation Timing
20.7 Interrupt Sources

IEB interrupt sources include the following:

- Transmit start (TXS)
- Transmit normal completion (TXF)
- Arbitration loss (TXEAL)
- Transmit timing error (TXETTME)
- Overflow of the maximum number of transmit bytes in one frame (TXERO)
- Acknowledge bits (TXEACK)
- Receive busy (RXBSY)
- Receive start (RXS)
- Receive normal completion (RXF)
- Broadcast Receive Error (RXEDE)
- Receive overrun flag (RXEOVE)
- Receive timing error (RXERTME)
- Overflow of the maximum number of receive bytes in one frame (RXEDLE)
- Parity error (RXEPE)

Each source has bits corresponding to the IEBus transmit interrupt enable register (IEIET) and the IEBus receive interrupt enable register (IEIER) and can enable/disable interrupts. Each source also has status flags corresponding to the IEBus transmit status register (IETSR) and IEBus receive status register (IERSR). Reading the status flags enables determination of the interrupt sources.

Figure 20.17 shows the relations between the IEB interrupt sources.

Figure 20.17 Relations between IEB Interrupt Sources

20.8 Usage Notes

20.8.1 Notes when the Communications have not been Completed within the Maximum Number of Transmit Bytes

(1) Data Transmission

During the data transmission, when the maximum number of bytes defined by the communications mode have been transmitted because a NAK has been received from the receive unit, or when transmission has not been completed within the maximum number of transmit bytes because the message length value exceeds it, the IEB sets the error flag of IETSR and enters the wait state. In this case, the IEB transmits data for the maximum number of transmit bytes + one byte. Then, if a NAK has been received in an acknowledge bit of data for the maximum number of transmit bytes + one byte, the TXERO flag is set. If an ACK is received, the TXF flag is set.

Figure 20.18 shows the operation timing when the transmission has not been completed within the maximum number of transmit bytes.

Figure 20.18 Operation Timing when Transmission has not been Completed within Maximum Number of Transmit Bytes

(2) Data Reception

During the data reception, when reception has not been completed within the maximum number of bytes defined by the communications mode because of a parity error or overrun error causing the retransfer of data, or when reception has not been completed because the message length value exceeds the maximum number of receive bytes, the IEB sets the error flag of IERSR and enters the wait state. In this case, the IEB waits for data for the maximum number of receive bytes + one byte. Then, if data for the maximum number of receive bytes + one byte is not received, a receive timing error is detected and the RXERTME flag is set. In this case, the RXEDLE flag is not set. The RXEDLE flag is set when data for the maximum number of receive bytes + one byte is received.

The IEB also waits for data for the maximum number of receive bytes + one byte, when the maximum number of receive bytes have been received but the parity error is not cleared. If data for the maximum number of receive bytes + one byte is not received, the RXERTME flag is set. In this case, the RXEPE flag is not set. The RXEPE flag is set when data for the maximum number of receive bytes + one byte is received.

Figure 20.19 shows the operation timing when the reception has not been completed within the maximum number of receive bytes.

Figure 20.19 Operation Timing when Reception has not been Completed within Maximum Number of Receive Bytes

Section 21 CD-ROM Decoder (ROM-DEC)

The CD-ROM decoder (ROM-DEC) decodes streams of data transferred from the CD-DSP. When the medium is $CD-DA^*$, the data stream is not input to the $CD-ROM$ decoder because it consists of PCM data. In the case of $CD-ROM*^2$, the stream of data is input and the CD-ROM decoder performs sync code detection and maintenance, descrambling, ECC correction, and EDC checking, and outputs the resulting stream of data.

However, since the stream received by the CD-ROM decoder is assumed to consist of data from a CD-ROM transferred via the SSI, the decoder does not bother with the subcodes defined in the CD-DA standard.

Notes: 1. Compliant with JIS S 8605 (Red Book)

2. Compliant with JIS X 6281 (Yellow Book)

21.1 Features

• Sync-code detection and maintenance

Detects sync codes from the CD-ROM and is capable of providing sync-code maintenance (automatic interpolation of sync codes) when the sync code cannot be detected because of defects such as scratches on the disc.

Supported maintenance modes are automatic sync maintenance mode, external sync mode, interpolated sync mode, and interpolated sync plus external sync mode.

- Descrambling
- ECC support

P-parity-based correction, Q-parity-based correction, PQ correction, and QP correction are available.

PQ correction and QP correction can be applied repeatedly up to three times. This, however, depends on the speed of the CD. For example, three iterations are possible when the CD-ROM decoder is operating at 60 MHz with a double-speed CD drive.

Two buffers are provided due to the need for ECC correction. This allows parallel operation, where ECC correction is performed in one buffer while the data stream is being received in the other.

• EDC checking

The EDC is checked before and after correction based on the ECC. Furthermore, an operating mode is available in which, if the result of pre-correction EDC checking indicates no errors, ECC correction is not performed regardless of the result of syndrome calculation.

• Data buffering control

The CD-ROM decoder outputs data to the buffer area in a specific format where the sync code is at the head of the data for each sector.

21.1.1 Formats Supported by ROM-DEC

The CD-ROM decoder of this LSI supports the five formats shown in figure 21.1.

Figure 21.1 Formats Supported by ROM-DEC

21.2 Block Diagrams

Figure 21.2 is a block diagram of the CD-ROM decoder functions of this LSI and the bus bridge for connection to the peripheral bus, that is, of the elements required to implement the CD-ROM decoder function.

Figure 21.2 ROM-DEC Block Diagram

The core of the CD-ROM decoder executes a series of processing required for CD-ROM decoding, including descrambling, sync code detection, ECC correction (P- and Q-parity-based correction), and EDC checking. The core includes sufficient memory to hold two sectors.

Input data come from the peripheral bus and output data go out via the peripheral bus along a single line each, but the bus bridge logic sets up branches for the register access port and stream data port.

The stream data from the CD-DSP are transferred via the SSI to the stream data input control block. They are then subjected to descrambling, ECC correction, and EDC checking as they pass through the CD-ROM decoder. After these processes, data from one sector are obtained. The data are subsequently transferred to the stream-data buffer via the stream-data output control block. Data can be transferred by either the DMAC or the CPU.

Figure 21.3 is a block diagram of the bus-bridge logic.

Since the input stream is transferred over the SSI, transfer is relatively slow. On the other hand, data from the output stream can be transferred at high speeds because they are already in the core of the CD-ROM decoder. Since the data for output are buffered in SDRAM or other memory, they must be transferred at high speeds in order to reduce the busy rate of the SDRAM. For this reason, the data for the output stream are read out before the CD-ROM decoder receives an output stream data read request from the peripheral bus. This allows the accumulation of streaming data in the registers of the bus bridge, so that the data are ready for immediate output to the peripheral bus upon a request from the peripheral bus. Accordingly, the reception of a request to read from registers other than the stream-data registers after the stream data has already been read out and stored in the register of the bus bridge is possible. To cope with this, the CD-ROM decoder is provided with separate intermediary registers for the output stream-data register and the other registers.

Figure 21.3 Schematic Diagram of the Bus Bridge

Figure 21.4 is a schematic diagram of the stream-data input control block. The stream-data input controller contains logic that controls the stream of input data and a register that is used to change the control mode of the CD-ROM decoder.

The SSI mode used to transfer the stream data may affect the order (through the endian setting) or lead to padding before the data is transferred. To handle the different arrangements of data appropriately, the stream-data input control block includes a register for changing the operating mode and generates signals to control the core of the CD-ROM decoder.

The data holding registers for the input stream consists of two 16-bit registers. The data holding registers are controlled according to the mode set in the control register. For example, controlling the order in which 16-bit data is supplied to the core of the CD-ROM decoder (sending the second 16-bytes first or vice versa). It is also possible to stop the supply of padding data to the core of the CD-ROM decoder.

Figure 21.4 Schematic Diagram of the Stream-Data Input Control Block

Figure 21.5 is a schematic diagram of the stream-data output control block.

On recognizing that one sector of CD-ROM data is ready in the core of the CD-ROM decoder, this block ensures that the output stream-data register in the bus bridge section is empty and then starts to acquire the data for output from the core of the CD-ROM decoder

Figure 21.5 Schematic Diagram of the Stream-Data Output Control Block

This block has functions related to INTC interrupts and DMAC activation control such as suspending and masking of interrupts, turning interrupt flags off after they are read, asserting the activation signal to the DMAC, and negating the activation signal according to the detected amount of data that has been transferred.

21.3 Register Descriptions

The ROM-DEC has the following registers.

Table 21.1 Register Configuration

21.3.1 ROM-DEC Enable Control Register (CROMEN)

CROMEN enables subcode processing and CD-ROM decoding, and stops CD-ROM decoding forcibly.

21.3.2 Sync Code-Based Synchronization Control Register (CROMSY0)

CROMSY0 selects the sync code maintenance function.

Table 21.2 Register Settings for Sync Code Maintenance Modes

21.3.3 Decoding Mode Control Register (CROMCTL0)

CROMCTL0 enables/disables the various functions, selects criteria for mode or form determination, and specifies the sector type. The setting of this register becomes valid at the sector-to-sector transition

Note: The setting of this register is reapplied on each sector-to-sector transition.

21.3.4 EDC/ECC Check Control Register (CROMCTL1)

CROMCTL1 controls EDC/ECC checking. The setting of this register becomes valid at the sectorto-sector transition

Note: The setting of this register is reapplied on each sector-to-sector transition.

21.3.5 Automatic Decoding Stop Control Register (CROMCTL3)

CROMCTL3 is used to select abnormal conditions on which decoding will be automatically stopped. When decoding is stopped in response to any of the selected conditions, an IBUF interrupt is generated and the condition is indicated in the CBUFST1 register. The setting of this register becomes valid at the sector-to-sector transition

Note: The setting of this register is reapplied on each sector-to-sector transition.

21.3.6 Decoding Option Setting Control Register (CROMCTL4)

CROMCTL4 enables/disables buffering control at link block detection, specifies the information indicated by the status register, and controls the ECC correction mode. The setting of this register becomes valid at the sector-to-sector transition

Note: The setting of this register is reapplied on each sector-to-sector transition.

21.3.7 HEAD20 to HEAD22 Representation Control Register (CROMCTL5)

CROMCTL5 specifies the representation mode for HEAD20 to HEAD22.

21.3.8 Sync Code Status Register (CROMST0)

CROMST0 indicates various status information in sync code maintenance modes.

Bit:	7	6	5	4	3	2	1	0	
—	-	$\begin{array}{r}\n$ ST	ST						
Initial value:	0	0	0	0	0	0	0	0	
RW:	R	R	R	R	R	R	R	R	

21.3.9 Post-ECC Header Error Status Register (CROMST1)

CROMST1 indicates error status in the post-ECC header.

21.3.10 Post-ECC Subheader Error Status Register (CROMST3)

CROMST3 indicates error status in the post-ECC subheader.

21.3.11 Header/Subheader Validity Check Status Register (CROMST4)

CROMST4 indicates errors relating to the automatic mode determination or form determination for Mode 2.

21.3.12 Mode Determination and Link Sector Detection Status Register (CROMST5)

CROMST5 indicates the result of automatic mode determination and link block detection.

21.3.13 ECC/EDC Error Status Register (CROMST6)

CROMST6 indicates ECC processing error or EDC check error before/after ECC correction.

21.3.14 Buffer Status Register (CBUFST0)

CBUFST0 indicates that the system is searching for the first sector to be buffered, or that buffering is in progress.

21.3.15 Decoding Stoppage Source Status Register (CBUFST1)

CBUFST1 indicates that decoding/buffering has been stopped due to some errors.

A bit in this register can only be set when the corresponding bit in the CROMCTL3 register is set to 1.

21.3.16 Buffer Overflow Status Register (CBUFST2)

CBUFST2 indicates that a sector-to-sector transition occurred before data transfer to the buffer is completed.

21.3.17 Pre-ECC Correction Header: Minutes Data Register (HEAD00)

HEAD00 indicates the minutes value in the header before ECC correction.

21.3.18 Pre-ECC Correction Header: Seconds Data Register (HEAD01)

HEAD01 indicates the seconds value in the header before ECC correction.

21.3.19 Pre-ECC Correction Header: Frames (1/75 Second) Data Register (HEAD02)

HEAD02 indicates the frames value (1 frame = 1/75 second) in the header before ECC correction.

21.3.20 Pre-ECC Correction Header: Mode Data Register (HEAD03)

HEAD03 indicates the mode value in the header before ECC correction.

21.3.21 Pre-ECC Correction Subheader: File Number (Byte 16) Data Register (SHEAD00)

SHEAD00 indicates the file number value in the subheader before ECC correction (byte 16).

21.3.22 Pre-ECC Correction Subheader: Channel Number (Byte 17) Data Register (SHEAD01)

SHEAD01 indicates the channel number value in the subheader before ECC correction (byte 17).

21.3.23 Pre-ECC Correction Subheader: Sub-Mode (Byte 18) Data Register (SHEAD02)

SHEAD02 indicates the sub-mode value in the subheader before ECC correction (byte 18).

21.3.24 Pre-ECC Correction Subheader: Data Type (Byte 19) Data Register (SHEAD03)

SHEAD03 indicates the data type value in the subheader before ECC correction (byte 19).

21.3.25 Pre-ECC Correction Subheader: File Number (Byte 20) Data Register (SHEAD04)

SHEAD04 indicates the file number value in the subheader before ECC correction (byte 20).

21.3.26 Pre-ECC Correction Subheader: Channel Number (Byte 21) Data Register (SHEAD05)

SHEAD05 indicates the channel number value in the subheader before ECC correction (byte 21).

21.3.27 Pre-ECC Correction Subheader: Sub-Mode (Byte 22) Data Register (SHEAD06)

SHEAD06 indicates the sub-mode value in the subheader before ECC correction (byte 22).

21.3.28 Pre-ECC Correction Subheader: Data Type (Byte 23) Data Register (SHEAD07)

SHEAD07 indicates the data type value in the subheader before ECC correction (byte 23).

21.3.29 Post-ECC Correction Header: Minutes Data Register (HEAD20)

HEAD20 indicates the minutes value in the header after ECC correction.

21.3.30 Post-ECC Correction Header: Seconds Data Register (HEAD21)

HEAD21 indicates the seconds value in the header after ECC correction.

21.3.31 Post-ECC Correction Header: Frames (1/75 Seconds) Data Register (HEAD22)

HEAD22 indicates the frames value (1 frame = $1/75$ seconds) in the header after ECC correction.

21.3.32 Post-ECC Correction Header: Mode Data Register (HEAD23)

HEAD23 indicates the mode value in the header after ECC correction.

21.3.33 Post-ECC Correction Subheader: File Number (Byte 16) Data Register (SHEAD20)

SHEAD20 indicates the file number value in the subheader after ECC correction (byte 16).

Bit: 7 6 5 4 3 2 1 0 Initial value: R/W: 0 0 0 0 0 0 0 0 R R R R R R R R SHEAD20[7:0]

21.3.34 Post-ECC Correction Subheader: Channel Number (Byte 17) Data Register (SHEAD21)

SHEAD21 indicates the channel number value in the subheader after ECC correction (byte 17).

21.3.35 Post-ECC Correction Subheader: Sub-Mode (Byte 18) Data Register (SHEAD22)

SHEAD22 indicates the sub-mode value in the subheader after ECC correction (byte 18).

21.3.36 Post-ECC Correction Subheader: Data Type (Byte 19) Data Register (SHEAD23)

SHEAD23 indicates the data type value in the subheader after ECC correction (byte 19).

21.3.37 Post-ECC Correction Subheader: File Number (Byte 20) Data Register (SHEAD24)

SHEAD24 indicates the file number value in the subheader after ECC correction (byte 20).

Bit: 7 6 5 4 3 2 1 0 Initial value: R/W: 0 0 0 0 0 0 0 0 R R R R R R R R SHEAD24[7:0]

21.3.38 Post-ECC Correction Subheader: Channel Number (Byte 21) Data Register (SHEAD25)

SHEAD25 indicates the channel number value in the subheader after ECC correction (byte 21).

21.3.39 Post-ECC Correction Subheader: Sub-Mode (Byte 22) Data Register (SHEAD26)

SHEAD26 indicates the sub-mode value in the subheader after ECC correction (byte 22).

21.3.40 Post-ECC Correction Subheader: Data Type (Byte 23) Data Register (SHEAD27)

SHEAD27 indicates the data type value in the subheader after ECC correction (byte 23).

21.3.41 Automatic Buffering Setting Control Register (CBUFCTL0)

21.3.42 Automatic Buffering Start Sector Setting: Minutes Control Register (CBUFCTL1)

CBUFCTL1 indicates the minutes value in the header for the first sector to be buffered.

21.3.43 Automatic Buffering Start Sector Setting: Seconds Control Register (CBUFCTL2)

CBUFCTL2 indicates the seconds value in the header for the first sector to be buffered.

21.3.44 Automatic Buffering Start Sector Setting: Frames Control Register (CBUFCTL3)

CBUFCTL3 indicates the frames (1 frame = 1/75 second) value in the header for the first sector to be buffered.

21.3.45 ISY Interrupt Source Mask Control Register (CROMST0M)

CROMST0M masks the ISY interrupt sources specified by the bits in the sync code status register (CROMST0).

21.3.46 CD-ROM Decoder Reset Control Register (ROMDECRST)

ROMDECRST resets the random logic of the CD-ROM decoder, clears the RAM in the CD-ROM decoder, and masks reset of the CD-ROM decoder.

Note: Before setting LOGICRST to 1, make sure that both the RAMRST and INHBUSCAN bits are clear and then make the setting by writing B'10000000 to this register.

21.3.47 CD-ROM Decoder Reset Status Register (RSTSTAT)

RSTSTAT indicates that the RAM in the CD-ROM decoder has been cleared.

21.3.48 SSI Data Control Register (SSI)

SSI provides various settings related to the data stream. For the operation corresponding to the setting of this register, see section 21.4.1, Endian Conversion for Data in the Input Stream.

21.3.49 Interrupt Flag Register (INTHOLD)

INTHOLD consists of various interrupt flags.

21.3.50 Interrupt Source Mask Control Register (INHINT)

INHINT controls masking of various interrupt requests in the CD-ROM decoder.

21.3.51 Buffer Control Register (RINGBUFCTL)

RINGBUFCTL introduces a delay in the DMA transfer request signal in order to reduce the time period over which the CD-ROM decoder is unable to respond due to ECC correction.

Data from the CD-ROM decoder may be transferred by DMA or by the CPU; the former is triggered by asserting the DMA transfer request signal, while the latter is driven by the IREADY interrupt. When either the DMAC or CPU is trying to get output data for which the input data contained errors and is thus being subjected to ECC correction, a wait will be necessary because the CD-ROM decoder cannot respond during ECC correction.

21.3.52 CD-ROM Decoder Stream Data Input Register (STRMDIN0)

STRMDIN0 indicates the 1 byte (from MSB) of the 4 bytes of data that is to be input to the CD-ROM decoder.

Bit: 7 6 5 4 3 2 1 0 Initial value: R/W: 0 0 0 0 0 0 0 0 R/W R/W R/W R/W R/W R/W R/W R/W STRMDIN[31:24]

21.3.53 CD-ROM Decoder Stream Data Input Register (STRMDIN1)

STRMDIN1 indicates the next 1 byte, in order from STRMDIN0, of the 4-byte input to the CD-ROM decoder.

21.3.54 CD-ROM Decoder Stream Data Input Register (STRMDIN2)

STRMDIN2 indicates the next 1 byte, in order from STRMDIN1, of the 4-byte input to the CD-ROM decoder.

Bit: 7 6 5 4 3 2 1 0 Initial value: R/W: 0 0 0 0 0 0 0 0 R/W R/W R/W R/W R/W R/W R/W R/W STRMDIN[15:8]

21.3.55 CD-ROM Decoder Stream Data Input Register (STRMDIN3)

STRMDIN3 indicates the 1 byte (from LSB) of the 4 bytes of data that is to be input to the CD-ROM decoder.

Bit: 7 6 5 4 3 2 1 0 Initial value: R/W: 0 0 0 0 0 0 0 0 R/W R/W R/W R/W R/W R/W R/W R/W STRMDIN[7:0]

21.3.56 CD-ROM Decoder Stream Data Output Register (STRMDOUT0)

STRMDOUT0 indicates the 1 byte (from MSB) of the 2 bytes of data that is to be output to the CD-ROM decoder.

21.3.57 CD-ROM Decoder Stream Data Output Register (STRMDOUT1)

STRMDOUT1 indicates the 1 byte (from LSB) of the 2 bytes of data that is to be output to the CD-ROM decoder.

21.4 Operation

21.4.1 Endian Conversion for Data in the Input Stream

Stream data must be input to the core of the CD-ROM decoder in order according to the CD-ROM data format specifications. In some systems, however, the order of the data from the SSI may have to be changed or the data will have been padded before transfer. To cope with this, the stream data input control section is capable of swapping the order of the data and preventing the input of padding data to the core of the CD-ROM decoder. These functions are controlled through the SSI data control register (SSI).

Figure 21.6 shows a case where the upper and lower 16 bits of the data, consisting of padding data plus the first 2 bytes of sync code, that is, H'000000FF, are swapped (H'00FF0000) and input to the CD-ROM decoder as the stream data.

Figure 21.6 Example of Padded Stream Data Control by the SSI Register

Figure 21.7 shows a case of input stream data that has no padding (H'12345678). The upper and lower 16 bits of data are swapped (H'56781234) for input to the CD-ROM decoder.

21.4.2 Sync Code Maintenance Function

Each sector of CD-ROM data consists of 2352 bytes starting with

H'00FFFFFFFFFFFFFFFFFFFF00 (sync code). However, a scratch on the disc or some other factor might lead to erroneous recognition of the sync code sequence at the wrong time. Conversely, a sync code might not be detected at a point where it should be detected. As a solution to these problems, the CD-ROM decoder of this LSI has a sync-code maintenance function, which operates to ignore sync codes detected at abnormal times and maintain the appearance of the sync code at the expected times when it is not actually detected on the disc.

The operating modes of the sync-code maintenance function are listed below. For details on the settings, refer to section 21.3.2, Sync Code-Based Synchronization Control Register (CROMSY0), and table 21.2.

- Automatic sync maintenance mode
- External sync mode
- Interpolated sync mode
- Interpolated sync plus external sync mode

(1) Automatic Sync Maintenance Mode

In automatic sync maintenance mode, the sync code is ignored if detected within the one-sector (2352-byte) period. Furthermore, if a sync code is not detected at the point where a next sector should start, sync code maintenance is applied. If synchronization timing has changed, resynchronization is performed at the point where a sync code is detected within 2352 bytes after the change.

Therefore, this mode is effective in rejecting abnormal sync patterns and following changes in synchronization timing. Note, however, that this mode cannot achieve synchronization with the first sector after a change to the synchronization timing.

Figure 21.8 shows operation in the case of normal sync-code detection, figure 21.9 shows a case where a sync code is detected before a current one-sector period has elapsed, and figure 21.10 shows the case where the actual sync code is only detected some time after a full one-sector period has elapsed.

Figure 21.8 Operation in Automatic Sync Maintenance Mode (Normal Timing)

Figure 21.9 Operation in Automatic Sync Maintenance Mode (When an Abnormally Short Sector is Encountered)

Figure 21.10 Operation in Automatic Sync Maintenance Mode (When an Abnormally Long Sector is Encountered)

(2) External Sync Mode

In external sync mode, synchronization is always based on the sync codes in the incoming data. Even if the next sync code is not detected at the 2352nd byte, decoding does not proceed until the next sync code is detected.

Accordingly, this mode is effective in that it strictly follows the external synchronization timing. Note, however, that decoding will not be performed normally when the sync-code pattern is input with abnormal timing.

Figure 21.11 shows the operation in external sync mode.

Figure 21.11 Operation in External Sync Mode

(3) Interpolated Sync Mode

In interpolated sync mode, synchronization is always driven by the internal counter after a sync code pattern has been detected at the start of decoding. Accordingly, this mode is effective when the sync patterns have been damaged.

However, decoding becomes incorrect after a change to the synchronization timing, since the change in timing is not followed.

Figure 21.12 shows the operation in interpolated sync mode.

Figure 21.12 Operation in Interpolated Sync Mode

(4) Interpolated Sync Plus External Sync Mode

In interpolated sync plus external sync mode, synchronization is based on the detected sync code patterns as long as they are present, and if a sync pattern is not detected at the 2352nd byte, the sync code maintenance is applied. Synchronization in this mode is more quickly responsive to changes in synchronization timing than synchronization in the automatic sync maintenance mode.

However, decoding still becomes incorrect when a sync pattern is input with abnormal timing.

Figures 21.13 and 21.14 show the operation in interpolated sync plus external sync mode in the cases of abnormally short and long sectors, respectively.

Figure 21.13 Operation in Interpolated Sync Plus External Sync Mode (When an Abnormally Short Sector is Encountered)

Figure 21.14 Operation in Interpolated Sync Plus External Sync Mode (When an Abnormally Long Sector is Encountered)

21.4.3 Error Correction

The CD-ROM decoder handles data in the formats containing information relevant to error correction, including the EDC, P parity, and Q parity. The CD-ROM decoder includes the following features for use in error correction.

- Syndrome calculation
- ECC correction
- EDC checking

(1) Syndrome Calculation

After the data of a sector in Mode 1 or Form 1 of Mode 2 has been input, the ECC is used in correction if any error is detected (the result of syndrome calculation is non-zero). After correction, the results of syndrome operation for the corrected data are output to bits ST_ECCP (P parity) and ST_ECCQ (Q parity) in the CROMST6 register, respectively.

(2) ECC correction and EDC Checking

For CD-ROM format data that contains EDC, P-parity, and Q-parity fields, the CD-ROM decoder performs EDC checking and ECC correction. Supported correction modes are P correction, Q correction, PQ correction (P correction followed by Q correction), and QP correction (Q correction followed by P correction). In PQ and QP correction modes, up to three iterations of correction are possible (the number of iterations is limited by the playback speed).

The EDC check is performed twice, before and after correction.

The mode of ECC correction and EDC checking is specified by bits MD_DEC[2:0] in the CROMCTL1 register. When the PQ or QP correction mode is selected, the number of iterations is specified by bits MD_PQREP[1:0] in the CROMCTL1 register.

When the automatic mode/form detection function is in use, the sector mode determines whether or not ECC correction and EDC checking can be performed. For sectors in Mode 0 and Mode 2 (non-XA), which include neither parity bits nor EDC, ECC correction and EDC checking are not performed. For sectors in Form 2 of Mode 2, ECC correction is not performed.

(a) ECC Correction

When ECC correction is in use and an error in a sector is identified as non-correctable, the CD-ROM decoder generates an IERR interrupt and sets the ST_ECCNG bit of the CROMST6 register to 1. The CD-ROM detector also sets this bit to 1 on detecting a short sector.

While the NO ECC bit of the CROMCTL4 register is set to 1, a 'pass' result in pre-correction EDC checking makes the CD-ROM decoder skip ECC correction, regardless of the results of the syndrome operation.

(b) EDC Checking

When EDC checking is in use, checking is in line with the specified or detected sector mode and form, depending on whether or not automatic sector mode and form detection is selected.

The results of EDC checking before and after correction are reflected in the ST_EDC1 and ST_EDC2 bits of the CROMST6 register, respectively. If EDC checking after ECC correction indicates that an error remains, an IERR interrupt is generated.

21.4.4 Automatic Decoding Stop Function

Decoding can be stopped automatically in response to an error during the decoding of CD-ROM data.

The possible conditions for automatically stopping the decoding process are listed below. The applicable conditions are specified in the CROMCTL3 register.

- An error is found to be not correctable by ECC correction.
- Post-correction EDC checking indicates that an error remains.
- A change of the sector mode or form.
- A non-sequential MSF (minutes, seconds, frames (1/75 second)) value.

When automatic stopping is set up and any of the above conditions is encountered in a certain sector, the decoding is stopped after the results of decoding for that sector have been output.

After decoding has been stopped in response to a condition specified in the CROMCTL3 register, the condition can be identified by reading the CBUFST1 register.

The CD-ROM decoder has buffer space for two sectors. If input of the data stream continues and the output stream of data is not read, the CD-ROM decoder stops at the point where the data of a third sector starts to be input. At this time, the BUF_NG bit in the CBUFST2 register is set to 1, but no interrupt is generated. Once the BUF_NG bit in the CBUFST2 register has been set to 1, recovery can only be accomplished by using the LOGICRST bit in the ROMDECRST register to reset the CD-ROM decoder function. When the LOGICRST bit in the ROMDECRST register is set to 1, a reset signal is output and any registers in which settings have been made are cleared to their initial values.

21.4.5 Buffering Format

Figure 21.15 shows the format of the output data stream produced by CD-ROM decoding.

A 2-byte-wide window register is provided for the output. When this window register is accessed after decoding of a CD-ROM sector has finished, the bytes of data are output in order from the sync code.

Figure 21.15 Output Data Stream Format

The meanings of bits in the two-byte status field shown in figure 21.15 are given below. The values of the non-assigned bits are undefined.

PERR: Indicates that a P-parity error remains.

QERR: Indicates that a Q-parity error remains.

EDCE: Indicates that a remaining error was detected in post-correction EDC checking.

SD: Indicates that a short sector was encountered

SY: Indicates that a sync code was interpolated.

- FM: Indicates the data format
	- 001: Mode 0
	- 010: Mode 1
	- 011: Long (format with no EDC and ECC)
	- 100: Mode 2 (non-XA)
	- 101: Mode 2 Form 1
	- 110: Mode 2 Form 2
- HD: Header continuity (minutes, seconds, and frames (1/75) are non-sequential)

The value of the storage flag field in figure 21.15 is incremented every time the data for one sector are output. The value starts at H'0000 and wraps back around to H'0000 after incrementation reaches H'FFFF. Note that the upper byte and lower byte in the storage flag are swapped.

21.4.6 Target-Sector Buffering Function

In the CD-ROM decoder, the sector for output can be designated in two ways: automatic buffering, where the CD-ROM decoder itself detects the presence of target sectors, and manual buffering, where the target sector for output is designated by software and the software also recognizes the sectors buffered in the CD-ROM decoder.

The following describes the procedures for setting the registers in the CD-ROM decoder to set up automatic or manual buffering.

(1) Setting Up Automatic Buffering

Figure 21.16 shows an example of setting up the automatic buffering. Set the relevant CD-ROM decoder registers and start input of the data stream; the CD-ROM decoder then detects the target sector and starts the output of the stream data.

Figure 21.16 Example of Setting Up Automatic Buffering

(2) Setting Up Manual Buffering

Figure 21.17 shows an example of setting up manual buffering. Each time an ISEC interrupt is generated, the software checks whether or not the sector is the target sector and starts buffering when the target sector is found .

Figure 21.17 Example of Setting Up Manual Buffering

21.5 Interrupt Sources

21.5.1 Interrupt and DMA Transfer Request Signals

Table 21.3 lists the interrupt signals and DMA transfer request signal generated by the CD-ROM decoder, along with the meanings and the modules to which the signals are connected.

Name	Description	Connected To
ISEC	Transitions from sector to sector	INTC
ITARG	Access to a CD-ROM sector that is not the expected target sector	INTC
ISY	A sync code from the CD-ROM with abnormal timing	INTC
IERR	An error that was not correctable by ECC correction or an error INTC indicated by EDC checking after ECC correction	
IBUF	State changes in data transfer to the buffer	INTC
IREADY	Request for data transfer to the buffer for CD-ROM	INTC
DMA transfer request	Request for data transfer to the buffer for CD-ROM	DMAC

Table 21.3 Interrupt and DMA Transfer Request Signals

The above interrupt signals are generated by the following sources

(1) ISEC

This interrupt is generated when the sync code indicates a transition from sector to sector.

(2) ITARG

This interrupt reports that the stream data transferred from the CD-DSP is not the data of the target sector. The CD-ROM decoder checks the time data in the subcode. In correct operation, data transfer is expected to start slightly before the target sector. An ITARG interrupt is generated in the following cases.

- When data of a sector preceding the target sector by quite a few sectors have been transferred
- When data of a sector that comes after the target sector have been transferred

For the generation of this interrupt, ITARG is detected from the subcode. However, this interrupt has no meaning in this LSI because CD-ROM data are transferred from the SSI module.

(3) ISY

This interrupt can be generated in the following cases.

- When a sync code was detected at a position where the value in the word counter (counter for checking sync code intervals) was not correct and the sync code was ignored
- When a sync code has not been detected although the word counter has reached the final value and a sync code has been interpolated (for sync maintenance)
- When a sync code was detected at a position where the value in the word counter (counter for checking sync code intervals) was not correct and the sync code was used in resynchronization
- When a sync code has not been detected although the word counter has reached the final value, so the period taken up by the sector has been prolonged
- When the sector has been processed as a short sector with the aid of interpolated sync codes
- When the sector has been processed as a long sector with the aid of interpolated sync codes

(4) IERR

This interrupt is generated in the following cases.

- When ECC correction was incapable of correcting an error
- When ECC correction was OK but the subsequent EDC check indicated an error

(5) IBUF

This interrupt is generated when the following transitions occur.

- Data transfer to the buffer \rightarrow Data transfer complete (searching for data for the next transfer)
- Data for transfer to the buffer are being searched for \rightarrow Data transfer started

(6) IREADY

This interrupt is generated when decoding of data for one sector is completed. This interrupt should be used to start the CPU buffering stream data for output to SDRAM.

(7) DMA Transfer Request

The source of DMA activation is the same as that of IREADY. An interrupt request is generated when output stream data for one sector becomes ready, and after the 2768 bytes of data shown in figure 21.15 have been transferred, the request signal is negated once. This is because a certain amount of time is required before the output data for the next sector is ready, so the transfer request from the DMAC should be turned off between transfers.

21.5.2 Timing of Status Registers Updates

The status information registers of the CD-ROM decoder are updated on each ISEC interrupt. The sector for which information is reflected in the status registers is selected by the ER0SEL bit of the CROMCTL4 register.

21.6 Usage Notes

21.6.1 Stopping and Resuming Buffering Alone During Decoding

When the data of the output stream are being not read out but operation of the CD-ROM decoder has continued until the buffers are full, the BUF_NG bit in the CBUFST2 register is set to 1; after that, the CD-ROM decoder becomes incapable of operation.

To stop buffering alone, clear the CBUF_EN bit in the CBUFCTL0 register to 0. If the automatic buffering function is in use, clear the CBUF_AUT in the CBUFCTL0 register to 0 at the same time. In this case, the sectors currently in the buffers must be read out.

To resume automatic buffering, set the CBUF_AUT and CBUF_EN bits in the CBUFCTL0 register at the same time.

21.6.2 When CROMST0 Status Register Bits are Set

- 1. When the ST_SECS bit in the CROMST0 register becomes set, stop decoding immediately and retry from one sector before the sector that was being decoded.
- 2. When the ST_SECL bit in the CROMST0 register becomes set, stop decoding immediately and retry from two sectors before the sector that was being decoded.

21.6.3 Link Blocks

The CD-ROM decoder uses the header information before ECC correction to detect link blocks. Accordingly, an input data stream that contains an error may be erroneously detected as a link block. To prevent this, the following measures should be implemented in software.

- During buffering (BUF $ACT = 1$ in the CBUFST0 register), check the LINK OUT1 bit in the CROMST5 register on each ISEC interrupt. If it is set to 1, check to see if an IERR interrupt has also occurred; if an IERR interrupt has not occurred, save the MFS values from the HEAD20 to HEAD23 registers. If an IERR interrupt has occurred, do not save the MSF values.
- Perform the following processing for seven sectors (indicated by ISEC being generated seven times) after finding that the LINK_OUT1 bit has been set to 1.

In either of cases 1 and 2 below,

- 1. LINK_ON = 1 (in the CROMST5 register) is confirmed at each ISEC interrupt, and LINK $ON = 1$ is detected again within the subsequent two-sector period
- 2. LINK $ON = 1$ was not detected at any ISEC interrupt

Forcibly stop decoding, set the CROMSY0 register to place the decoder in external sync mode, and retry decoding by specifying the MSF value stored above + 7 as the MSF value for the target sector. The start sector address will be the address where RUN_OUT is stored + 7 when CBUF_LINK = 0 , and the address where RUN_OUT is stored when CBUF_LINK = 1 .

21.6.4 Reading from the STRMDOUT0 and STRMDOUT1 Registers

When the input stream of data contains an error and ECC correction is executed, the process of reading from the STRMDOUT0 and STRMDOUT1 registers will be kept waiting by the execution of ECC correction if the read request is made immediately after an IREADY interrupt or DMA transfer trigger signal has been generated. This only applies to the first time the registers are read in the reading of one sector.

In cases where CD-ROM decoding cannot be performed if kept waiting (for example, when the STRMDOUT0 and STRMDOUT1 registers are read by DMA transfer and input and output for the SSI are also handled by DMA transfer, so that DMA transfer for the SSI must be carried out with a fixed period), use the DMAREQDELAY[1:0] bits in the RINGBUFCTL register to delay the DMA activation signal. This can reduce the length of the wait before reading from the STRMDOUT0 and STRMDOUT1 registers. Table 21.4 shows the waiting times for the first reading of STRMDOUT0 and STRMDOUT1 as approximate numbers of cycles. The number of wait cycles varies according to the positions of any errors for which ECC processing is performed and so on.

Number of Wait Cycles when Delay is Introduced

Table 21.4 Number of Wait Cycles for Reading STRMDOUT0 and STRMDOUT1 Registers

When the number of the ECC correction is 1 and the DMAREQDELAY[1:0] setting is 3, the period of waiting is up to 376 cycles. If the CD-ROM decoding is not completed in this waiting time, take the waiting time into account in software for DMAC activation.

In addition, the effect of DMAREQDELAY[1:0] in reducing the wait only applies to the DMA activation signal. Since the IREADY interrupt cannot be delayed, waiting time must be taken into account in the software if the STRMDOUT0 and STRMDOUT1 registers are to be read by the CPU.

21.6.5 Stopping and Resuming CD-DSP Operation

When stopping and resuming the stream data input to the CD-ROM decoder, note that the input data stream does not stop immediately before a sync code and that the CD-ROM decoder may recognize the data as incorrect when the input stream is resumed. This happens because the system holds a combination of the data up to the point where input was stopped and data that is input from the point of resumption. Take care on this point when stopping and resuming input.
Section 22 A/D Converter (ADC)

This LSI includes a 10-bit successive-approximation A/D converter allowing selection of up to eight analog input channels.

22.1 Features

- Resolution: 10 bits
- Input channels: 8
- Minimum conversion time: 3.9 μs per channel
- Operating modes: 3
	- Single mode: A/D conversion on one channel
	- Multi mode: A/D conversion on one to four channels or on one to eight channels
	- Scan mode: Continuous A/D conversion on one to four channels or on one to eight channels
- Data registers: 8

Conversion results are held in a 16-bit data register for each channel

- Sample-and-hold function
- A/D conversion start methods: 4
	- Software
	- ⎯ Conversion start trigger from multi-function timer pulse unit 2 (MTU2)
	- Conversion start trigger from the 8-bit timer (TMR)
	- External trigger signal
- Interrupt source

An A/D conversion end interrupt (ADI) request can be generated on completion of A/D conversion.

Module standby mode can be set

Figure 22.1 Block Diagram of A/D Converter

22.2 Input/Output Pins

Table 22.1 summarizes the A/D converter's input pins.

Table 22.1 Pin Configuration

22.3 Register Configuration

The A/D converter has the following registers.

Table 22.2 Register Configuration

22.3.1 A/D Data Registers A to H (ADDRA to ADDRH)

The sixteen A/D data registers, ADDRA to ADDRH, are 16-bit read-only registers that store the results of A/D conversion.

An A/D conversion produces 10-bit data, which is transferred for storage into the ADDR corresponding to the selected channel. The 10 bits of the result are stored in the upper bits (bits 15 to 6) of ADDR. Bits 5 to 0 of ADDR are reserved bits that are always read as 0.

Access to ADDR in 8-bit units is prohibited. ADDR must always be accessed in 16-bit units.

ADDR is initialized to H'0000 by a power-on reset as well as in deep standby mode, software standby mode or module standby mode.

Table 22.3 indicates the pairings of analog input channels and ADDR.

Table 22.3 Analog Input Channels and ADDR

22.3.2 A/D Control/Status Register (ADCSR)

ADCSR is a 16-bit readable/writable register that selects the mode, controls the A/D converter, and enables or disables starting of A/D conversion by external trigger input.

ADCSR is initialized to H'0040 by a power-on reset as well as in deep standby mode, software standby mode or module standby mode.

Note: $*$ Only 0 can be written to clear the flag after 1 is read.

[Legend]

x: Don't care

Notes: 1. Only 0 can be written to clear the flag after 1 is read.

Please note that ADF flag becomes "0" in the following cases, too.

- (1) Reading the state of $ADF = 1$ with CPU.
- (2) Clearing ADF flag by having read ADDR with DMAC
- (3) Set of ADF flag according to A/D conversion end
- (4) Writing 0 in the ADF flag with CPU
- 2. To satisfy the absolute accuracy of the A/D converter characteristics, set a value greater than the minimum conversion time.

22.4 Operation

The A/D converter uses the successive-approximation method, and the resolution is 10 bits. It has three operating modes: single mode, multi mode, and scan mode. Switching the operating mode or analog input channels must be done while the ADST bit in ADCSR is 0 to prevent incorrect operation. The ADST bit can be set at the same time as the operating mode or analog input channels are changed.

22.4.1 Single Mode

Single mode should be selected when only A/D conversion on one channel is required.

In single mode, A/D conversion is performed once for the specified one analog input channel, as follows:

- 1. A/D conversion for the selected channel starts when the ADST bit in ADCSR is set to 1 by software, MTU2, TMR, or external trigger input.
- 2. When A/D conversion is completed, the A/D conversion result is transferred to the A/D data register corresponding to the channel.
- 3. After A/D conversion has completed, the ADF bit in ADCSR is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated.
- 4. The ADST bit that remains 1 during A/D conversion is automatically cleared to 0 when A/D conversion is completed, and the A/D converter becomes idle.

When the operating mode or analog input channel selection must be changed during A/D conversion, to prevent incorrect operation, first clear the ADST bit to 0 to halt A/D conversion. After making the necessary changes, set the ADST bit to 1 to start A/D conversion again. The ADST bit can be set at the same time as the mode or channel selection is switched.

Typical operations when a single channel (AN1) is selected in single mode are described next. Figure 22.2 shows a timing diagram for this example (the bits which are set in this example belong to ADCSR).

- 1. Single mode is selected, input channel AN1 is selected (CH2 = 0, CH1 = 0, CH0 = 1), the A/D interrupt is enabled (ADIE = 1), and A/D conversion is started (ADST = 1).
- 2. When A/D conversion is completed, the A/D conversion result is transferred into ADDRB. At the same time the ADF flag is set to 1, the ADST bit is cleared to 0, and the A/D converter becomes idle.
- 3. Since $ADF = 1$ and $ADIE = 1$, an ADI interrupt is requested.
- 4. The A/D interrupt handling routine starts.
- 5. The routine reads $ADF = 1$, and then writes 0 to the ADF flag.
- 6. The routine reads and processes the A/D conversion result (ADDRB).
- 7. Execution of the A/D interrupts handling routine ends. Then, when the ADST bit is set to 1, A/D conversion starts and steps 2. to 7. are executed.

Figure 22.2 Example of A/D Converter Operation (Single Mode, One Channel (AN1) Selected)

22.4.2 Multi Mode

Multi mode should be selected when performing A/D conversion once on one or more channels.

In multi mode, A/D conversion is performed once for a maximum of eight specified analog input channels, as follows:

- 1. A/D conversion starts from the analog input channel with the lowest number (e.g. AN0, AN1, …, AN3) when the ADST bit in ADCSR is set to 1 by software, MTU2, TMR, or external trigger input.
- 2. When A/D conversion is completed on each channel, the A/D conversion result is sequentially transferred to the A/D data register corresponding to that channel.
- 3. After A/D conversion on all selected channels has completed, the ADF bit in ADCSR is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated.
- 4. The ADST bit that remains 1 during A/D conversion is automatically cleared to 0 when A/D conversion is completed, and the A/D converter becomes idle. If the ADST bit is cleared to 0 during A/D conversion, A/D conversion is halted and the A/D converter becomes idle. The ADF bit is cleared by reading ADF while $ADF = 1$, then writing 0 to the ADF bit.

A/D conversion is to be performed once on all the specified channels. The conversion results are transferred for storage into the A/D data registers corresponding to the channels.

When the operating mode or analog input channel selection must be changed during A/D conversion, to prevent incorrect operation, first clear the ADST bit to 0 to halt A/D conversion. After making the necessary changes, set the ADST bit to 1. A/D conversion will start again from the first channel in the group. The ADST bit can be set at the same time as the mode or channel selection is changed.

Typical operations when three channels (AN0 to AN2) are selected in multi mode are described next. Figure 22.3 shows a timing diagram for this example.

- 1. Multi mode is selected (MDS2 = 1, MDS1 = 0), analog input channels AN0 to AN2 are selected (CH2 = 0, CH1 = 1, CH0 = 0), and A/D conversion is started (ADST = 1).
- 2. A/D conversion of the first channel (AN0) starts. When A/D conversion is completed, the A/D conversion result is transferred into ADDRA.
- 3. Next, the second channel (AN1) is selected automatically and A/D conversion starts.
- 4. Conversion proceeds in the same way through the third channel (AN2).
- 5. When conversion of all selected channels (AN0 to AN2) is completed, the ADF flag is set to 1 and the ADST bit cleared to 0.

If the ADIE bit is set to 1 at this time, an ADI interrupt is requested.

If the ADIE bit is set to 1 at this time, an ADI interrupt is requested.

22.4.3 Scan Mode

Scan mode is useful for monitoring analog inputs in a group of one or more channels at all times. In scan mode, A/D conversion is performed sequentially for a maximum of eight specified analog input channels, as follows:

- 1. A/D conversion for the selected channels starts from the analog input channel with the lowest number (e.g. AN0, AN1, ..., AN3) when the ADST bit in ADCSR is set to 1 by software, MTU2, TMR, or external trigger input.
- 2. When A/D conversion is completed on each channel, the A/D conversion result is sequentially transferred to the A/D data register corresponding to that channel.
- 3. After A/D conversion on all selected channels has completed, the ADF bit in ADCSR is set to 1. If the ADIE bit is set to 1 at this time, an ADI interrupt request is generated. The A/D converter starts A/D conversion again from the channel with the lowest number.
- 4. The ADST bit is not cleared automatically, so steps 2. and 3. are repeated as long as the ADST bit remains set to 1. When the ADST bit is cleared to 0, A/D conversion halts and the A/D converter becomes idle.

The ADF bit is cleared by reading ADF while $ADE = 1$, then writing 0 to the ADF bit.

When the operating mode or analog input channel selection must be changed during A/D conversion, to prevent incorrect operation, first clear the ADST bit to 0 to halt A/D conversion. After making the necessary changes, set the ADST bit to 1. A/D conversion will start again from the first channel in the group. The ADST bit can be set at the same time as the mode or channel selection is changed.

Typical operations when three channels (AN0 to AN2) are selected in scan mode are described as follows. Figure 22.4 shows a timing diagram for this example.

- 1. Scan mode is selected (MDS2 = 1, MDS1 = 1), analog input channels AN0 to AN2 are selected (CH2 = 0, CH1 = 1, CH0 = 0), and A/D conversion is started (ADST = 1).
- 2. A/D conversion of the first channel (AN0) starts. When A/D conversion is completed, the A/D conversion result is transferred into ADDRA.
- 3. Next, the second channel (AN1) is selected automatically and A/D conversion starts.
- 4. Conversion proceeds in the same way through the third channel (AN2).
- 5. When conversion of all the selected channels (AN0 to AN2) is completed, the ADF flag is set to 1 and conversion of the first channel (AN0) starts again. If the ADIE bit is set to 1 at this time, an ADI interrupt is requested.

6. The ADST bit is not cleared automatically, so steps 2. to 4. are repeated as long as the ADST bit remains set to 1. When steps 2. to 4. are repeated, the ADF flag is kept to 1. When the ADST bit is cleared to 0, A/D conversion stops. The ADF bit is cleared by reading ADF while $ADF = 1$, then writing 0 to the ADF bit.

If both the ADF flag and ADIE bit are set to 1 while steps 2. to 4. are repeated, an ADI interrupt is requested at all times. To generate an interrupt on completing conversion of the third channel, clear the ADF bit to 0 after an interrupt is requested.

Figure 22.4 Example of A/D Converter Operation (Scan Mode, Three Channels (AN0 to AN2) Selected)

SH7261 Group

Section 22 A/D Converter (ADC)

22.4.4 A/D Converter Activation by External Trigger, MTU2, or TMR

The A/D converter can be independently activated by an A/D conversion request from the external trigger, MTU2, or TMR. To activate the A/D converter by the external trigger, MTU2, or TMR, set the A/D trigger enable bits (TRGS[3:0]). After this bit setting has been made, the ADST bit is automatically set to 1 and A/D conversion is started when an A/D conversion request from the external trigger, MTU2, or TMR occurs. The channel combination is determined by the CH[2:0] bits in ADCSR. The timing from setting of the ADST bit until the start of A/D conversion is the same as when 1 is written to the ADST bit by software.

22.4.5 Input Sampling and A/D Conversion Time

The A/D converter has a built-in sample-and-hold circuit. The A/D converter samples the analog input at the A/D conversion start delay time (t_D) after the ADST bit in ADCSR is set to 1, then starts conversion. Figure 22.5 shows the A/D conversion timing. Table 22.4 indicates the A/D conversion time.

As indicated in figure 22.5, the A/D conversion time (t_{conv}) includes t_{D} and the input sampling time(t_{SPL}). The length of t_{D} varies depending on the timing of the write access to ADCSR. The total conversion time therefore varies within the ranges indicated in table 22.4.

In multi mode and scan mode, the values given in table 22.4 apply to the first conversion. In the second and subsequent conversions, time is the values given in table 22.5.

Table 22.4 A/D Conversion Time (Single Mode)

Note: Values in the table are the numbers of states.

Table 22.5 A/D Conversion Time (Multi Mode and Scan Mode)

Note: Values in the table are the numbers of states.

22.4.6 External Trigger Input Timing

A/D conversion can also be externally triggered. When the TRGS[3:0] bits in ADCSR are set to B'1001, an external trigger is input to the ADTRG pin. The ADST bit in ADCSR is set to 1 at the falling edge of the ADTRG pin, thus starting A/D conversion. Other operations, regardless of the operating mode, are the same as when the ADST bit has been set to 1 by software. Figure 22.6 shows the timing.

Figure 22.6 External Trigger Input Timing

22.5 Interrupt Sources and DMAC Transfer Request

The A/D converter generates an A/D conversion end interrupt (ADI) at the end of A/D conversion. An ADI interrupt request is generated if the ADIE bit is set to 1 when the ADF bit in ADCSR is set to 1 on completion of A/D conversion. Note that the direct memory access controller (DMAC) can be activated by an ADI interrupt depending on the interrupt controller (INTC) setting. In this case, an interrupt is not issued to the CPU. If the setting to activate the DMAC has not been made, an interrupt request is sent to the CPU. Having the converted data read by the DMAC in response to an ADI interrupt enables continuous conversion to be achieved without imposing a load on software.

In single mode, set the DMAC so that DMA transfer initiated by an ADI interrupt is performed only once. In the case of A/D conversion on multiple channels in scan mode or multi mode, setting the DMA transfer count to one causes DMA transfer to finish after transferring only one channel of data. To make the DMAC transfer all conversion data, set the ADDR where A/D conversion data is stored as the transfer source address, the number of converted channels \times 2 as the transfer byte count, and continuous operand transfer or non-stop transfer as the DMA transfer condition.

When the DMAC is activated by ADI, the ADF bit in ADCSR is automatically cleared to 0 when data is transferred by the DMAC.

22.6 Definitions of A/D Conversion Accuracy

The A/D converter compares an analog value input from an analog input channel with its analog reference value and converts it to 10-bit digital data. The absolute accuracy of this A/D conversion is the deviation between the input analog value and the output digital value. It includes the following errors:

- Offset error
- Full-scale error
- Quantization error
- Nonlinearity error

These four error quantities are explained below with reference to figure 22.7. In the figure, the 10 bits of the A/D converter have been simplified to 3 bits. Offset error is the deviation between actual and ideal A/D conversion characteristics when the digital output value changes from the minimum (zero voltage) B'0000000000 (000 in the figure) to B'000000001 (001 in the figure) (figure 22.7, item (1)). Full-scale error is the deviation between actual and ideal A/D conversion characteristics when the digital output value changes from B'1111111110 (110 in the figure) to the maximum B'1111111111111111 (111 in the figure) (figure 22.7, item (2)). Quantization error is the intrinsic error of the A/D converter and is expressed as 1/2 LSB (figure 22.7, item (3)). Nonlinearity error is the deviation between actual and ideal A/D conversion characteristics between zero voltage and full-scale voltage (figure 22.7, item (4)). Note that it does not include offset, full-scale, or quantization error.

Figure 22.7 Definitions of A/D Conversion Accuracy

22.7 Usage Notes

When using the A/D converter, note the following points.

22.7.1 Module Standby Mode Setting

Operation of the A/D converter can be disabled or enabled using the standby control register. The initial setting is for operation of the A/D converter to be halted. Register access is enabled by clearing module standby mode. For details, see section 27, Power-Down Modes.

22.7.2 Setting Analog Input Voltage

Permanent damage to the LSI may result if the following voltage ranges are exceeded.

1. Analog input range

During A/D conversion, voltages on the analog input pins ANn should not go beyond the following range: $AVss \leq ANn \leq AVcc$ (n = 0 to 7).

2. AVcc and AVss input voltages

Input voltages AVcc and AVss should be PVcc -0.3 V \leq AVcc \leq PVcc and AVss = PVss. Do not leave the AVcc and AVss pins open when the A/D converter or D/A converter is not in use and in software standby mode. When not in use, connect AVcc to the power supply (PVcc) and AVss to the ground (PVss).

3. Setting range of AVref input voltage Set the reference voltage range of the AVref pin as $3.0 \text{ V} \leq \text{A} \text{V}$ and $\text{A} \text{V}$ and $\text{A$

22.7.3 Notes on Board Design

In board design, digital circuitry and analog circuitry should be as mutually isolated as possible, and layout in which digital circuit signal lines and analog circuit signal lines cross or are in close proximity should be avoided as far as possible. Failure to do so may result in incorrect operation of the analog circuitry due to inductance, adversely affecting A/D conversion values.

Digital circuitry must be isolated from the analog input signals (AN0 to AN7), analog reference voltage (AVref), and analog power supply (AVcc) by the analog ground (AVss). Also, the analog ground (AVss) should be connected at one point to a stable digital ground (PVss) on the board.

22.7.4 Processing of Analog Input Pins

To prevent damage from voltage surges at the analog input pins (AN0 to AN7), connect an input protection circuit like the one shown in figure 22.8. The circuit shown also includes a CR filter to suppress noise. This circuit is shown as an example; the circuit constants should be selected according to actual application conditions.

Figure 22.9 shows an equivalent circuit diagram of the analog input ports and table 22.7 lists the analog input pin specifications.

Figure 22.9 Analog Input Pin Equivalent Circuit

Table 22.7 Analog Input Pin Ratings

22.7.5 Permissible Signal Source Impedance

This LSI's analog input is designed such that conversion precision is guaranteed for an input signal for which the signal source impedance is $5 \text{ k}\Omega$ or less. This specification is provided to enable the A/D converter's sample-and-hold circuit input capacitance to be charged within the sampling time; if the sensor output impedance exceeds 5 kΩ, charging may be insufficient and it may not be possible to guarantee A/D conversion precision. However, for A/D conversion in single mode with a large capacitance provided externally for A/D conversion in single mode, the input load will essentially comprise only the internal input resistance of 3 k Ω , and the signal source impedance is ignored. However, as a low-pass filter effect is obtained in this case, it may not be possible to follow an analog signal with a large differential coefficient (e.g., 5 mV/µs or greater) (see figure 22.10). When converting a high-speed analog signal, a low-impedance buffer should be inserted.

Figure 22.10 Example of Analog Input Circuit

22.7.6 Influences on Absolute Precision

Adding capacitance results in coupling with GND, and therefore noise in GND may adversely affect absolute precision. Be sure to connect AVss, etc. to an electrically stable GND.

Care is also required to insure that filter circuits do not communicate with digital signals on the mounting board (i.e., acting as antennas).

22.7.7 Note on Usage in Scan Mode and Multi Mode

Starting conversion immediately after having stopped scan mode or multi mode operation may lead to erroneous results of conversion.

To perform continuous conversion in such cases, set ADST to 0, wait for at least the A/D conversion time for a single channel to elapse, and then start conversion (ADST = 1). (The A/D conversion time for a single channel will vary according to the settings of the ADC registers.)

Section 23 D/A Converter (DAC)

23.1 Features

- Resolution: 8 bits
- Input channels: 2
- Minimum conversion time: Max. 10 µs (with 20 pF load)
- Output voltage: 0 V to AVref
- D/A output hold function in software standby mode
- Module standby mode can be set

Figure 23.1 Block Diagram of D/A Converter

23.2 Input/Output Pins

Table 23.1 shows the pin configuration of the D/A converter.

Table 23.1 Pin Configuration

23.3 Register Descriptions

The D/A converter has the following registers.

Table 23.2 Register Configuration

23.3.1 D/A Data Registers 0 and 1 (DADR0 and DADR1)

DADR is an 8-bit readable/writable register that stores data to which D/A conversion is to be performed. Whenever analog output is enabled, the values in DADR are converted and output to the analog output pins.

DADR is initialized to H'00 by a power-on reset in deep standby mode or module standby mode.

23.3.2 D/A Control Register (DACR)

DACR is an 8-bit readable/writable register that controls the operation of the D/A converter.

DACR is initialized to H'1F by a power-on reset in deep standby mode or module standby mode.

23.4 Operation

The D/A converter includes D/A conversion circuits for two channels, each of which can operate independently. When the DAOE bit in DACR is set to 1, D/A conversion is enabled and the conversion result is output.

An operation example of D/A conversion on channel 0 is shown below. Figure 23.2 shows the timing of this operation.

- 1. Write the conversion data to DADR0.
- 2. Set the DAOE0 bit in DACR to 1 to start D/A conversion. The conversion result is output from the analog output pin DA0 after the conversion time t_{PCONV} has elapsed. The conversion result continues to be output until DADR0 is written to again or the DAOE0 bit is cleared to 0. The output value is expressed by the following formula:

Contents of DADR 256 ⋅ AVref

- 3. If DADR0 is written to again, the conversion is immediately started. The conversion result is output after the conversion time t_{pconv} has elapsed.
- 4. If the DAOE0 bit is cleared to 0, analog output is disabled.

Figure 23.2 Example of D/A Converter Operation

23.5 Usage Notes

23.5.1 Module Standby Mode Setting

Operation of the D/A converter can be disabled or enabled using the standby control register. The initial setting is for operation of the D/A converter to be halted. Register access is enabled by canceling module standby mode. For details, see section 27, Power-Down Modes.

23.5.2 D/A Output Hold Function in Software Standby Mode

When this LSI enters software standby mode with D/A conversion enabled, the D/A outputs are retained, and the analog power supply current is equal to as during D/A conversion. If the analog power supply current needs to be reduced in software standby mode, clear the DAOE0, DAOE1, and DAE bits to 0 to disable the D/A outputs.

23.5.3 D/A Conversion and D/A Output in Deep Standby Mode

When this LSI enters deep standby mode with D/A conversion enabled, the D/A conversion is stopped and thus the D/A outputs are also stopped. Before entering deep standby mode, clear the DAOE0, DAOE1, and DAE bits to 0 to disable the D/A outputs.

23.5.4 Setting Analog Input Voltage

The reliability of this LSI may be adversely affected if the following voltage ranges are exceeded.

1. AVcc and AVss input voltages

Input voltages AVcc and AVss should be PVcc -0.3 V \leq AVcc \leq PVcc and AVss = PVss. Do not leave the AVcc and AVss pins open when the A/D converter or D/A converter is not in use and in software standby mode. When not in use, connect AVcc to the power supply (PVcc) and AVss to the ground (PVss).

2. Setting range of AVref input voltage Set the reference voltage range of the AVref pin as $3.0 \text{ V} \leq \text{AVec}$.

Section 24 I/O Ports

This LSI has six ports: A to F.

All port pins are multiplexed with other pin functions. The functions of the multiplexed pins are selected using the pin function controller (PFC).

Each port is provided with a data register for storing the pin data and a port read register for reading out the pin values.

24.1 Port A

Port A is an I/O port with 32 pins shown in figure 24.1.

Figure 24.1 Port A

24.1.1 Register Configuration

Table 24.1 lists the port A registers.

24.1.2 Port A Data Registers H and L (PADRH and PADRL)

PADRH and PADRL are 16-bit readable/writable registers that store port A data. Bits PA31DR to PA0DR correspond to pins PA31 to PA0, respectively.

If a pin is set to the general output function, the pin will output the value written to the corresponding bit in PADRH or PADRL, and the register value is read from PADRH or PADRL regardless of the state of the pin.

If a pin is set to the general input function, the pin state, not the register value, will be returned if PADRH or PADRL is read. Also, if a value is written to PADRH or PADRL, although the value will actually be written, it will have no influence on the state of the pin. Table 24.2 summarizes the PADRH and PADRL read/write operations.

PADRH and PADRL are initialized to H'0000 by a power-on reset or in deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

• Port A Data Register H (PADRH)

• Port A Data Register L (PADRL)

Table 24.2 Port A Data Registers H and L (PADRH and PADRL) Read/Write Operations

24.1.3 Port A Port Registers H and L (PAPRH and PAPRL)

PAPRH and PAPRL are 16-bit read-only registers in which bits PA31PR to PA0PR correspond to pins PA31 to PA0. PAPRH and PAPRL are always read as the states of the pins regardless of the PFC setting.

• Port A Port Register H (PAPRH)

• Port A Port Register L (PAPRL)

24.2 Port B

Port B is an I/O port with 32 pins shown in figure 24.2.

Figure 24.2 Port B

24.2.1 Register Configuration

Table 24.3 lists the port B registers.

24.2.2 Port B Data Registers H and L (PBDRH and PBDRL)

PBDRH and PBDRL are 16-bit readable/writable registers that store port B data. Bits PB31DR to PB0DR correspond to pins PB31 to PB0, respectively.

If a pin is set to the general output function, the pin will output the value written to the corresponding bit in PBDRH or PBDRL, and the register value is read from PBDRH or PBDRL regardless of the state of the pin.

If a pin is set to the general input function, the pin state, not the register value, will be returned if PBDRH or PBDRL is read. Also, if a value is written to PBDRH or PBDRL, although the value will actually be written, it will have no influence on the state of the pin. Table 24.4 summarizes the PBDRH and PBDRL read/write operations.

PBDRH and PBDRL are initialized to H'0000 by a power-on reset or in deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

• Port B data register H (PBDRH)

• Port B data register L (PBDRL)

Table 24.4 Port B Data Registers H and L (PBDRH and PBDRL) Read/Write Operations

24.2.3 Port B Port Registers H and L (PBPRH and PBPRL)

PBPRH and PBPRL are 16-bit read-only registers in which bits PB31PR to PB0PR correspond to pins PB31 to PB0. PBPRH and PBPRL are always read the states of the pins regardless of the PFC setting.

• Port B Port Register H (PBPRH)

• Port B Port Register L (PBPRL)

24.3 Port C

Port C is an I/O port with 26 pins and is shown in figure 24.3.

Figure 24.3 Port C

24.3.1 Register Configuration

Table 24.5 lists the port C registers.

Table 24.5 Register Configuration

24.3.2 Port C Data Registers H and L (PCDRH and PCDRL)

PCDRH and PCDRL are 16-bit readable/writable registers that store port C data. Bits PC21DR to PC0DR correspond to pins PC21 to PC0, respectively.

If a pin is set to the general output function, the pin will output the value written to the corresponding bit in PCDRH or PCDRL, and the register value is read from PCDRH and PCDRL regardless of the state of the pin.

If a pin is set to the general input function, the pin state, not the register value, will be returned if PCDRH or PCDRL is read. Also, if a value is written to PCDRH or PCDRL, although the value will actually be written, it will have no influence on the state of the pin. Table 24.6 summarizes the PCDRH and PCDRL read/write operations.

Bits 15 to 6 in PCDRH are reserved. These bits are read as 0. The write value should always be 0.

PCDRH and PCDRL are initialized to H'0000 by a power-on reset or in deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

• Port C data register H (PCDRH)

• Port C data register L (PCDRL)

Table 24.6 Port C Data Registers H and L (PCDRH and PCDRL) Read/Write Operations

24.3.3 Port C Port Registers H and L (PCPRH and PCPRL)

PCPRH and PCPRL are 16-bit read-only registers in which bits PC25PR to PC0PR correspond to pins PC25 to PC0. PCPRH and PCPRL are always read as the states of the pins regardless of the PFC setting.

Bits 15 to 10 in PCPRH are reserved. These bits are read as 0. The write value should always be 0.

• Port C Port Register H (PCPRH)

• Port C Port Register L (PCPRL)

24.4 Port D

Port D is an I/O port with 17 pins shown in figure 24.4.

Figure 24.4 Port D

24.4.1 Register Configuration

Table 24.7 lists the port D registers.

Table 24.7 Register Configuration

24.4.2 Port D Data Register (PDDR)

PDDR is a 16-bit readable/writable register that stores port D data. Bits PD14DR to PD0DR correspond to pins PD14 to PD0, respectively.

If a pin is set to the general output function, that pin will output the value written to the corresponding bit in PDDR, and the register value is read from PDDR regardless of the state of the pin.

If a pin is set to the general input function, the pin state, not the register value, will be returned if PDDR is read. Also, if a value is written to PDDR, although the value will actually be written, it will have no influence on the state of the pin. Table 24.8 summarizes the PDDR read/write operations.

Bit 15 in PDDR is reserved. This bit is read as 0. The write value should always be 0.

PDDR is initialized to H'0000 by a power-on reset or in deep standby mode. This register is not initialized either by a manual reset or by switching to sleep mode or software standby mode.

Table 24.8 Port D Data Register (PDDR) Read/Write Operations

24.4.3 Port D Port Registers H and L (PDPRH and PDPRL)

PDPRH and PDPRL are 16-bit read-only registers in which bits PD16PR to PD0PR correspond to pins PD16 to PD0. PDPRH and PDPRL are always read as the states of the pins regardless of the PFC setting.

Bits 15 to 1 in PDPRH are reserved. These bits are read as 0. The write value should always be 0.

• Port D Port Register H (PDPRH)

• Port D Port Register L (PDPRL)

24.5 Port E

Port E is an I/O port with 8 pins shown in figure 24.5.

24.5.1 Register Configuration

Table 24.9 lists the port E registers.

Table 24.9 Register Configuration

24.5.2 Port E Port Register (PEPR)

PEPR is a 16-bit read-only register. Bits PE7PR to PE0PR correspond to pins PE7 to PE0, respectively. The pin values can always be read from PEPR, regardless of the PFC settings.

24.6 Port F

Port F is an I/O port with 8 pins shown in figure 24.6.

Figure 24.6 Port F

24.6.1 Register Configuration

Table 24.10 lists the port F registers.

Table 24.10 Register Configuration

24.6.2 Port F Data Register (PFDR)

PFDR is a 16-bit read-only register that stores the port F data. Bits PF7DR to PF0DR correspond to pins PF7 to PF0, respectively.

If a pin is set to the general output function, the pin will output the value written to the corresponding bit in PFDR, and the register value is read from PFDR regardless of the state of the pin.

If a pin is set to the general input function, the pin state, not the register value, will be returned if PFDR is read. Also, if a value is written to PFDR, although the value will actually be written, it will have no influence on the state of the pin. Table 24.11 summarizes the PFDR read/write operations.

PFDR is initialized to H'0000 by a power-on reset or in deep standby mode. This register is not initialized either by a manual reset or by switching to sleep mode or software standby mode.

24.6.3 Port F Port Register (PFPR)

PFPR is a 16-bit read-only register in which bits PF7PR to PF0PR correspond to pins PF7 to PF0. PFPR are always read as the states of the pins regardless of the PFC setting.

Section 25 Pin Function Controller (PFC)

The pin function controller (PFC) consists of registers that select the functions of the multiplexed pins and their I/O directions. Tables 25.1 to 25.6 list the multiplexed pins of this LSI.

Table 25.1 Multiplexed Pin Table (Port A)

Function 5 PBnMD[2:0] = 100 (Related modules)

(TMR)

(TMR)

Port (Related modules) Function 1 PBnMD[2:0] = 000 Function 2 PBnMD[2:0] = 001 (Related modules) Function 3 PBnMD[2:0] = 010 (Related modules) Function 4 PBnMD[2:0] = 011 (Related modules) B PB31 I/O (Port) D31 I/O (BSC) PINT7A input (INTC) PB30 I/O (Port) D30 I/O (BSC) PINT6A input (INTC) SCK3 I/O (SCIF) TMCI0 input PB29 I/O (Port) D29 I/O (BSC) PINT5A input (INTC) RxD3 input (SCIF) TMRI0 input PB28 I/O (Port) D28 I/O (BSC) PINT4A input (INTC) PB27 I/O (Port) D27 I/O (BSC) PINT3A input (INTC) PB26 I/O (Port) D26 I/O (BSC) PINT2A input (INTC) TIC5W input (MTU2) PB25 I/O (Port) D25 I/O (BSC) PINT1A input (INTC) TIC5V input (MTU2) PB24 I/O (Port) D24 I/O (BSC) PINT0A input (INTC) TIC5U input (MTU2) PB23 I/O (Port) D23 I/O (BSC) IRQ7A input (INTC) TIOC4D I/O (MTU2)

Table 25.2 Multiplexed Pin Table (Port B)

Table 25.4 Multiplexed Pin Table (Port D)

Port (Related modules) Function 1 PEnMD[2:0] = 000 Function 2 PEnMD[2:0] = 001 (Related modules) Function 3 PEnMD[2:0] = 010 (Related modules) Function 4 PEnMD[2:0] = 011 (Related modules) Function 5 PEnMD[2:0] = 100 (Related modules) E PE7 input (Port) IRQ7B input (INTC) PE6 input (Port) IRQ6B input (INTC) PE5 input (Port) IRQ5B input (INTC) PE4 input (Port) IRQ4B input (INTC) PE3 input (Port) PINT7B input (INTC) PE2 input (Port) PINT6B input (INTC) PE1 input (Port) PINT5B input (INTC)

Table 25.5 Multiplexed Pin Table (Port E)

Table 25.6 Multiplexed Pin Table (Port F)

PE0 input (Port) PINT4B input (INTC)

25.1 Register Descriptions

The PFC includes the following registers.

Table 25.7 Register Configuration

25.1.1 Port A I/O Registers H and L (PAIORH and PAIORL)

PAIORH and PAIORL are 16-bit readable/writable registers that select the I/O direction for the port A pins. Bits PA31IOR to PA0IOR correspond to pins PA31 to PA0, respectively. PAIORH and PAIORL are enabled when the function of the port A pins is set to general-purpose I/O (PA31 to PA0) by PACR, and are disabled in other cases. When a bit in PAIORH and PAIORL is set to 1, the corresponding pin is set to output, and when set to 0, the pin is set to input.

PAIORH and PAIORL are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port A I/O Register H (PAIORH)

(2) Port A I/O Register L (PAIORL)

25.1.2 Port A Control Registers 1 to 8 (PACR1 to PACR8)

PACR1 to PACR8 are 16-bit readable/writable registers that select the functions of the multiplexed port A pins. When PINT3B to PINT0B are selected, do not set A input for the same interrupt.

PACR8 and PACR7 are initialized to H'0000 by a power-on reset or by switching to deep standby mode. PACR1 to PACR6 are initialized to H'1111 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port A Control Register 8 (PACR8)

(2) Port A Control Register 7 (PACR7)

(3) Port A Control Register 6 (PACR6)

(4) Port A Control Register 5 (PACR5)

(5) Port A Control Register 4 (PACR4)

(6) Port A Control Register 3 (PACR3)

(7) Port A Control Register 2 (PACR2)

(8) Port A Control Register 1 (PACR1)

25.1.3 Port B I/O Registers H and L (PBIORH and PBIORL)

PBIORH and PBIORL are 16-bit readable/writable registers that select the I/O direction for the port B pins. Bits PB31IOR to PB0IOR correspond to pins PB31 to PB0, respectively. PBIORH and PBIORL are enabled when the function of the port B pins is set to general-purpose I/O (PB31 to PB0) and to TIOC I/O (MTU2) by PBCR, and are disabled in other cases. When a bit in PBIORH and PBIORL is set to 1, the corresponding pin is set to output, and when set to 0, the pin is set to input.

PBIORH and PBIORL are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port B I/O Register H (PBIORH)

(2) Port B I/O Register L (PBIORL)

25.1.4 Port B Control Registers 1 to 8 (PBCR1 to PBCR8)

PBCR1 to PBCR8 are 16-bit readable/writable registers that select the functions of the multiplexed port B pins. When IRQ7A to IRQ0A or PINT7A to PINT0A are selected, do not set B input for the same interrupt.

PBCR1 and PBCR2 are initialized to H'1111 by a power-on reset or by switching to deep standby mode. PBCR3 to PBCR8 are initialized to the values shown in table 25.8 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

Table 25.8 Port B Control Register Initial Values

(1) Port B Control Register 8 (PBCR8)

(2) Port B Control Register 7 (PBCR7)

(3) Port B Control Register 6 (PBCR6)

(4) Port B Control Register 5 (PBCR5)

(5) Port B Control Register 4 (PBCR4)

Note: * The initial value dependes on the LSI's clock operating mode.

(6) Port B Control Register 3 (PBCR3)

Note: * The initial value dependes on the LSI's clock operating mode.

(7) Port B Control Register 2 (PBCR2)

(8) Port B Control Register 1 (PBCR1)

25.1.5 Port C I/O Registers H and L (PCIORH and PCIORL)

PCIORH and PCIORL are 16-bit readable/writable registers that select the I/O direction for the port C pins. Bits PC21IOR to PC0IOR correspond to pins PC21 to PC0, respectively. PCIORH and PCIORL are enabled when the function of the port C pins is set to general-purpose I/O (PC21 to PC0) and to TIOC I/O (MTU2) by PCCR, and are disabled in other cases. When a bit in PCIORH and PCIORL is set to 1, the corresponding pin is set to output, and when set to 0, the pin is set to input.

Bits 15 to 6 in PCIORH are reserved. These bits are always read as 0. The write value should always be 0.

PCIORH and PCIORL are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port C I/O Register H (PCIORH)

(2) Port C I/O Register L (PCIORL)

25.1.6 Port C Control Registers 1 to 7 (PCCR1 to PCCR7)

PCCR1 to PCCR7 are 16-bit readable/writable registers that select the functions of the multiplexed port C pins. When IRQ3B to IRQ0B are selected, do not set A input for the same interrupt.

PCCR2, PCCR5, PCCR6, and PCCR7 are initialized to H'0000 by a power-on reset or by switching to deep standby mode. PCCR1 is initialized to H'0001 and PCCR3 and PCCR4 are initialized to the values shown in table 25.9 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

Table 25.9 Port C Control Register Initial Values

(1) Port C Control Register 7 (PCCR7)

(2) Port C Control Register 6 (PCCR6)

(3) Port C Control Register 5 (PCCR5)

(4) Port C Control Register 4 (PCCR4)

Note: * The initial value depends on the LSI's operating mode.

(5) Port C Control Register 3 (PCCR3)

(6) Port C Control Register 2 (PCCR2)

(7) Port C Control Register 1 (PCCR1)

25.1.7 Port D I/O Register (PDIOR)

PDIOR are 16-bit readable/writable registers that select the I/O direction for the port D pins. Bits PD14IOR to PD0IOR correspond to pins PD14 to PD0, respectively. PDIOR are enabled when the function of the port D pins is set to general-purpose I/O (PD14 to PD0) and to TIOC I/O (MTU2) by PDCR, and are disabled in other cases. When a bit in PDIOR is set to 1, the corresponding pin is set to output, and when set to 0, the pin is set to input.

Bits 15 in PDIOR is reserved. These bits are always read as 0. The write value should always be 0.

PDIOR are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

25.1.8 Port D Control Registers 1 to 5 (PDCR1 to PDCR5)

PDCR1 to PDCR5 are 16-bit readable/writable registers that select the functions of the multiplexed port D pins.

PDCR1 to PDCR5 are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port D Control Register 5 (PDCR5)

(2) Port D Control Register 4 (PDCR4)

(3) Port D Control Register 3 (PDCR3)

(4) Port D Control Register 2 (PDCR2)

(5) Port D Control Register 1 (PDCR1)

25.1.9 Port E Control Registers 1 and 2 (PECR1 and PECR2)

PECR1 and PECR2 are 16-bit readable/writable registers that select the functions of the multiplexed port E pins. The pins states are set by the corresponding module for A/D converter input and for D/A converter output. When IRQ7B to IRQ4B or PINT7B to PINT4B are selected, do not set A input for the same interrupt.

PECR1 and PECR2 are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port E Control Register 2 (PECR2)

(2) Port E Control Register 1 (PECR1)

25.1.10 Port F I/O Register (PFIOR)

PFIOR is a 16-bit readable/writable register that selects the I/O direction for the port F pins. Bits PF7IOR to PF0IOR correspond to pins PF7 to PF0, respectively. PFIOR is enabled when the function of the port F pins is set to general-purpose I/O (PF7 to PF0) by PFCR, and are disabled in other cases. When a bit in PFIOR is set to 1, the corresponding pin is set to output, and when set to 0, the pin is set to input.

Bits 15 to 8 in PFIOR are reserved. These bits are always read as 0. The write value should always be 0.

PFIOR is initialized to H'0000 by a power-on reset or by switching to deep standby mode. This register is not initialized either by a manual reset or by switching to sleep mode or software standby mode.

25.1.11 Port F Control Registers 1 and 2 (PFCR1 and PFCR2)

PFCR1 and PFCR2 are 16-bit readable/writable registers that select the functions of the multiplexed port F pins.

PFCR1 and PFCR2 are initialized to H'0000 by a power-on reset or by switching to deep standby mode. These registers are not initialized either by a manual reset or by switching to sleep mode or software standby mode.

(1) Port F Control Register 2 (PFCR2)

(2) Port F Control Register 1 (PFCR1)

25.2 Usage Note

The settings of the port control registers are used as the output pin select signals, and are not basically used as the input pin select signals. This causes the signals input from the pins to propagate to all the modules having the relevant multiplexed pins. So, unnecessary input signals must be disabled by the settings of the respective modules.

Settings of port control registers are decoded to enable/disable pins IRQ7A to IRQ0A and IRQ7B to IRQ0B or pins PINT7A to PINT0A and PINT7B to PINT0B. Be sure to select either one of them.

Section 26 On-Chip RAM

This LSI has an on-chip RAM module that achieves high-speed access and can store instructions or data.

On-chip RAM operation and write access to the RAM can be enabled or disabled through the RAM enable bits and RAM write enable bits.

26.1 Features

• Pages

Two pages (pages 0 and 1) are provided.

• Memory map

The on-chip RAM is located in the address spaces shown in table 26.1.

Table 26.1 On-Chip RAM Address Spaces

• Ports

Each page has two independent read and write ports and is connected to the internal bus (I bus), CPU instruction fetch bus (F bus), and CPU memory access bus (M bus). (Note that the F bus is connected only to the read ports.)

The F bus and M bus are used for access by the CPU, and the I bus is used for access by the DMAC via the internal DMA write bus/internal DMA read bus and bus bridge.

• Priority

When requests for access to the same page from different buses coincide, the access is processed in priority order. The priority is I bus $> M$ bus $> F$ bus.

26.2 Usage Notes

26.2.1 Page Conflict

When the same page is accessed from different buses simultaneously, a conflict on the page occurs. Although each access is completed correctly, this kind of conflict degrades the memory access speed. Therefore, it is advisable to provide software measures to prevent such conflicts as far as possible. For example, no conflict will arise if different memory modules or different pages are accessed by each bus.

26.2.2 RAME and RAMWE Bits

Before disabling memory operation or write access through the RAME or RAMWE bit, be sure to read from any address and then write to the same address in each page; otherwise, the last written data in each page may not be actually written to the RAM.

```
 // For RAM page 0 
  MOV.L #H'FFF80000,R0 
  MOV.L @R0,R1 
  MOV.L R1,@R0 
// For RAM page 1 
  MOV.L #H'FFF84000,R0 
  MOV.L @R0,R1 
  MOV.L R1,@R0
```
Figure 26.1 Examples of Read/Write before Disabling RAM

Section 27 Power-Down Modes

This LSI supports sleep mode, software standby mode, deep standby mode, and module standby mode. In power-down modes, functions of CPU, clocks, on-chip memory, or part of on-chip peripheral modules are halted or the power-supply is turned off, through which low power consumption is achieved. These modes are canceled by a reset or interrupt.

27.1 Features

27.1.1 Power-Down Modes

This LSI has the following power-down modes and function:

- 1. Sleep mode
- 2. Software standby mode
- 3. Deep standby mode
- 4. Module standby function

Table 27.1 shows the transition conditions for entering the modes from the program execution state, as well as the CPU and peripheral module states in each mode and the procedures for canceling each mode.

Table 27.1 States of Power-Down Modes

Notes: 1. The pin state is retained or set to high impedance. For details, see appendix A, Pin States.

- 2. RTC operates when the START bit in the RCR2 register is set to 1. For details, see section 15, Realtime Clock (RTC).
- 3. Setting bits RAMKP3 to RAMKP0 in the RAMKP register to 1 enables the retention of data in the corresponding area in the on-chip RAM during the transition to deep standby mode. However, when deep standby mode is canceled by a power-on reset, the contents in the corresponding on-chip RAM area are not retained.
- 4. Deep standby mode can be canceled by an interrupt (NMI or IRQ) or a reset (manual reset or power-on reset). However, IRQ is reset only by PE7 to PE0 and PC25 to PC22. When deep standby mode is canceled by NMI interrupt or IRQ interrupt, reset exception handling is executed instead of interrupt exception handling. These are power-on reset exception handlings including a manual reset.

27.2 Register Descriptions

The following registers are used in power-down modes.

Table 27.2 Register Configuration

27.2.1 Standby Control Register (STBCR)

STBCR is an 8-bit readable/writable register that specifies the state of the power-down mode. This register is initialized to H'00 by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

Note: When writing to this register, see section 27.4, Usage Note.

[Legend]

x: Don't care

27.2.2 Standby Control Register 2 (STBCR2)

STBCR2 is an 8-bit readable/writable register that controls the operation of modules in powerdown modes. STBCR2 is initialized to H'1E by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

27.2.3 Standby Control Register 3 (STBCR3)

STBCR3 is an 8-bit readable/writable register that controls the operation of modules in powerdown modes. STBCR3 is initialized to H'3F by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

27.2.4 Standby Control Register 4 (STBCR4)

STBCR4 is an 8-bit readable/writable register that controls the operation of modules in powerdown modes. STBCR4 is initialized to H'FF by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

27.2.5 Standby Control Register 5 (STBCR5)

STBCR5 is an 8-bit readable/writable register that controls the operation of modules in powerdown modes. STBCR5 is initialized to H'FF by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

27.2.6 System Control Register 1 (SYSCR1)

SYSCR1 is an 8-bit readable/writable register that enables or disables access to the on-chip RAM. SYSCR1 is initialized to H'FF by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

When an RAME bit is set to 1, the corresponding on-chip RAM area is enabled. When an RAME bit is cleared to 0, the corresponding on-chip RAM area cannot be accessed. In this case, an undefined value is returned when reading data or fetching an instruction from the on-chip RAM, and writing to the on-chip RAM is ignored. The initial value of an RAME bit is 1.

Note that when clearing the RAME bit to 0 to disable the on-chip RAM, be sure to execute an instruction to read from or write to the same arbitrary address in each page before setting the RAME bit. If such an instruction is not executed, the data last written to each page may not be written to the on-chip RAM. Furthermore, an instruction to access the on-chip RAM should not be located immediately after the instruction to write to SYSCR1. If an on-chip RAM access instruction is set, normal access is not guaranteed.

7 6 5 4 3 2 1 0

Note: When writing to this register, see section 27.4, Usage Note.

Bit: 7

27.2.7 System Control Register 2 (SYSCR2)

SYSCR2 is an 8-bit readable/writable register that enables or disables write to the on-chip RAM. SYSCR2 is initialized to H'FF by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

When an RAMWE bit is set to 1, the corresponding on-chip RAM area is enabled. When an RAMWE bit is cleared to 0, the corresponding on-chip RAM area cannot be written to. In this case, writing to the on-chip RAM is ignored. The initial value of an RAMWE bit is 1.

Note that when clearing the RAMWE bit to 0 to disable the on-chip RAM, be sure to execute an instruction to read from or write to the same arbitrary address in each page before setting the RAMWE bit. If such an instruction is not executed, the data last written to each page may not be written to the on-chip RAM. Furthermore, an instruction to access the on-chip RAM should not be placed immediately after the instruction to write to SYSCR2. If an on-chip RAM access instruction is set, normal access is not guaranteed.

27.2.8 RAM Retaining Area Specifying Register (RAMKP)

RAMKP is an 8-bit readable/writable register that specifies whether or not to retain data in the corresponding on-chip RAM area in deep standby mode. RAMKP is initialized to H'00 by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

When an RAMKP bit is set to 1, data in the corresponding on-chip RAM area is retained in deep standby mode. When an RAMWE bit is cleared to 0, data in the corresponding on-chip RAM is not retained in deep standby mode.

Deep standby mode is canceled by an interrupt (NMI or IRQ) or a reset (manual reset or power-on reset). However, when deep standby mode is canceled by a power-on reset, the contents in the corresponding on-chip RAM area are not retained even with the RAMKP bit set to 1.

27.2.9 Deep Standby Oscillation Settling Clock Select Register (DSCNT)

DSCNT is an 8-bit readable/writable register that selects the clock used to count the oscillation settling time when the system returns from deep standby mode. DSCNT is initialized to H'00 by a power-on reset or in deep standby mode but retains its previous value by a manual reset or in software standby mode. Only byte access is valid.

Since the frequency control register for the CPG (FRQCR) is initialized in deep standby mode, the frequency of the peripheral clock ($P\phi$) specified by the CKS[2:0] bits in DSCNT is determined by the FRQCR's initial value.

Notes: 1. Do not use this setting.

 2. Set the clock so that it is equal to or longer than the oscillation settling time 2 on return from standby $(t_{\rm osc})$.

27.2.10 Deep Standby Cancel Source Flag Register (DSFR)

DSFR is a 16-bit readable/writable register composed of two types of bits. One is the flag that confirms which interrupt canceled deep standby mode. The other is the bit that releases the retaining state of pins after canceling the deep standby mode. DSFR is initialized to H'0000 by a power-on reset by the RES pin but retains its previous value through a power-on reset caused by a WDT overflow, a manual reset, or a period in software standby mode. When deep standby mode is canceled by interrupts (NMI or IRQ) and a manual reset, this register retains the previous data although power-on reset exception handling is executed. Only word access is valid.

Since interrupt inputs for the NMI and IRQ pins specified by the interrupt controller (INTC) and the pin function controller (PFC) are always detected, these interrupts set flags even during normal operation. Therefore, all flags must be cleared immediately before the transition to deep standby mode.

If an interrupt occurred immediately before executing the SLEEP instruction after the flag clear, the system enters deep standby mode with the flag set again. To prevent this, clear the flag in DSFR even in interrupt exception handling routine.

Note: $*$ Only 0 can be written after reading 1 to clear the flag. Even when IRQ is input after a manual reset has been accepted as a source canceling deep standby, the IRQ flag is not set.

R/W Description

Bit Bit Name

Initial

27.3 Operation

27.3.1 Sleep Mode

(1) Transition to Sleep Mode

Executing the SLEEP instruction when the STBY bit in STBCR is 0 causes a transition from the program execution state to sleep mode. Although the CPU halts immediately after executing the SLEEP instruction, the contents of its internal registers remain unchanged. The on-chip modules continue to run in sleep mode. Clock pulses continue to be output on the CKIO pin.

(2) Canceling Sleep Mode

Sleep mode is canceled by an interrupt (NMI, H-UDI, IRQ, PINT, and on-chip peripheral module), a bus error, or a reset (manual reset or power-on reset).

• Canceling with an interrupt

When an NMI, H-UDI, IRQ, PINT, or on-chip peripheral module interrupt occurs, sleep mode is canceled and interrupt exception handling is executed. When the priority level of the generated interrupt is equal to or lower than the interrupt mask level that is set in the status register (SR) of the CPU, or the interrupt by the on-chip peripheral module is disabled on the module side, the interrupt request is not accepted and sleep mode is not canceled.

• Canceling with a bus error

When a bus error occurs, sleep mode is canceled and bus error exception handling is executed.

Canceling with a reset

Sleep mode is canceled by a power-on reset or a manual reset.

27.3.2 Software Standby Mode

(1) Transition to Software Standby Mode

The LSI switches from a program execution state to software standby mode by executing the SLEEP instruction when the STBY bit and DEEP bit in STBCR are 1 and 0 respectively. In software standby mode, not only the CPU but also the clock and on-chip peripheral modules halt. The clock output from the CKIO pin also halts in clock mode 0 or 2.

The contents of the CPU and cache registers remain unchanged. Some registers of on-chip peripheral modules are, however, initialized. As for the states of on-chip peripheral module registers in software standby mode, see section 30.3, Register States in Each Operating Mode.

The CPU takes one cycle to finish writing to STBCR, and then executes processing for the next instruction. However, it takes one or more cycles to actually write. Therefore, execute a SLEEP instruction after reading STBCR to have the values written to STBCR by the CPU to be definitely reflected in the SLEEP instruction.

The procedure for switching to software standby mode is as follows:

- 1. Clear the TME bit in the WDT's timer control register (WTCSR) to 0 to stop the WDT.
- 2. Set the WDT's timer counter (WTCNT) to 0 and the CKS[2:0] bits in WTCSR to appropriate values to secure the specified oscillation settling time.
- 3. After setting the STBY and DEEP bits in STBCR to 1 and 0 respectively, read STBCR. Then, execute a SLEEP instruction.

(2) Canceling Software Standby Mode

Software standby mode is canceled by interrupts (NMI or IRQ) or a reset (manual reset or poweron reset). In clock modes 0 and 2, a clock signal starts to be output from the CKIO pin.

Canceling with an interrupt

When the falling edge or rising edge of the NMI pin (selected by the NMI edge select bit (NMIE) in interrupt control register 0 (ICR0) of the interrupt controller (INTC)) or the falling edge or rising edge of an IRQ pin (IRQ7 to IRQ0) (selected by the IRQn sense select bits (IRQn1S and IRQn0S) in interrupt control register 1 (ICR1) of the interrupt controller (INTC)) is detected, clock oscillation is started. This clock pulse is supplied only to the oscillation settling counter (WDT) used to count the oscillation settling time.

After the elapse of the time set in the clock select bits (CKS[2:0]) in the watchdog timer control/status register (WTCSR) of the WDT before the transition to software standby mode, the WDT overflow occurs. Since this overflow indicates that the clock has been stabilized, the clock pulse will be supplied to the entire chip after this overflow. Software standby mode is thus cleared and NMI interrupt exception handling (IRQ interrupt exception handling in case of IRQ) is started. However, when the IRQ interrupt priority level is lower than the interrupt mask level set in the status register (SR) of the CPU, the interrupt request is not accepted and software standby mode is not canceled.

When canceling software standby mode by the NMI interrupt or IRQ interrupt, set the CKS[2:0] bits so that the WDT overflow period will be equal to or longer than the oscillation settling time.

The clock output phase of the CKIO pin may be unstable immediately after detecting an interrupt and until software standby mode is canceled. When software standby mode is canceled by the falling edge of the NMI pin, the NMI pin should be high when the CPU enters software standby mode (when the clock pulse stops) and should be low when the CPU returns from software standby mode (when the clock is initiated after the oscillation settling). When software standby mode is canceled by the rising edge of the NMI pin, the NMI pin should be low when the CPU enters software standby mode (when the clock pulse stops) and should be high when the CPU returns from software standby mode (when the clock is initiated after the oscillation settling) (This is the same with the IRQ pin.)

Canceling with a reset

When the $\overline{\text{RES}}$ pin is driven low, software standby mode is canceled and the LSI enters the power-on reset state. After that, if the RES pin is driven high, the power-on reset exception handling is started.

When the MRES pin is driven low, software standby mode is canceled and the LSI enters the manual reset state. After that, if the MRES pin is driven high, the manual reset exception handling is started.

Keep the RES or MRES pin low until the clock oscillation settles. The internal clock will continue to be output to the CKIO pin in clock mode 0 or 2.

(3) Note on Making a Transition To Software Standby Mode

If the SLEEP instruction is executed to make a transition to software standby mode during transfer by the DMAC, the DMAC stops its operation without waiting for the completion of the transfer. Thus, the DMA transfer is not guaranteed. Therefore, when making a transition to software standby mode, wait for the completion of the DMA transfer or stop the DMA transfer to execute the SLEEP instruction.

27.3.3 Software Standby Mode Application Example

This example describes a transition to software standby mode on the falling edge of the NMI signal, and cancellation on the rising edge of the NMI signal. The timing is shown in figure 27.1.

When the NMI pin is changed from high to low level while the NMI edge select bit (NMIE) in ICR is set to 0 (falling edge detection), the NMI interrupt is accepted. When the NMIE bit is set to 1 (rising edge detection) by the NMI exception service routine, the STBY and DEEP bits in STBCR are set to 1 and 0 respectively, and a SLEEP instruction is executed, software standby mode is entered. Thereafter, software standby mode is canceled when the NMI pin is changed from low to high level.

Figure 27.1 NMI Timing in Software Standby Mode (Application Example)

27.3.4 Deep Standby Mode

(1) Transition to Deep Standby Mode

The LSI switches from a program execution state to deep standby mode by executing the SLEEP instruction when the STBY bit and DEEP bit in STBCR are set to 1. In deep standby mode, not only the CPU, clocks, and on-chip peripheral modules but also power supply is turned off excluding the on-chip RAM retaining area specified by the RAMKP3 to RAMKP0 bits in the RAMKP register and RTC. This can significantly reduce power consumption. Therefore, data in the registers of the CPU, cache, and on-chip peripheral modules are not retained. Pin state values immediately before the transition to deep standby mode can be retained.

The CPU takes one cycle to finish writing to DSFR, and then executes processing for the next instruction. However, it actually takes one or more cycles to write. Therefore, execute a SLEEP instruction after reading DSFR to have the values written to DSFR by the CPU to be definitely reflected in the SLEEP instruction.

The procedure for switching to deep standby mode is as follows. Figure 27.2 also shows its flowchart.

- 1. Set the RAMKP3 to RAMKP0 bits in the RAMKP register for the corresponding on-chip RAM retaining area.
- 2. Execute read and write of an arbitrary but the same address for each page in the retaining RAM area. When this is not executed, data last written may not be stored in the on-chip RAM. If there is a write to the on-chip RAM after this time, execute this processing after the last write to the on-chip RAM.
- 3. Set the CKS[2:0] bits in the DSCNT register so that the initial value of FRQCR in the CPG becomes larger than the oscillation settling time.
- 4. Set the STBY and DEEP bits in the STBCR register to 1.
- 5. Read out the DSFR register after clearing the flag in the DSFR register. Then execute the SLEEP instruction.

Figure 27.2 Flowchart of Transition to Deep Standby Mode

(2) Canceling Deep Standby Mode

Deep standby mode is canceled by interrupts (NMI or IRQ) or a reset (manual reset or power-on reset). However, IRQ is canceled only by PE7 to PE4 and PC25 to PC22. To cancel deep standby mode by interrupt NMI or IRQ, a power-on reset exception handling instead of an interrupt exception handling is executed. In the same way, a power-on reset exception handling is executed by a power-on reset. Figure 27.3 shows the flowchart of canceling deep standby mode.

Figure 27.3 Flowchart of Canceling Deep Standby Mode

Canceling with an interrupt

When the falling edge or rising edge of the NMI pin (selected by the NMI edge select bit (NMIE) in interrupt control register 0 (ICR0) of the interrupt controller (INTC)) or the falling edge or rising edge of an IRQ pin (IRQ7 to IRQ0: PE7 to PE4 and PC25 to PC22) (selected by the IRQn sense select bits (IRQn1S and IRQn0S) in interrupt control register 1 (ICR1) of the interrupt controller (INTC)) is detected, clock oscillation is started after the wait time for the oscillation settling time. This clock pulse is supplied only to the oscillation settling counter (DSCNT) used to count the oscillation settling time.

After the elapse of the time set in the clock select bits (CKS[2:0]) in DSCNT before the transition to deep standby mode, an overflow occurs. Since this overflow indicates that the clock has been stabilized, the clock pulse will be supplied to the entire chip after this overflow. Deep standby mode is thus cleared and reset exception handling is started.

When canceling deep standby mode by the NMI interrupt or IRO interrupt, set the CKS[2:0] bits so that the overflow period will be equal to or longer than the oscillation settling time.

The clock output phase of the CKIO pin may be unstable immediately after detecting an interrupt and until deep standby mode is canceled. When deep standby mode is canceled by the falling edge of the NMI pin, the NMI pin should be high when the CPU enters deep standby mode (when the clock pulse stops) and should be low when the CPU returns from deep standby mode (when the clock is initiated after the oscillation settling). When deep standby mode is canceled by the rising edge of the NMI pin, the NMI pin should be low when the CPU enters deep standby mode (when the clock pulse stops) and should be high when the CPU returns from deep standby mode (when the clock is initiated after the oscillation settling). (This is the same with the IRQ pin.)

Canceling with a reset

When the $\overline{\text{RES}}$ or $\overline{\text{MRES}}$ pin is driven low, this LSI enters the power-on reset state and deep standby mode is canceled.

Keep the RES or MRES pin low until the clock oscillation settles. When deep standby mode is canceled by the RES pin, the contents in the on-chip RAM area are not retained.

(3) Operation after Canceling Deep Standby Mode

When deep standby mode is canceled by interrupts (NMI or IRQ) or a manual reset, the deep standby cancel source flag register (DSFR) can be used to confirm which interrupt has canceled the mode.

Pins retain the state immediately before the transition to deep standby mode. However, in canceling deep standby mode, only the pins in buses listed in the table 27.3 can fetch programs while canceling pin states. Pins other than those retain the pin states after canceling deep standby mode, in which DSFR can confirm which interrupt has triggered returning to deep standby mode. Reconfiguration of peripheral functions is required to return to the previous state of deep standby mode. Peripheral functions include every function such as CPG, INTC, BSC, I/O ports, PFC, and peripheral modules. After the reconfiguration, pin-retaining state can be canceled by reading 1 in the IOKEEP bit of DSFR then writing 0 to it.

(4) Note on Making a Transition To Deep Standby Mode

If the SLEEP instruction is executed to make a transition to deep standby mode during transfer by the DMAC, the DMAC stops its operation without waiting for the completion of the transfer. Thus, the DMA transfer is not guaranteed. Therefore, when making a transition to deep standby mode, wait for the completion of the DMA transfer or stop the DMA transfer to execute the SLEEP instruction.

27.3.5 Module Standby Function

(1) Transition to Module Standby Function

Setting the standby control register MSTP bits to 1 halts the supply of clocks to the corresponding on-chip peripheral modules. This function can be used to reduce the power consumption in normal mode and sleep mode. Disable a module before placing it in the module standby mode. In addition, do not access the module's registers while it is in the module standby state.

The register states are the same as those in software standby mode.

However, the state of DAC registers are exceptional. In the DAC, all registers retain their previous values in software standby mode, but are initialized in module standby mode.

(2) Canceling Module Standby Function

The module standby function can be canceled by clearing the MSTP bits to 0, or by a power-on reset (only possible for RTC, H-UDI, UBC, DMAC, and AUD-II). When taking a module out of the module standby state by clearing the corresponding MSTP bit to 0, read the MSTP bit to confirm that it has been cleared to 0.

27.4 Usage Note

27.4.1 Note on Setting Registers

When writing to the registers related to power-down modes, note the following.

When writing to the register related to power-down modes, the CPU, after executing a write instruction, executes the next instruction without waiting for the write operation to complete. Therefore, to reflect the change specified by writing to the register while the next instruction is executed, insert a dummy read of the same register between the register write instruction and the next instruction.

27.4.2 Note on Canceling Standby Mode when an External Clock is being Input

When release from standby mode is initiated by an interrupt (NMI or IRQ) while an external clock from the EXTAL pin or CKIO pin is in use, make sure that the external clock is being input before input of the interrupt. If this is not the case, correct counting of the oscillation settling time will not be possible.

Section 28 User Debugging Interface (H-UDI)

This LSI incorporates a user debugging interface (H-UDI) for emulator support.

28.1 Features

The user debugging interface (H-UDI) has reset and interrupt request functions.

The H-UDI in this LSI is used for emulator connection. Refer to the emulator manual for the method of connecting the emulator.

Figure 28.1 shows a block diagram of the H-UDI.

Figure 28.1 Block Diagram of H-UDI

28.2 Input/Output Pins

Table 28.1 Pin Configuration

Note: $*$ The pin with the pull-up function.
28.3 Register Descriptions

The H-UDI has the following registers.

Table 28.2 Register Configuration

28.3.1 Bypass Register (SDBPR)

SDBPR is a 1-bit register that cannot be accessed by the CPU. When SDIR is set to BYPASS mode, SDBPR is connected between H-UDI pins UDTDI and UDTDO. The initial value is undefined.

28.3.2 Instruction Register (SDIR)

SDIR is a 16-bit read-only register. It is initialized by UDTRST assertion, in the TAP test-logicreset state or in deep standby mode, and can be written to by the H-UDI irrespective of the CPU mode. Operation is not guaranteed if a reserved command is set in this register. The initial value is H'EFFD.

Note: * The initial value of the TI[7:0] bits is a reserved value. When setting a command, the TI[7:0] bits must be set to another value.

Note: * The initial value of the TI[7:0] bits is a reserved value. When setting a command, the TI[7:0] bits must be set to another value.

Table 28.3 H-UDI Commands

28.4 Operation

28.4.1 TAP Controller

Figure 28.2 shows the internal states of the TAP controller.

Figure 28.2 TAP Controller State Transitions

Note: The transition condition is the UDTMS value at the rising edge of UDTCK. The UDTDI value is sampled at the rising edge of UDTCK; shifting occurs at the falling edge of UDTCK. For details on change timing of the UDTDO value, see section 28.4.3, UDTDO Output Timing. The UDTDO is at high impedance, except with shift-DR and shift-IR states. There is a transition to test-logic-reset asynchronously with UDTCK by UDTRST assertion or deep standby mode.

28.4.2 Reset Types

Table 28.4 Reset Types

Note: * Fix ASEMD to high.

28.4.3 UDTDO Output Timing

The initial value of the UDTDO change timing is to perform data output from the UDTDO pin on the UDTCK falling edge. However, setting a UDTDO change timing switch command in SDIR via the H-UDI pin and passing the Update-IR state synchronizes the UDTDO change timing to the UDTCK rising edge. Hereafter, to synchronize the UDTDO change timing with the UDTCK falling edge, the UDTRST pin must be asserted simultaneously with a power-on reset or deep standby mode must be entered. In the case of a power-on reset by the RES pin, the LSI falls in reset state for a certain period after the RES pin negation. Therefore, when the UDTRST pin is asserted immediately after the RES pin negation, a UDTDO change timing switch command is cleared and the UDTDO change timing becomes synchronous with the output of UDTCK falling edge. To prevent this, at least 20 tcyc must be set between the change timings of the RES pin and UDTRST pin.

Figure 28.3 H-UDI Data Transfer Timing

28.4.4 H-UDI Reset

An H-UDI reset is executed by setting an H-UDI reset assert command in SDIR. An H-UDI reset is of the same kind as a power-on reset. An H-UDI reset is released by setting an H-UDI reset negate command. The required time between the H-UDI reset assert command and H-UDI reset negate command is the same as time for keeping the RES pin low to apply a power-on reset.

Figure 28.4 H-UDI Reset

28.4.5 H-UDI Interrupt

The H-UDI interrupt function generates an interrupt by setting a command from the H-UDI in SDIR. An H-UDI interrupt is a general exception/interrupt operation, resulting in fetching the exception service routine start address from the exception handling vector table, jumping to that address, and starting program execution from that address. This interrupt request has a fixed priority level of 15.

H-UDI interrupts are accepted in sleep mode, but not in software standby or deep standby mode.

28.5 Usage Notes

- 1. An H-UDI command, once set, will not be modified as long as another command is not set again from the H-UDI. If the same command is to be set continuously, the command must be set after a command (BYPASS mode, etc.) that does not affect chip operations is once set.
- 2. In software standby mode or H-UDI module standby state, all of the functions in the H-UDI cannot be used. To retain the TAP status before and after software standby mode or H-UDI module standby state, keep UDTCK high before entering software standby mode or H-UDI module standby state.
- 3. In deep standby mode, all of the functions in the H-UDI cannot be used. H-UDI is initialized in deep standby mode.
- 4. If the UDTRST pin is asserted immediately after the setting of the UDTDO transition timing switching command and the negation of the RES pin, the UDTDO transition timing switching command is cleared. To avoid this case, make sure to put 20 tcyc or longer between the signal transition timing of the RES and UDTRST pins. For details, see section 28.4.3, UDTDO Output Timing.
- 5. When starting the TAP controller after the negation of the UDTRST pin, make sure to allow 200 ms or more after the negation.

Section 29 Advanced User Debugger II (AUD-II)

The AUD-II offers functions that support user program debugging with the LSI mounted and operated in actual performance. Use of the AUD-II simplifies the construction of a simple emulator, with functions such as monitoring/tuning of on-chip RAM data.

29.1 Features

The AUD-II can be used in RAM monitor mode by setting AUDMD.

RAM monitor mode:

- Functions to read/write modules connected to internal/external buses (except cache and H-UDI)
- Outputs data corresponding to an address that is externally written to AUDATA
- Transmits data to the address in AUDATA to which address and data are written

29.2 Input/Output Pins

(1) Description of Pins

Table 29.2 Description of Pins

29.3 RAM Monitor Mode

In this mode, all the modules connected to this LSI's internal or external bus can be read and written to (except cache and H-UDI), allowing RAM monitoring and tuning to be carried out.

29.3.1 Communication Protocol

The AUD-II latches the AUDATA input when AUDSYNC is asserted. The following AUDATA input format should be used.

Figure 29.1 AUDATA Input Format

29.3.2 Operation

Operation starts in RAM monitor mode when AUDRST is asserted, AUDMD is driven high, and then AUDRST is negated.

Figure 29.2 shows an example of a read operation, and figure 29.3 shows an example of a write operation.

When AUDSYNC is asserted, input from the AUDATA pins begins. When a command, address, or data (writing only) is input in the format shown in figure 29.1, execution of read/write access to the specified address is started. During internal execution, the AUD returns Not Ready (B'0000). When execution is completed, the Ready flag (B'0001) is returned (figures 29. 2 and 29. 3). Table 29.3 shows the Ready flag format.

In a read, data of the specified size is output when AUDSYNC is negated following detection of this flag (figure 29. 2).

If a command other than the above is input in DIR, the AUD-II treats this as a command error, disables processing, and sets bit 1 in the Ready flag to 1. If a read/write operation initiated by the command specified in DIR causes a bus error, the AUD-II disables processing and sets bit 2 in the Ready flag to 1 (figure 29. 4).

Bus error conditions are shown below.

- 1. Word access to address 4n+1 or 4n+3
- 2. Longword access to address 4n+1, 4n+2, or 4n+3

Table 29.3 Ready Flag Format

Figure 29.3 Example of Write Operation (Longword Write)

Figure 29.4 Example of Error Occurrence (Longword Read)

29.3.3 Usage Notes (RAM Monitor Mode)

(1) Guidelines for initialization of the RAM monitor mode

The buffers in this debugger and the processing status are initialized under the following conditions.

- Power-on reset
- When the AUDRST pin is driven low
- Module standby
- Deep standby mode

(2) Guidelines for AUDCK

• AUDCK is for inputting the external clock. Input the clock to satisfy $B\phi/2 \geq \text{AUDCK}$.

(3) Other Limitations

- Do not negate $\overline{\text{AUDSYNC}}$ until the command is input to AUDATA and the Ready is returned.
- The RAM monitor functions in sleep mode but is not available in software standby or deep standby mode.

Section 30 List of Registers

The address map gives information on the on-chip I/O registers and is configured as described below.

- 1. Register Addresses (address order)
- Registers are described by functional module, in order of the corresponding section numbers.
- Access to reserved addresses that are not described in this register address list is prohibited.
- When addresses consist of 16 or 32 bits, the addresses of the MSBs are given when big-endian mode is selected.
- 2. Register Bits
- Bit configurations of the registers are described in the same order as the Register Addresses (address order).
- Reserved bits are indicated by in the bit name.
- No entry in the bit-name column indicates that the whole register is allocated as a counter or for holding data.
- When registers consist of 16 or 32 bits, the bits are given from the MSB side. The listing order of bytes is based on big-endian mode.
- 3. Register States in Each Operating Mode
- Register states are described in the same order as the Register Addresses (address order).
- For the initial state of each bit, refer to the description of the register in the corresponding section.
- The register states described are for the basic operating modes. If there is a specific reset for an on-chip peripheral module, refer to the section on that on-chip peripheral module.

30.1 Register Addresses (Address Order)

Entries under Access Size indicate numbers of bits.

Note: Access to undefined or reserved addresses is prohibited. Since operation or continued operation is not guaranteed when these registers are accessed, do not attempt such access.

Page 1178 of 1336 RO1UH0025EJ0300 Rev. 3.00

30.2 Register Bits

Register addresses and bit names of the on-chip peripheral modules are described below.

Each line covers eight bits, and 16-bit and 32-bit registers are shown as 2 or 4 lines, respectively.

⎯ ⎯ ⎯ ⎯ ⎯ ⎯ ⎯ ⎯

e e

SH7261 Group

Section 30 List of Registers

IPR08 IP33 IP32 IP31 IP30 IP23 IP22 IP21 IP20

IPR09 IP33 IP32 IP31 IP30 IP23 IP22 IP21 IP20

IPR10 IP33 IP32 IP31 IP30 IP23 IP22 IP21 IP20

IPR11 IP33 IP32 IP31 IP30 IP23 IP22 IP21 IP20

IPR12 IP33 IP32 IP31 IP30 IP23 IP22 IP21 IP20

IPR13 IP33 IP32 IP31 IP30 IP23 IP22 IP21 IP20

IP13 IP12 IP11 IP10 IP03 IP02 IP01 IP00

16/8/0 Module

Page 1210 of 1336 RO1UH0025EJ0300 Rev. 3.00 Sep 24, 2010

Section 30 List of Registers

SH7261 Group

SH7261 Group

SH7261 Group

—

j.

j.

l.

j.

j.

j.

Sep 24, 2010

30.3 Register States in Each Operating Mode

Section 30 List of Registers

Notes: 1. Not initialized in deep standby mode. But initialized after deep standby mode is released because a power-on reset exception handling is executed.

2. Initialized by UDTRST assertion or in the Test-Logic-Reset state of the TAP controller.

- 3. Bits BN[3:0] are initialized.
- 4. Retains the previous value after an internal power-on reset by means of the WDT.
- 5. Counting up continues.
- 6. Bits RTCEN and START are retained.
- 7. Bits BC[3:0] are initialized.
- 8. The ENB bit is initialized.

Section 31 Electrical Characteristics

31.1 Absolute Maximum Ratings

Table 31.1 lists the absolute maximum ratings.

Table 31.1 Absolute Maximum Ratings

Caution: Permanent damage to the LSI may result if absolute maximum ratings are exceeded.

31.2 DC Characteristics

Tables 31.2 and 31.3 list DC characteristics.

Table 31.2 DC Characteristics (1) [Common Items] [Regular Specifications]

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{cc}} - 0.3 \text{ V} \leq AV_{\text{cc}} \leq PV_{\text{cc}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{cc} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Caution: When the A/D converter or D/A converter is not in use, the AV_{cc} and AV_{ss} pins should not be open.

Note: * Supply current values are values when all of the output pins and pins with the pull-up function (UDTRST, UDTMS, UDTDI, UDTCK, ASEBRK/ASEBRKAK) are unloaded and represent the total current supplied to the PV_{cc} , $V_{\text{cc}}R$, and $PLLV_{\text{cc}}$ systems. Reference values are given under "Typ."

Table 31.2 DC Characteristics (2) [Common Items] [Wide-Range Specifications]

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V \leq AV_{cc} \leq PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Caution: When the A/D converter or D/A converter is not in use, the AV_{cc} and AV_{ss} pins should not be open.

Note: * Supply current values are values when all of the output pins and pins with the pull-up function (UDTRST, UDTMS, UDTDI, UDTCK, ASEBRK/ASEBRKAK) are unloaded and represent the total current supplied to the PV_{cc}, V_{cc}R, and PLLV_{cc} systems. Reference values are given under "Typ."

Table 31.2 DC Characteristics (3) [Except for I²C-Related Pins* **1]**

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V \leq AV_{cc} \leq PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

PC24/IRQ2/SCL1, PC25/IRQ3/SDA1, PD15/SDA2, and PD16/SCL2

2. Except (PC22/)IRQ0, (PC23/)IRQ1, (PC24/)IRQ2, and (PC25/)IRQ3

Table 31.2 DC Characteristics (4) [I²C-Related Pins***]**

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{cc} ≤ PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: * Pins (open-drain pins): PC22/IRQ0/SCL0/DREQ2, PC23/IRQ1/SDA0, PC24/IRQ2/SCL1, PC25/IRQ3/SDA1, PD15/SDA2, and PD16/SCL2

Table 31.3 Permissible Output Currents (1) [Common Items]

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{CC} ≤ PV_{CC}, AV_{ref} = 3.0 V to AV_{CC}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Caution: To protect the LSI's reliability, do not exceed the output current values in the table above.

Table 31.3 Permissible Output Currents (2) [Wide-Range Specifications]

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{CC} ≤ PV_{CC}, AV_{ref} = 3.0 V to AV_{CC}, $PV_{\text{ss}} = V_{\text{ss}}R = PLLV_{\text{ss}} = AV_{\text{ss}} = 0 V$, $I\phi \le 80 MHz$

Caution: To protect the LSI's reliability, do not exceed the output current values in the table above.

Table 31.3 Permissible Output Currents (3) [Wide-Range Specifications]

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{cc} ≤ PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$, 80 MHz < I $\phi \le 100$ MHz

Caution: To protect the LSI's reliability, do not exceed the output current values in the table above.

SH7261 Group

31.3 AC Characteristics

Signals input to this LSI are basically handled as signals in synchronization with a clock. The setup and hold times for input pins must be followed.

Table 31.4 Maximum Operating Frequency

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{CC} ≤ PV_{CC}, AV_{ref} = 3.0 V to AV_{CC}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

31.3.1 Clock Timing

Table 31.5 Clock Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{cc} ≤ PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Figure 31.1 EXTAL, AUDIO_X1, and AUDIO_CLK Clock Input Timing

Figure 31.2 CKIO Clock Input Timing

Figure 31.3 CKIO Clock Output Timing

Figure 31.4 Power-On Oscillation Settling Time

Figure 31.5 Oscillation Settling Time on Return from Standby (Return by Reset)

Figure 31.6 Oscillation Settling Time on Return from Standby (Return by NMI or IRQ)

Figure 31.7 RTC Clock Oscillation Settling Time

31.3.2 Control Signal Timing

Table 31.6 Control Signal Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Notes: 1. The RES, MRES, NMI, IRQ7 to IRQ0 and PINT7 to PINT0 signals are asynchronous signals. When the setup time is satisfied, change of signal level is detected at the rising edge of the clock. If not, the detection can be delayed until the rising edge of the next clock.

- 2. In software standby mode, deep standby mode or when the clock multiplication ratio is changed, $t_{\text{new}} = t_{\text{osc}}$ (min).
- 3. In software standby mode or deep standby mode, $t_{MRESW} = t_{OSC2}$ (min).
- 4. In software standby mode or deep standby mode, $t_{\text{mmv}}/t_{\text{ROW}} = t_{\text{osc3}}$ (min).

Figure 31.8 Reset Input Timing

Figure 31.9 Interrupt Signal Input Timing

31.3.3 Bus Timing

Table 31.7 Bus Timing* 1

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Notes: 1. When writing to the external address space or making SDRAM settings in power-on reset exception handling or cancellation of deep standby mode, be sure to set bits ACOSW[3:0] in ACSWR to B'0011 beforehand.

2. The maximum value (f_{max}) of B ϕ (bus clock) depends on the number of wait cycles and the system configuration of your board.

Figure 31.10 (1) External Address Space: Basic Bus Timing (Normal Access, Read/Write Cycle Wait = 3, CS Assert Wait = 1, Write Data Output Wait = 1, WR/RD Assert Wait = 2, Write Data Output Delay Cycles = 0, Read/Write CS Delay Cycles = 1)

Figure 31.10 (2) External Address Space: Basic Bus Timing (Normal Access, Data Recovery Cycles = 0, Read/Write Cycle Wait = 1, Read/Write CS Delay Cycles = 1, Other Wait Settings = 0)

Figure 31.11 External Address Space: Basic Bus Timing (Page Read Access, Normal Access Compatible Mode , Read Cycle Wait = 2, Page Read Cycle Wait = 2, CS Assert Wait = 1, RD Assert Wait = 1, Read CS Delay Cycles = 1)

Figure 31.12 External Address Space: Basic Bus Timing (Page Read Access, External Read Data Continuous Assert Mode, Read Cycle Wait = 2, Page Read Cycle Wait = 1, CS Assert Wait = 1, RD Assert Wait = 1, Read CS Delay Cycles = 1)

Figure 31.13 External Address Space: Basic Bus Timing (Page Write Access, Write Cycle Wait = 2, CS Assert Wait = 1, WR Assert Wait = 1, Write Data Output Delay Cycles = 1, Other Wait Settings = 0)

Figure 31.14 External Address Space: Timing with External Wait (Page Read Access to 16-Bit Width Channel, External Read Data Continuous Assert Mode, Read Cycle Wait = 3, Page Read Cycle Wait = 3, Other Wait Settings = 0, External Wait Cycles = 2)

Figure 31.15 Single Read Bus Timing for SDRAM Space (DCL = 2 (Two Cycles), DRCD = 1 (Two Cycles), DPCG = 1 (Two Cycles))

Figure 31.16 Single Write Bus Timing for SDRAM Space (DCL = 2 (Two Cycles), DRCD = 1 (Two Cycles), DPCG = 1 (Two Cycles))

Figure 31.17 Multiple Read Bus Timing for SDRAM Space (Four Data Access, DCL = 2 (Two Cycles), DRCD = 1 (Two Cycles), DPCG = 1 (Two Cycles))

Figure 31.18 Multiple Write Bus Timing for SDRAM Space (Four Data Access, DCL = 2 (Two Cycles), DRCD = 1 (Two Cycles), DPCG = 1 (Two Cycles))

Figure 31.19 Multiple Read Row Span Bus Timing for SDRAM Space (Eight Data Access, DCL = 2 (Two Cycles), DRCD = 1 (Two Cycles), DPCG = 1 (Two Cycles))

Figure 31.20 Bus Timing for SDRAM Space Mode Register Setting

Figure 31.21 Bus Timing for SDRAM Space Self Refresh

31.3.4 DMAC Module Timing

Table 31.8 DMAC Module Timing

Figure 31.23 DACK, DACT, DTEND Output Timing

31.3.5 UBC Trigger Timing

Table 31.9 UBC Trigger Timing

CKIO	^I UBCTGD
UBCTRG	

Figure 31.24 UBC Trigger Timing

31.3.6 MTU2 Module Timing

Table 31.10 MTU2 Module Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: $_{\sf pcyc}$ indicates peripheral clock (P \upphi) cycle.

Figure 31.25 MTU2 Input/Output Timing

31.3.7 8-Bit Timer Timing

Table 31.11 8-Bit Timer Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{CC} ≤ PV_{CC}, AV_{ref} = 3.0 V to AV_{CC}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: Above is the case in which the clock ratio B:P = n:1 (n = 1, 2, 3, 4, 6, 8, or 12)

 t_{pcyc} indicates peripheral clock (P ϕ) cycle.

Figure 31.27 8-Bit Timer Output Timing

Figure 31.28 8-Bit Timer Reset Input Timing

31.3.8 Watchdog Timer Timing

Table 31.12 shows the timing of the watchdog timer.

Table 31.12 Watchdog Timer Timing

Figure 31.30 Watchdog Timer Timing

31.3.9 SCIF Module Timing

Table 31.13 SCIF Module Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: t_{pcyc} indicates a peripheral clock (Pφ) cycle.

Figure 31.31 SCK Input Clock Timing

Figure 31.32 SCIF Input/Output Timing in Clocked Synchronous Mode

31.3.10 IIC3 Module Timing

Table 31.14 I²C Bus Interface 3 Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Notes: 1. t_{poyc} indicates the peripheral clock (P ϕ) cycle.

2. Depends on the value of NF2CYC.

3. Indicates the I/O buffer characteristics.

Figure 31.33 I²C Bus Interface 3 Input/Output Timing

31.3.11 SSI Module Timing

Table 31.15 SSI Module Timing

Figure 31.35 SSI Transmit Timing (1)

Figure 31.36 SSI Transmit Timing (2)

Figure 31.37 SSI Receive Timing (1)

31.3.12 RCAN-ET Module Timing [R5S72611] [R5S72613]

Table 31.16 RCAN-ET Module Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{cc} ≤ PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

31.3.13 A/D Trigger Input Timing

Table 31.17 A/D Trigger Input Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: Above is the case in which the clock ratio B:P = n:1 (n = 1, 2, 3, 4, 6, 8, or 12)

Figure 31.41 A/D Converter External Trigger Input Timing

31.3.14 I/O Port Timing

Table 31.18 I/O Port Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{cc} ≤ PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Figure 31.42 I/O Port Timing

31.3.15 H-UDI-Related Pin Timing

Table 31.19 H-UDI-Related Pin Timing

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{CC}} - 0.3 \text{ V} \leq AV_{\text{CC}} \leq PV_{\text{CC}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{CC} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: $*$ Should be greater than the peripheral clock $(P\phi)$ cycle time.

Figure 31.44 UDTRST **Input Timing**

Figure 31.45 H-UDI Data Transfer Timing

31.3.16 AUD-II Timing

Table 31.20 AUD-II Timing

Figure 31.46 AUD Reset Timing

Figure 31.47 RAM Monitor Timing

31.3.17 AC Characteristics Measurement Conditions

- Input signal reference levels: high level = V_{H} min, low level = V_{H} max
- Output signal reference level: PVCC/2 (PVCC = 3.0 to 3.6 V)
- Input pulse level: PVSS to 3.0 V (where RES, MRES, NMI, MD1, MD0, MD_CLK1, MD_CLK0, ASEMD, UDTRST, and Schmitt trigger input pins are within PVSS to PVCC)
- Input rise and fall times: 1 ns

Figure 31.48 Measurement Circuit

31.4 A/D Converter Characteristics

Table 31.21 lists the A/D converter characteristics.

Table 31.21 A/D Converter Characteristics

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, PV_{cc} – 0.3 V ≤ AV_{cc} ≤ PV_{cc}, AV_{ref} = 3.0 V to AV_{cc}, $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0$ V

Notes: 1. Reference values

2. To satisfy the absolute accuracy, the conversion time should be 3.9 μ s or longer.

31.5 D/A Converter Characteristics

Table 31.22 lists the D/A converter characteristics.

Table 31.22 D/A Converter Characteristics

Conditions: $PV_{cc} = V_{cc}R = PLLV_{cc} = 3.0 V$ to 3.6 V, $AV_{cc} = 3.0 V$ to 3.6 V, $PV_{\text{cc}} - 0.3 \text{ V} \leq AV_{\text{cc}} \leq PV_{\text{cc}}$, $AV_{\text{ref}} = 3.0 \text{ V}$ to AV_{cc} , $PV_{ss} = V_{ss}R = PLLV_{ss} = AV_{ss} = 0 V$

Note: * Reference values

31.6 Usage Note

Mount a multilayer ceramic capacitor between a pair of pins PVcc and PVss, VccR and VssR, or PLLVcc and PLLVss as a bypass capacitor. These capacitors must be placed as close as the power supply pins of the LSI. Also, a capacitor must be connected between the VCL and VSS pins to stabilize the power supply voltage that is internally lowered.

Figure 31.49 Example of Externally Allocated Capacitors

Appendix

A. Pin States

Table A.1 Pin States

[Legend]

I: Input

O: Output

- H: High-level output
- L: Low-level output
- Z: High-impedance

K: Input pins become high-impedance, and output pins retain their state.

- Notes: 1. When pins for the connection with a crystal resonator are not used, the EXTAL and AUDIO_X1 pins must be pulled up and the XTAL and AUDIO_X2 pins must be open. The RTC X1 pin must be connected to GND and the RTC X2 must be open.
	- 2. Power-on reset by low-level input to the RES pin. The pin states after a power-on reset by the H-UDI reset assert command or WDT overflow are the same as the initial pin states at normal operation (see section 25, Pin Function Controller (PFC)).
	- 3. IRQ pins that can release deep standby mode are limited to PE7 to PE4 and PC25 to PC22.
	- 4. Z when the TAP controller of the H-UDI is neither the Shift-DR nor Shift-IR state.
	- 5. L when the CKIO output is specified and Z when the CKIO output is stopped with the setting of CKIOCR.

B. Package Dimensions

The package dimension that is shown in the Renesas Semiconductor Package Data Book has priority.

Main Revisions for This Edition

Main Revisions for This Edition

Index

Numerics

A

B

C

D

E

F

G

H

I

J

L

M

N

O

\mathbf{P}

$\overline{\mathbf{Q}}$

$\overline{\mathbf{R}}$

j

S

T

U

V

W

Renesas 32-Bit RISC Microcomputer SH7261 Group User's Manual: Hardware

SALES OFFICES

Renesas Electronics Corporation

http://www.renesas.com

Renesas Electronics America Inc.
2880 Scott Boulevard Santa Clara, CA 95050-2554, U.S.A.
Tel: +1-408-588-6000, Fax: +1-408-588-6130 **Renesas Electronics Canada Limited** 1101 Nicholson Road, Newmarket, Ontario L3Y 9C3, Canada Tel: +1-905-898-5441, Fax: +1-905-898-3220 **Renesas Electronics Europe Limited** Dukes Meadow, Millboard Road, Bourne End, Buckinghamshire, SL8 5FH, U.K Tel: +44-1628-585-100, Fax: +44-1628-585-900 **Renesas Electronics Europe GmbH** Arcadiastrasse 10, 40472 Düsseldorf, Germany Tel: +49-211-65030, Fax: +49-211-6503-1327 **Renesas Electronics (China) Co., Ltd.** 7th Floor, Quantum Plaza, No.27 ZhiChunLu Haidian District, Beijing 100083, P.R.China Tel: +86-10-8235-1155, Fax: +86-10-8235-7679

Renesas Electronics (Shanghai) Co., Ltd. Unit 204, 205, AZIA Center, No.1233 Lujiazui Ring Rd., Pudong District, Shanghai 200120, China Tel: +86-21-5877-1818, Fax: +86-21-6887-7858 / -7898 **Renesas Electronics Hong Kong Limited**

Unit 1601-1613, 16/F., Tower 2, Grand Century Place, 193 Prince Edward Road West, Mongkok, Kowloon, Hong Kong Tel: +852-2886-9318, Fax: +852 2886-9022/9044

Renesas Electronics Taiwan Co., Ltd. 7F, No. 363 Fu Shing North Road Taipei, Taiwan Tel: +886-2-8175-9600, Fax: +886 2-8175-9670

Renesas Electronics Singapore Pte. Ltd. 1 harbourFront Avenue, #06-10, keppel Bay Tower, Singapore 098632 Tel: +65-6213-0200, Fax: +65-6278-8001

Refer to "http://www.renesas.com/" for the latest and detailed information.

Renesas Electronics Malaysia Sdn.Bhd.
Unit 906, Block B, Menara Amcorp, Amcorp Trade Centre, No. 18, Jln Persiaran Barat, 46050 Petaling Jaya, Selangor Darul Ehsan, Malaysia
Tel: +60-3-7955-9390, Fax: +60-3-7955-9510

Renesas Electronics Korea Co., Ltd.
11F., Samik Lavied' or Bldg., 720-2 Yeoksam-Dong, Kangnam-Ku, Seoul 135-080, Korea
Tel: +82-2-558-3737, Fax: +82-2-558-5141

SH7261 Group User's Manual: Hardware

