DP8390D,NS32490D

DP8390D NS32490D NIC Network Interface Controller

Literature Number: SNLS377A

DP8390D/NS32490D NIC Network Interface Controller

General Description

The DP8390D/NS32490D Network Interface Controller (NIC) is a microCMOS VLSI device designed to ease interfacing with CSMA/CD type local area networks including Ethernet, Thin Ethernet (Cheapernet) and StarLAN. The NIC implements all Media Access Control (MAC) layer functions for transmission and reception of packets in accordance with the IEEE 802.3 Standard. Unique dual DMA channels and an internal FIFO provide a simple yet efficient packet management design. To minimize system parts count and cost, all bus arbitration and memory support logic are integrated into the NIC.

The NIC is the heart of a three chip set that implements the complete IEEE 802.3 protocol and node electronics as shown below. The others include the DP8391 Serial Network Interface (SNI) and the DP8392 Coaxial Transceiver Interface (CTI).

Features

- Compatible with IEEE 802.3/Ethernet II/Thin Ethernet/ StarLAN
- Interfaces with 8-, 16- and 32-bit microprocessor systems
- \blacksquare Implements simple, versatile buffer management
- Requires single 5V supply
- Utilizes low power microCMOS process
- \blacksquare Includes
	- Two 16-bit DMA channels
- 16-byte internal FIFO with programmable threshold Ð Network statistics storage
- Supports physical, multicast, and broadcast address filtering
- Provides 3 levels of loopback
- \blacksquare Utilizes independent system and network clocks

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July 1995

3.0 Functional Description

(Refer to Figure ¹)

RECEIVE DESERIALIZER

The Receive Deserializer is activated when the input signal Carrier Sense is asserted to allow incoming bits to be shifted into the shift register by the receive clock. The serial receive data is also routed to the CRC generator/checker. The Receive Deserializer includes a synch detector which detects the SFD (Start of Frame Delimiter) to establish where byte boundaries within the serial bit stream are located. After every eight receive clocks, the byte wide data is transferred to the 16-byte FIFO and the Receive Byte Count is incremented. The first six bytes after the SFD are checked for valid comparison by the Address Recognition Logic. If the Address Recognition Logic does not recognize the packet, the FIFO is cleared.

CRC GENERATOR/CHECKER

During transmission, the CRC logic generates a local CRC field for the transmitted bit sequence. The CRC encodes all fields after the synch byte. The CRC is shifted out MSB first following the last transmit byte. During reception the CRC logic generates a CRC field from the incoming packet. This local CRC is serially compared to the incoming CRC appended to the end of the packet by the transmitting node. If the local and received CRC match, a specific pattern will be generated and decoded to indicate no data errors. Transmission errors result in a different pattern and are detected, resulting in rejection of a packet.

TRANSMIT SERIALIZER

The Transmit Serializer reads parallel data from the FIFO and serializes it for transmission. The serializer is clocked by the transmit clock generated by the Serial Network Interface (DP8391). The serial data is also shifted into the CRC generator/checker. At the beginning of each transmission, the Preamble and Synch Generator append 62 bits of 1,0 preamble and a 1,1 synch pattern. After the last data byte of the packet has been serialized the 32-bit FCS field is shifted directly out of the CRC generator. In the event of a collision the Preamble and Synch generator is used to generate a 32-bit JAM pattern of all 1's

ADDRESS RECOGNITION LOGIC

The address recognition logic compares the Destination Address Field (first 6 bytes of the received packet) to the Physical address registers stored in the Address Register Array. If any one of the six bytes does not match the pre-programmed physical address, the Protocol Control Logic rejects the packet. All multicast destination addresses are filtered using a hashing technique. (See register description.) If the multicast address indexes a bit that has been set in the filter bit array of the Multicast Address Register Array the packet is accepted, otherwise it is rejected by the Protocol Control Logic. Each destination address is also checked for all 1's which is the reserved broadcast address.

FIFO AND FIFO CONTROL LOGIC

The NIC features a 16-byte FIFO. During transmission the DMA writes data into the FIFO and the Transmit Serializer reads data from the FIFO and transmits it. During reception the Receive Deserializer writes data into the FIFO and the DMA reads data from the FIFO. The FIFO control logic is used to count the number of bytes in the FIFO so that after a preset level, the DMA can begin a bus access and write/ read data to/from the FIFO before a FIFO underflow//overflow occurs.

3.0 Functional Description (Continued)

Because the NIC must buffer the Address field of each incoming packet to determine whether the packet matches its Physical Address Registers or maps to one of its Multicast Registers, the first local DMA transfer does not occur until 8 bytes have accumulated in the FIFO.

To assure that there is no overwriting of data in the FIFO, the FIFO logic flags a FIFO overrun as the 13th byte is written into the FIFO; this effectively shortens the FIFO to 13 bytes. In addition, the FIFO logic operates differently in Byte Mode than in Word Mode. In Byte Mode, a threshold is indicated when the $n + 1$ byte has entered the FIFO; thus, with an 8-byte threshold, the NIC issues Bus Request (BREQ) when the 9th byte has entered the FIFO. For Word Mode, BREQ is not generated until the $n + 2$ bytes have entered the FIFO. Thus, with a 4 word threshold (equivalent to an 8-byte threshold), BREQ is issued when the 10th byte has entered the FIFO.

PROTOCOL PLA

The protocol PLA is responsible for implementing the IEEE 802.3 protocol, including collision recovery with random backoff. The Protocol PLA also formats packets during transmission and strips preamble and synch during reception.

DMA AND BUFFER CONTROL LOGIC

The DMA and Buffer Control Logic is used to control two 16-bit DMA channels. During reception, the Local DMA stores packets in a receive buffer ring, located in buffer memory. During transmission the Local DMA uses programmed pointer and length registers to transfer a packet from local buffer memory to the FIFO. A second DMA channel is used as a slave DMA to transfer data between the local buffer memory and the host system. The Local DMA and Remote DMA are internally arbitrated, with the Local DMA channel having highest priority. Both DMA channels use a common external bus clock to generate all required bus timing. External arbitration is performed with a standard bus request, bus acknowledge handshake protocol. -system breational, briticles issued when the Diff. This a b-bit withic hits b-bit wave interest a 64-bit are proposed
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4.0 Transmit/Receive Packet Encapsulation/Decapsulation

A standard IEEE 802.3 packet consists of the following fields: preamble, Start of Frame Delimiter (SFD), destination address, source address, length, data, and Frame Check Sequence (FCS). The typical format is shown in *Figure 2.* The packets are Manchester encoded and decoded by the DP8391 SNI and transferred serially to the NIC using NRZ data with a clock. All fields are of fixed length except for the data field. The NIC generates and appends the preamble, SFD and FCS field during transmission. The Preamble and SFD fields are stripped during reception. (The CRC is passed through to buffer memory during reception.)

PREAMBLE AND START OF FRAME DELIMITER (SFD)

The Manchester encoded alternating 1,0 preamble field is used by the SNI (DP8391) to acquire bit synchronization with an incoming packet. When transmitted each packet contains 62 bits of alternating 1,0 preamble. Some of this preamble will be lost as the packet travels through the network. The preamble field is stripped by the NIC. Byte alignment is performed with the Start of Frame Delimiter (SFD) pattern which consists of two consecutive 1's. The NIC does not treat the SFD pattern as a byte, it detects only the two bit pattern. This allows any preceding preamble within the SFD to be used for phase locking.

DESTINATION ADDRESS

The destination address indicates the destination of the packet on the network and is used to filter unwanted packets from reaching a node. There are three types of address formats supported by the NIC: physical, multicast, and broadcast. The physical address is a unique address that corresponds only to a single node. All physical addresses have an MSB of ''0''. These addresses are compared to the internally stored physical address registers. Each bit in the destination address must match in order for the NIC to accept the packet. Multicast addresses begin with an MSB of "1". The DP8390D filters multicast addresses using a standard hashing algorithm that maps all multicast addresses into a 6-bit value. This 6-bit value indexes a 64-bit array that filters the value. If the address consists of all 1's it is a broadcast address, indicating that the packet is intended for all nodes. A promiscuous mode allows reception of all packets: the destination address is not required to match any filters. Physical, broadcast, multicast, and promiscuous address modes can be selected.

SOURCE ADDRESS

The source address is the physical address of the node that sent the packet. Source addresses cannot be multicast or broadcast addresses. This field is simply passed to buffer memory.

LENGTH FIELD

The 2-byte length field indicates the number of bytes that are contained in the data field of the packet. This field is not interpreted by the NIC.

DATA FIELD

The data field consists of anywhere from 46 to 1500 bytes. Messages longer than 1500 bytes need to be broken into multiple packets. Messages shorter than 46 bytes will require appending a pad to bring the data field to the minimum length of 46 bytes. If the data field is padded, the number of valid data bytes is indicated in the length field. The NIC does not strip or append pad bytes for short packets, or check for oversize packets.

FCS FIELD

The Frame Check Sequence (FCS) is a 32-bit CRC field calculated and appended to a packet during transmission to allow detection of errors when a packet is received. During reception, error free packets result in a specific pattern in the CRC generator. Packets with improper CRC will be rejected. The AUTODIN II (X32 + X26 + X23 + X22 + X16 + $X^{12} + X^{11} + X^{10} + X^{8} + X^{7} + X^{5} + X^{4} + X^{2} + X^{1} + 1$ polynomial is used for the CRC calculations.

6.0 Direct Memory Access Control (DMA)

The DMA capabilities of the NIC greatly simplify use of the DP8390D in typical configurations. The local DMA channel transfers data between the FIFO and memory. On transmission, the packet is DMA'd from memory to the FIFO in bursts. Should a collision occur (up to 15 times), the packet is retransmitted with no processor intervention. On reception, packets are DMAed from the FIFO to the receive buffer ring (as explained below).

A remote DMA channel is also provided on the NIC to accomplish transfers between a buffer memory and system memory. The two DMA channels can alternatively be combined to form a single 32-bit address with 8- or 16-bit data.

DUAL DMA CONFIGURATION

An example configuration using both the local and remote DMA channels is shown below. Network activity is isolated

on a local bus, where the NIC's local DMA channel performs burst transfers between the buffer memory and the NIC's FIFO. The Remote DMA transfers data between the buffer memory and the host memory via a bidirectional I/O port. The Remote DMA provides local addressing capability and is used as a slave DMA by the host. Host side addressing must be provided by a host DMA or the CPU. The NIC allows Local and Remote DMA operations to be interleaved.

SINGLE CHANNEL DMA OPERATION

If desirable, the two DMA channels can be combined to provide a 32-bit DMA address. The upper 16 bits of the 32 bit address are static and are used to point to a 64k byte (or 32k word) page of memory where packets are to be received and transmitted.

7.0 Packet Reception (Continued)

for storing packets is controlled by Buffer Management Logic in the NIC. The Buffer Management Logic provides three basic functions: linking receive buffers for long packets, recovery of buffers when a packet is rejected, and recirculation of buffer pages that have been read by the host.

At initialization, a portion of the 64k byte (or 32k word) address space is reserved for the receive buffer ring. Two eight bit registers, the Page Start Address Register (PSTART) and the Page Stop Address Register (PSTOP) define the physical boundaries of where the buffers reside. The NIC treats the list of buffers as a logical ring; whenever the DMA address reaches the Page Stop Address, the DMA is reset to the Page Start Address.

INITIALIZATION OF THE BUFFER RING

Two static registers and two working registers control the operation of the Buffer Ring. These are the Page Start Register, Page Stop Register (both described previously), the Current Page Register and the Boundary Pointer Register. The Current Page Register points to the first buffer used to store a packet and is used to restore the DMA for writing status to the Buffer Ring or for restoring the DMA address in the event of a Runt packet, a CRC, or Frame Alignment error. The Boundary Register points to the first packet in the Ring not yet read by the host. If the local DMA address ever reaches the Boundary, reception is aborted. The Boundary Pointer is also used to initialize the Remote DMA for removing a packet and is advanced when a packet is removed. A simple analogy to remember the function of these registers is that the Current Page Register acts as a Write Pointer and the Boundary Pointer acts as a Read Pointer.

Note 1: At initialization, the Page Start Register value should be loaded into both the Current Page Register and the Boundary Pointer Register. Note 2: The Page Start Register must not be initialized to 00H.

Receive Buffer Ring At Initialization

BEGINNING OF RECEPTION

When the first packet begins arriving the NIC begins storing the packet at the location pointed to by the Current Page

Register. An offset of 4 bytes is saved in this first buffer to allow room for storing receive status corresponding to this packet.

LINKING RECEIVE BUFFER PAGES

If the length of the packet exhausts the first 256 byte buffer, the DMA performs a forward link to the next buffer to store the remainder of the packet. For a maximal length packet the buffer logic will link six buffers to store the entire packet. Buffers cannot be skipped when linking, a packet will always be stored in contiguous buffers. Before the next buffer can be linked, the Buffer Management Logic performs two comparisons. The first comparison tests for equality between the DMA address of the next buffer and the contents of the Page Stop Register. If the buffer address equals the Page Stop Register, the buffer management logic will restore the DMA to the first buffer in the Receive Buffer Ring value programmed in the Page Start Address Register. The second comparison tests for equality between the DMA address of the next buffer address and the contents of the Boundary Pointer Register. If the two values are equal the reception is aborted. The Boundary Pointer Register can be used to protect against overwriting any area in the receive buffer ring that has not yet been read. When linking buffers, buffer management will never cross this pointer, effectively avoiding any overwrites. If the buffer address does not match either the Boundary Pointer or Page Stop Address, the link to the next buffer is performed. Tabler Rino Theose are the Prace Start Register. The three Rino At Institute Teams are the Boundary Donton Feature Teams and the boundary Donton Feature Teams and the boundary content in the observer to the CFT C France St

Linking Buffers

Before the DMA can enter the next contiguous 256 byte buffer, the address is checked for equality to PSTOP and to the Boundary Pointer. If neither are reached, the DMA is allowed to use the next buffer.

Buffer Ring Overflow

If the Buffer Ring has been filled and the DMA reaches the Boundary Pointer Address, reception of the incoming packet will be aborted by the NIC. Thus, the packets previously received and still contained in the Ring will not be destroyed.

In a heavily loaded network environment the local DMA may be disabled, preventing the NIC from buffering packets from the network. To guarantee this will not happen, a software reset must be issued during all Receive Buffer Ring overflows (indicated by the OVW bit in the Interrupt Status Register). The following procedure is required to recover from a Receiver Buffer Ring Overflow.

If this routine is not adhered to, the NIC may act in an unpredictable manner. It should also be noted that it is not permissible to service an overflow interrupt by continuing to empty packets from the receive buffer without implementing the prescribed overflow routine. A flow chart of the NIC's overflow routine can be found at the right.

Note: It is necessary to define a variable in the driver, which will be called ''Resend''.

- 1. Read and store the value of the TXP bit in the NIC's Command Register.
- 2. Issue the STOP command to the NIC. This is accomplished be setting the STP bit in the NIC's Command Register. Writing 21H to the Command Register will stop the NIC.
- Note: If the STP is set when a transmission is in progress, the RST bit may not be set. In this case, the NIC is guaranteed to be reset after the longest packet time (1500 bytes = 1.2 ms). For the DP8390D (but not for the DP8390B), the NIC will be reset within 2 microseconds after the STP bit is set and Loopback mode 1 is programmed.
- 3. Wait for at least 1.6 ms. Since the NIC will complete any transmission or reception that is in progress, it is necessary to time out for the maximum possible duration of an Ethernet transmission or reception. By waiting 1.6 ms this is achieved with some guard band added. Previously, it was recommended that the RST bit of the Interrupt Status Register be polled to insure that the pending transmission or reception is completed. This bit is not a reliable indicator and subsequently should be ignored.
- 4. Clear the NIC's Remote Byte Count registers (RBCR0 and RBCR1).

If this value is a 0, set the ''Resend'' variable to a 0 and jump to step 6.

If this value is a 1, read the NIC's Interrupt Status Register. If either the Packet Transmitted bit (PTX) or Transmit Error bit (TXE) is set to a 1, set the ''Resend'' variable to a 0 and jump to step 6. If neither of these bits is set, place a 1 in the ''Resend'' variable and jump to step 6.

This step determines if there was a transmission in progress when the stop command was issued in step 2. If there was a transmission in progress, the NIC's ISR is read to determine whether or not the packet was recognized by the NIC. If neither the PTX nor TXE bit was set,

7.0 Packet Reception (Continued)

then the packet will essentially be lost and re-transmitted only after a time-out takes place in the upper level software. By determining that the packet was lost at the driver level, a transmit command can be reissued to the NIC once the overflow routine is completed (as in step 11). Also, it is possible for the NIC to defer indefinitely, when it is stopped on a busy network. Step 5 also alleviates this problem. Step 5 is essential and should not be omitted from the overflow routine, in order for the NIC to operate correctly.

- 6. Place the NIC in either mode 1 or mode 2 loopback. This can be accomplished by setting bits D2 and D1, of the Transmit Configuration Register, to ''0,1'' or ''1,0'', respectively.
- 7. Issue the START command to the NIC. This can be accomplished by writing 22H to the Command Register. This is necessary to activate the NIC's Remote DMA channel.
- 8. Remove one or more packets from the receive buffer ring.
- 9. Reset the overwrite warning (OVW, overflow) bit in the Interrupt Status Register.
- 10. Take the NIC out of loopback. This is done by writing the Transmit Configuration Register with the value it contains during normal operation. (Bits D2 and D1 should both be programmed to 0.)
- 11. If the ''Resend'' variable is set to a 1, reset the ''Resend'' variable and reissue the transmit command. This is done by writing a value of 26H to the Command Register. If the ''Resend'' variable is 0, nothing needs to be done.
- Note: If Remote DMA is not being used, the NIC does not need to be started before packets can be removed from the receive buffer ring. Hence, step 8 could be done before step 7.

END OF PACKET OPERATIONS

At the end of the packet the NIC determines whether the received packet is to be accepted or rejected. It either branches to a routine to store the Buffer Header or to another routine that recovers the buffers used to store the packet.

SUCCESSFUL RECEPTION

If the packet is successfully received as shown, the DMA is restored to the first buffer used to store the packet (pointed

to by the Current Page Register). The DMA then stores the Receive Status, a Pointer to where the next packet will be stored (Buffer 4) and the number of received bytes. Note that the remaining bytes in the last buffer are discarded and reception of the next packet begins on the next empty 256 byte buffer boundary. The Current Page Register is then initialized to the next available buffer in the Buffer Ring. (The location of the next buffer had been previously calculated and temporarily stored in an internal scratchpad register.)

BUFFER RECOVERY FOR REJECTED PACKETS

If the packet is a runt packet or contains CRC or Frame Alignment errors, it is rejected. The buffer management logic resets the DMA back to the first buffer page used to store the packet (pointed to by CURR), recovering all buffers that had been used to store the rejected packet. This operation will not be performed if the NIC is programmed to accept either runt packets or packets with CRC or Frame Alignment errors. The received CRC is always stored in buffer memory after the last byte of received data for the packet.

Termination of Received Packet-Packet Rejected

Error Recovery

If the packet is rejected as shown, the DMA is restored by the NIC by reprogramming the DMA starting address pointed to by the Current Page Register.

REMOVING PACKETS FROM THE RING

Packets are removed from the ring using the Remote DMA or an external device. When using the Remote DMA the Send Packet command can be used. This programs the Remote DMA to automatically remove the received packet pointed to by the Boundary Pointer. At the end of the transfer, the NIC moves the Boundary Pointer, freeing additional buffers for reception. The Boundary Pointer can also be moved manually by programming the Boundary Register. Care should be taken to keep the Boundary Pointer at least one buffer behind the Current Page Pointer.

The following is a suggested method for maintaining the Receive Buffer Ring pointers.

- 1. At initialization, set up a software variable (next_pkt) to indicate where the next packet will be read. At the beginning of each Remote Read DMA operation, the value of next_pkt will be loaded into RSAR0 and RSAR1.
- 2. When initializing the NIC set: $BNDRY = PSTART$ $CURR = PSTART + 1$ $next_pkt = PSTART + 1$

7.0 Packet Reception (Continued)

3. After a packet is DMAed from the Receive Buffer Ring, the Next Page Pointer (second byte in NIC buffer header) is used to update BNDRY and next_pkt.

 $next_pkt = Next$ Page Pointer $BNDRY = Next Page Pointer - 1$

If BNDRY \leq PSTART then BNDRY = PSTOP -1

Note the size of the Receive Buffer Ring is reduced by one 256-byte buffer; this will not, however, impede the operation of the NIC.

In StarLAN applications using bus clock frequencies greater than 4 MHz, the NIC does not update the buffer header information properly because of the disparity between the network and bus clock speeds. The lower byte count is copied twice into the third and fourth locations of the buffer header and the upper byte count is not written. The upper byte count, however, can be calculated from the current next page pointer (second byte in the buffer header) and the previous next page pointer (stored in memory by the CPU). The following routine calculates the upper byte count and allows StarLAN applications to be insensitive to bus clock speeds. Next_pkt is defined similarly as above.

upper byte count = $(PSTOP - next_pkt) +$

(next page pointer $-$ PSTART) -1

if (lower byte count) > 0 fch then

upper byte count = upper byte count $+1$

STORAGE FORMAT FOR RECEIVED PACKETS

The following diagrams describe the format for how received packets are placed into memory by the local DMA channel. These modes are selected in the Data Configuration Register.

This format used with Series 32000 808X type processors.

This format used with 68000 type processors.

Note: The Receive Byte Count ordering remains the same for $BOS = 0$ or 1.

 $BOS = 0$, WTS = 0 in Data Configuration Register. This format used with general 8-bit CPUs.

8.0 Packet Transmission

The Local DMA is also used during transmission of a packet. Three registers control the DMA transfer during transmission, a Transmit Page Start Address Register (TPSR) and the Transmit Byte Count Registers (TBCR0,1). When the NIC receives a command to transmit the packet pointed to by these registers, buffer memory data will be moved into the FIFO as required during transmission. The NIC will generate and append the preamble, synch and CRC fields.

TRANSMIT PACKET ASSEMBLY

The NIC requires a contiguous assembled packet with the format shown. The transmit byte count includes the Destination Address, Source Address, Length Field and Data. It does not include preamble and CRC. When transmitting data smaller than 46 bytes, the packet must be padded to a minimum size of 64 bytes. The programmer is responsible for adding and stripping pad bytes.

8.0 Packet Transmission (Continued)

TRANSMISSION

Prior to transmission, the TPSR (Transmit Page Start Register) and TBCR0, TBCR1 (Transmit Byte Count Registers) must be initialized. To initiate transmission of the packet the TXP bit in the Command Register is set. The Transmit Status Register (TSR) is cleared and the NIC begins to prefetch transmit data from memory (unless the NIC is currently receiving). If the interframe gap has timed out the NIC will begin transmission.

CONDITIONS REQUIRED TO BEGIN TRANSMISSION

In order to transmit a packet, the following three conditions must be met:

- 1. The Interframe Gap Timer has timed out the first 6.4 μ s of the Interframe Gap (See appendix for Interframe Gap Flowchart)
- 2. At least one byte has entered the FIFO. (This indicates that the burst transfer has been started)
- 3. If the NIC had collided, the backoff timer has expired.

In typical systems the NIC has already prefetched the first burst of bytes before the 6.4 μ s timer expires. The time during which NIC transmits preamble can also be used to load the FIFO.

Note: If carrier sense is asserted before a byte has been loaded into the FIFO, the NIC will become a receiver.

COLLISION RECOVERY

During transmission, the Buffer Management logic monitors the transmit circuitry to determine if a collision has occurred. If a collision is detected, the Buffer Management logic will reset the FIFO and restore the Transmit DMA pointers for retransmission of the packet. The COL bit will be set in the TSR and the NCR (Number of Collisions Register) will be incremented. If 15 retransmissions each result in a collision the transmission will be aborted and the ABT bit in the TSR will be set.

Note: NCR reads as zeroes if excessive collisions are encountered.

TRANSMIT PACKET ASSEMBLY FORMAT

The following diagrams describe the format for how packets must be assembled prior to transmission for different byte ordering schemes. The various formats are selected in the Data Configuration Register.

 $BOS = 0$. WTS = 1 in Data Configuration Register.

This format is used with Series 32000, 808X type processors.

 $BOS = 1$, WTS = 1 in Data Configuration Register. This format is used with 68000 type processors.

 $BOS = 0$. WTS = 0 in Data Configuration Register.

This format is used with general 8-bit CPUs.

- Note: All examples above will result in a transmission of a packet in order of DA0, DA1, DA2, DA3 . . . bits within each byte will be transmitted least significant bit first.
	- $DA =$ Destination Address
	- $SA =$ Source Address
	- $T/L = Type/L$ ength Field

9.0 Remote DMA

The Remote DMA channel is used to both assemble packets for transmission, and to remove received packets from the Receive Buffer Ring. It may also be used as a general purpose slave DMA channel for moving blocks of data or commands between host memory and local buffer memory. There are three modes of operation, Remote Write, Remote Read, or Send Packet.

Two register pairs are used to control the Remote DMA, a Remote Start Address (RSAR0, RSAR1) and a Remote Byte Count (RBCR0, RBCR1) register pair. The Start Address Register pair points to the beginning of the block to be moved while the Byte Count Register pair is used to indicate the number of bytes to be transferred. Full handshake logic is provided to move data between local buffer memory and a bidirectional I/O port.

9.0 Remote DMA (Continued)

REMOTE WRITE

A Remote Write transfer is used to move a block of data from the host into local buffer memory. The Remote DMA will read data from the I/O port and sequentially write it to local buffer memory beginning at the Remote Start Address. The DMA Address will be incremented and the Byte Counter will be decremented after each transfer. The DMA is terminated when the Remote Byte Count Register reaches a count of zero.

REMOTE READ

A Remote Read transfer is used to move a block of data from local buffer memory to the host. The Remote DMA will sequentially read data from the local buffer memory, beginning at the Remote Start Address, and write data to the I/O port. The DMA Address will be incremented and the Byte Counter will be decremented after each transfer. The DMA is terminated when the Remote Byte Count Register reaches zero.

REMOTE DMA WRITE

Setting PRQ Using the Remote Read

Under certain conditions the NIC's bus state machine may issue /MWR and /PRD before PRQ for the first DMA transfer of a Remote Write Command. If this occurs this could cause data corruption, or cause the remote DMA count to be different from the main CPU count causing the system to ''lock up''.

To prevent this condition when implementing a Remote DMA Write, the Remote DMA Write command should first be preceded by a Remote DMA Read command to insure that the PRQ signal is asserted before the NIC starts its port read cycle. The reason for this is that the state machine that asserts PRQ runs independently of the state machine that controls the DMA signals. The DMA machine assumes that PRQ is asserted, but actually may not be. To remedy this situation, a single Remote Read cycle should be inserted before the actual DMA Write Command is given. This will ensure that PRQ is asserted when the Remote DMA Write is subsequently executed. This single Remote Read cycle is

called a ''dummy Remote Read.'' In order for the dummy Remote Read cycle to operate correctly, the Start Address should be programmed to a known, safe location in the buffer memory space, and the Remote Byte Count should be progammed to a value greater than 1. This will ensure that the master read cycle is performed safely, eliminating the possiblity of data corruption.

Remote Write with High Speed Buses

When implementing the Remote DMA Write solution in previous section with high speed buses and CPU's, timing problems may cause the system to hang. Therefore additional considerations are required.

The problem occurs when the system can execute the dummy Remote Read and then start the Remote Write before the NIC has had a chance to execute the Remote Read. If this happens the PRQ signal will not get set, and the Remote Byte Count and Remote Start Address for the Remote Write operation could be corrupted. This is shown by the hatched waveforms in the timing diagram below. The execution of the Remote Read can be delayed by the local DMA operations (particularly during end-of-packet processing).

To ensure the dummy Remote Read does execute, a delay must be inserted between writing the Remote Read Command, and starting to write the Remote Write Start Address. (This time is designated in figure below by the delay arrows.) The recommended method to avoid this problem is, after the Remote Read command is given, to poll both bytes of the Current Remote DMA Address Registers. When the address has incremented, PRQ has been set. Software should recognize this and then start the Remote Write.

An additional caution for high speed systems is that the polling must follow guidelines specified at the end of Section 13. That is, there must be at least 4 bus clocks between chip selects. (For example, when BSCK $= 20$ MHz, then this time should be 200 ns.)

The general flow for executing a Remote Write is:

1. Set Remote Byte Count to a value \geq 1 and Remote Start Address to unused RAM (one location before the transmit start address is usually a safe location).

9.0 Remote DMA (Continued)

- 2. Issue the ''dummy'' Remote Read command.
- 3. Read the Current Remote DMA Address (CRDA) (both bytes).
- 4. Compare to previous CRDA value if different go to 6.
- 5. Delay and jump to 3.
- 6. Set up for the Remote Write command, by setting the Remote Byte Count and the Remote Start Address (note that if the Remote Byte count in step 1 can be set to the tramsmit byte count plus one, and the Remote Start Address to one less, these will now be incremented to the correct values.)
- 7. Issue the Remote Write command.

FIFO AND BUS OPERATIONS

Overview

To accommodate the different rates at which data comes from (or goes to) the network and goes to (or comes from) the system memory, the NIC contains a 16-byte FIFO for buffering data between the bus and the media. The FIFO threshold is programmable, allowing filling (or emptying) the FIFO at different rates. When the FIFO has filled to its programmed threshold, the local DMA channel transfers these bytes (or words) into local memory. It is crucial that the local DMA is given access to the bus within a minimum bus latency time; otherwise a FIFO underrun (or overrun) occurs.

To understand FIFO underruns or overruns, there are two causes which produce this condition-

- 1) the bus latency is so long that the FIFO has filled (or emptied) from the network before the local DMA has serviced the FIFO.
- 2) the bus latency or bus data rate has slowed the throughput of the local DMA to point where it is slower than the network data rate (10 Mb/s). This second condition is also dependent upon DMA clock and word width (byte wide or word wide).

The worst case condition ultimately limits the overall bus latency which the NIC can tolerate.

FIFO Underrun and Transmit Enable

During transmission, if a FIFO underrun occurs, the Transmit enable (TXE) output may remain high (active). Generally, this will cause a very large packet to be transmitted onto the network. The jabber feature of the transceiver will terminate the transmission, and reset TXE.

To prevent this problem, a properly designed system will not allow FIFO underruns by giving the NIC a bus acknowledge within time shown in the maximum bus latency curves shown and described later.

FIFO at the Beginning of Receive

At the beginning of reception, the NIC stores entire Address field of each incoming packet in the FIFO to determine whether the packet matches its Physical Address Registers or maps to one of its Multicast Registers. This causes the FIFO to accumulate 8 bytes. Furthermore, there are some synchronization delays in the DMA PLA. Thus, the actual time that BREQ is asserted from the time the Start of Frame Delimiter (SFD) is detected is 7.8 μ s. This operation affects the bus latencies at 2 and 4 byte thresholds during the first receive BREQ since the FIFO must be filled to 8 bytes (4 words) before issuing a BREQ.

FIFO Operation at the End of Receive

When Carrier Sense goes low, the NIC enters its end of packet processing sequence, emptying its FIFO and writing the status information at the beginning of the packet, figure below. This NIC holds onto the bus for the entire sequence. The longest time BREQ may be extended occurs when a packet ends just as the NIC performs its last FIFO burst. The NIC, in this case, performs a programmed burst transfer followed by flushing the remaining bytes in the FIFO, and completes by writing the header information to memory. The following steps occur during this sequence.

- 1) NIC issues BREQ because the FIFO threshold has been reached.
- 2) During the burst, packet ends, resulting in BREQ extended.
- 3) NIC flushes remaining bytes from FIFO.
- 4) NIC performs internal processing to prepare for writing the header.
- 5) NIC writes 4-byte (2-word) header.
- 6) NIC deasserts BREQ.

End of Packet Processing (EOPP) times for 10 MHz and 20 MHz have been tabulated in the table below.

End of Packet Processing Times for Various FIFO Thresholds, Bus Clocks and Transfer Modes

Threshold Detection (Bus Latency)

To assure that no overwriting of data in the FIFO, the FIFO logic flags a FIFO overrun as the 13th byte is written into the FIFO, effectively shortening the FIFO to 13 bytes. The FIFO logic also operates differently in Byte Mode and in Word Mode. In Byte Mode, a threshold is indicated when the $n+1$

byte has entered the FIFO; thus, with an 8 byte threshold, the NIC issues Bus Request (BREQ) when the 9th byte has entered the FIFO. For Word Mode, BREQ is not generated until the $n+2$ bytes have entered the FIFO. Thus, with a 4 word threshold (equivalent to 8 byte threshold), BREQ is issued when the 10th byte has entered the FIFO. The two graphs, the figures above, indicate the maximum allowable bus latency for Word and Byte transfer modes.

The FIFO at the Beginning of Transmit

Before transmitting, the NIC performs a prefetch from memory to load the FIFO. The number of bytes prefetched is the programmed FIFO threshold. The next BREQ is not issued until after the NIC actually begins trasmitting data, i.e., after SFD. The Transmit Prefetch diagram illustrates this process.

SEND PACKET COMMAND

The Remote DMA channel can be automatically initialized to transfer a single packet from the Receive Buffer Ring.

The CPU begins this transfer by issuing a ''Send Packet'' Command. The DMA will be initialized to the value of the Boundary Pointer Register and the Remote Byte Count Register pair (RBCR0, RBCR1) will be initialized to the value of the Receive Byte Count fields found in the Buffer Header of each packet. After the data is transferred, the Boundary Pointer is advanced to allow the buffers to be used for new receive packets. The Remote Read will terminate when the Byte Count equals zero. The Remote DMA is then prepared to read the next packet from the Receive Buffer Ring. If the DMA pointer crosses the Page Stop Register, it is reset to the Page Start Address. This allows the Remote DMA to remove packets that have wrapped around to the top of the Receive Buffer Ring.

Note 1: In order for the NIC to correctly execute the Send Packet Command, the upper Remote Byte Count Register (RBCR1) must first be loaded with 0FH.

Note 2: The Send Packet command cannot be used with 68000 type processors.

Page 2 Address Assignments (PS1 = $1,$ PS0 = 0)

Note: Page 2 registers should only be accessed for diagnostic purposes. They should not be modified during normal operation.

Page 3 should never be modified.

10.3 Register Descriptions

COMMAND REGISTER (CR) 00H (READ/WRITE)

The Command Register is used to initiate transmissions, enable or disable Remote DMA operations and to select register pages. To issue a command the microprocessor sets the corresponding bit(s) (RD2, RD1, RD0, TXP). Further commands may be overlapped, but with the following rules: (1) If a transmit command overlaps with a remote DMA operation, bits RD0, RD1, and RD2 must be maintained for the remote DMA command when setting the TXP bit. Note, if a remote DMA command is re-issued when giving the transmit command, the DMA will complete immediately if the remote byte count register have not been reinitialized. (2) If a remote DMA operation overlaps a transmission, RD0, RD1, and RD2 may be written with the desired values and a ''0'' written to the TXP bit. Writing a ''0'' to this bit has no effect. (3) A remote write DMA may not overlap remote read operation or visa versa. Either of these operations must either complete or be aborted before the other operation may start. Bits PS1, PS0, RD2, and STP may be set any time.

10.3 Register Descriptions (Continued)

INTERRUPT STATUS REGISTER (ISR) 07H (READ/WRITE)

This register is accessed by the host processor to determine the cause of an interrupt. Any interrupt can be masked in the Interrupt Mask Register (IMR). Individual interrupt bits are cleared by writing a ''1'' into the corresponding bit of the ISR. The INT signal is active as long as any unmasked signal is set, and will not go low until all unmasked bits in this register have been cleared. The ISR must be cleared after power up by writing it with all 1's.

completed. D7 RST RESET STATUS: Set when NIC enters reset state and cleared when a Start Command is issued to the CR. This bit is also set when a Receive Buffer Ring overflow occurs and is cleared when one or more packets have been removed from the ring. Writing to this bit has no effect. NOTE: This bit does not generate an interrupt, it is merely a status indicator.

10.3 Register Descriptions (Continued)

INTERRUPT MASK REGISTER (IMR) 0FH (WRITE)

The Interrupt Mask Register is used to mask interrupts. Each interrupt mask bit corresponds to a bit in the Interrupt Status Register (ISR). If an interrupt mask bit is set an interrupt will be issued whenever the corresponding bit in the ISR is set. If any bit in the IMR is set low, an interrupt will not occur when the bit in the ISR is set. The IMR powers up all zeroes.

10.3 Register Descriptions (Continued)

TRANSMIT CONFIGURATION REGISTER (TCR) 0DH (WRITE)

The transmit configuration establishes the actions of the transmitter section of the NIC during transmission of a packet on the network. LB1 and LB0 which select loopback mode power up as 0.

10.3 Register Descriptions (Continued)

TRANSMIT STATUS REGISTER (TSR) 04H (READ)

This register records events that occur on the media during transmission of a packet. It is cleared when the next transmission is initiated by the host. All bits remain low unless the event that corresponds to a particular bit occurs during transmission. Each transmission should be followed by a read of this register. The contents of this register are not specified until after the first transmission.

10.3 Register Descriptions (Continued)

RECEIVE STATUS REGISTER (RSR) 0CH (READ)

This register records status of the received packet, including information on errors and the type of address match, either physical or multicast. The contents of this register are written to buffer memory by the DMA after reception of a good packet. If packets with errors are to be saved the receive status is written to memory at the head of the erroneous packet if an erroneous packet is received. If packets with errors are to be rejected the RSR will not be written to memory. The contents will be cleared when the next packet arrives. CRC errors, Frame Alignment errors and missed packets are counted internally by the NIC which relinquishes the Host from reading the RSR in real time to record errors for Network Management Functions. The contents of this register are not specified until after the first reception.

Note: Following coding applies to CRC and FAE bits

FAE CRC Type of Error

0 0 No Error (Good CRC and <6 Dribble Bits)
0 1 CRC Error

- 1 CRC Error 1 0 Illegal, will not occur
- 1 1 Frame Alignment Error and CRC Error

The DMA Registers are partitioned into three groups; Transmit, Receive and Remote DMA Registers. The Transmit registers are used to initialize the Local DMA Channel for transmission of packets while the Receive Registers are used to initialize the Local DMA Channel for packet Reception. The Page Stop, Page Start, Current and Boundary Registers are used by the Buffer Management Logic to supervise the Receive Buffer Ring. The Remote DMA Registers are used to initialize the Remote DMA.

Note: In the figure above, registers are shown as 8 or 16 bits wide. Although some registers are 16-bit internal registers, all registers are accessed as 8-bit registers. Thus the 16-bit Transmit Byte Count Register is broken into two 8-bit registers, TBCR0 and TBCR1. Also TPSR, PSTART, PSTOP, CURR and BNRY only check or control the upper 8 bits of address information on the bus. Thus they are shifted to positions 15-8 in the diagram above.

10.5 TRANSMIT DMA REGISTERS

TRANSMIT PAGE START REGISTER (TPSR)

This register points to the assembled packet to be transmitted. Only the eight higher order addresses are specified since all transmit packets are assembled on 256-byte page boundaries. The bit assignment is shown below. The values placed in bits D7–D0 will be used to initialize the higher order address (A8–A15) of the Local DMA for transmission. The lower order bits (A7–A0) are initialized to zero.

(A7–A0 Initialized to zero)

TRANSMIT BYTE COUNT REGISTER 0,1 (TBCR0, TBCR1) These two registers indicate the length of the packet to be transmitted in bytes. The count must include the number of bytes in the source, destination, length and data fields. The maximum number of transmit bytes allowed is 64k bytes. The NIC will not truncate transmissions longer than 1500 bytes. The bit assignment is shown below:

10.6 LOCAL DMA RECEIVE REGISTERS

PAGE START STOP REGISTERS (PSTART, PSTOP)

The Page Start and Page Stop Registers program the starting and stopping address of the Receive Buffer Ring. Since the NIC uses fixed 256-byte buffers aligned on page boundaries only the upper eight bits of the start and stop address are specified.

PSTART,PSTOP bit assignment

BOUNDARY (BNRY) REGISTER

This register is used to prevent overflow of the Receive Buffer Ring. Buffer management compares the contents of this register to the next buffer address when linking buffers together. If the contents of this register match the next buffer address the Local DMA operation is aborted.

CURRENT PAGE REGISTER (CURR)

This register is used internally by the Buffer Management Logic as a backup register for reception. CURR contains the address of the first buffer to be used for a packet reception and is used to restore DMA pointers in the event of receive errors. This register is initialized to the same value as PSTART and should not be written to again unless the controller is Reset.

CURRENT LOCAL DMA REGISTER 0,1 (CLDA0,1)

These two registers can be accessed to determine the current Local DMA Address.

10.7 REMOTE DMA REGISTERS

REMOTE START ADDRESS REGISTERS (RSAR0,1)

Remote DMA operations are programmed via the Remote Start Address (RSAR0,1) and Remote Byte Count (RBCR0,1) registers. The Remote Start Address is used to point to the start of the block of data to be transferred and the Remote Byte Count is used to indicate the length of the block (in bytes).

10.8 PHYSICAL ADDRESS REGISTERS (PAR0–PAR5)

The physical address registers are used to compare the destination address of incoming packets for rejecting or accepting packets. Comparisons are performed on a bytewide basis. The bit assignment shown below relates the sequence in PAR0–PAR5 to the bit sequence of the received packet.

Note:

 $P/S =$ Preamble, Synch

DA0 = Physical/Multicast Bit

10.9 MULTICAST ADDRESS REGISTERS (MAR0–MAR7)

The multicast address registers provide filtering of multicast addresses hashed by the CRC logic. All destination addresses are fed through the CRC logic and as the last bit of the destination address enters the CRC, the 6 most significant bits of the CRC generator are latched. These 6 bits are then decoded by a 1 of 64 decode to index a unique filter bit (FB0–63) in the multicast address registers. If the filter bit selected is set, the multicast packet is accepted. The system designer would use a program to determine which filter bits to set in the multicast registers. All multicast filter bits that correspond to multicast address accepted by the node are then set to one. To accept all multicast packets all of the registers are set to all ones.

Note: Although the hashing algorithm does not guarantee perfect filtering of multicast address, it will perfectly filter up to 64 multicast addresses if these addresses are chosen to map into unique locations in the multicast filter.

If address Y is found to hash to the value 32 (20H), then FB32 in MAR4 should be initialized to ''1''. This will cause the NIC to accept any multicast packet with the address Y.

NETWORK TALLY COUNTERS

Three 8-bit counters are provided for monitoring the number of CRC errors, Frame Alignment Errors and Missed Packets. The maximum count reached by any counter is 192 (C0H). These registers will be cleared when read by the CPU. The count is recorded in binary in CT0–CT7 of each Tally Register.

Frame Alignment Error Tally (CNTR0)

This counter is incremented every time a packet is received with a Frame Alignment Error. The packet must have been recognized by the address recognition logic. The counter is cleared after it is read by the processor.

CRC Error Tally (CNTR1)

This counter is incremented every time a packet is received with a CRC error. The packet must first be recognized by the address recognition logic. The counter is cleared after it is read by the processor.

Frames Lost Tally Register (CNTR2)

This counter is incremented if a packet cannot be received due to lack of buffer resources. In monitor mode, this counter will count the number of packets that pass the address recognition logic.

FIFO

This is an eight bit register that allows the CPU to examine the contents of the FIFO after loopback. The FIFO will contain the last 8 data bytes transmitted in the loopback packet. Sequential reads from the FIFO will advance a pointer in the FIFO and allow reading of all 8 bytes.

Note: The FIFO should only be read when the NIC has been programmed in loopback mode.

NUMBER OF COLLISIONS (NCR)

This register contains the number of collisions a node experiences when attempting to transmit a packet. If no collisions are experienced during a transmission attempt, the COL bit of the TSR will not be set and the contents of NCR will be zero. If there are excessive collisions, the ABT bit in the TSR will be set and the contents of NCR will be zero. The NCR is cleared after the TXP bit in the CR is set.

11.0 Initialization Procedures

The NIC must be initialized prior to transmission or reception of packets from the network. Power on reset is applied to the NIC's reset pin. This clears/sets the following bits:

The NIC remains in its reset state until a Start Command is issued. This guarantees that no packets are transmitted or received and that the NIC remains a bus slave until all appropriate internal registers have been programmed. After initialization the STP bit of the command register is reset and packets may be received and transmitted.

Initialization Sequence

The following initialization procedure is mandatory.

- 1) Program Command Register for Page 0 (Command $Register = 21H$)
- 2) Initialize Data Configuration Register (DCR)
- 3) Clear Remote Byte Count Registers (RBCR0, RBCR1)
- 4) Initialize Receive Configuration Register (RCR)
- 5) Place the NIC in LOOPBACK mode 1 or 2 (Transmit Configuration Register $=$ 02H or 04H)
- 6) Initialize Receive Buffer Ring: Boundary Pointer (BNDRY), Page Start (PSTART), and Page Stop (PSTOP)
- 7) Clear Interrupt Status Register (ISR) by writing 0FFh to it.
- 8) Initialize Interrupt Mask Register (IMR)
- 9) Program Command Register for page 1 (Command $Request = 61H$ i)Initialize Physical Address Registers (PAR0-PAR5) ii)Initialize Multicast Address Registers (MAR0-MAR7) iii)Initialize CURRent pointer
- 10) Put NIC in START mode (Command Register $= 22H$). The local receive DMA is still not active since the NIC is in LOOPBACK.
- 11) Initialize the Transmit Configuration for the intended value. The NIC is now ready for transmission and reception.

11.0 Initialization Procedures

(Continued)

Before receiving packets, the user must specify the location of the Receive Buffer Ring. This is programmed in the Page Start and Page Stop Registers. In addition, the Boundary and Current Page Registers must be initialized to the value of the Page Start Register. These registers will be modified during reception of packets.

12.0 Loopback Diagnostics

Three forms of local loopback are provided on the NIC. The user has the ability to loopback through the deserializer on the DP8390D NIC, through the DP8391 SNI, and to the coax to check the link through the transceiver circuitry. Because of the half duplex architecture of the NIC, loopback testing is a special mode of operation with the following restrictions:

Restrictions During Loopback

The FIFO is split into two halves, one used for transmission the other for reception. Only 8-bit fields can be fetched from memory so two tests are required for 16-bit systems to verify integrity of the entire data path. During loopback the maximum latency from the assertion of BREQ to BACK is 2.0 μ s. Systems that wish to use the loopback test yet do not meet this latency can limit the loopback packet to 7 bytes without experiencing underflow. Only the last 8 bytes of the loopback packet are retained in the FIFO. The last 8 bytes can be read through the FIFO register which will advance through the FIFO to allow reading the receive packet sequentially.

When in word-wide mode with Byte Order Select set, the loopback packet must be assembled in the even byte locations as shown below. (The loopback only operates with byte wide transfers.)

When in word-wide mode with Byte Order Select low, the following format must be used for the loopback packet.

Note: When using loopback in word mode 2n bytes must be programmed in
TBCR0, 1. Where n = actual number of bytes assembled in even or odd location.

To initiate a loopback the user first assembles the loopback packet then selects the type of loopback using the Transmit Configuration register bits LB0, LB1. The transmit configuration register must also be set to enable or disable CRC generation during transmission. The user then issues a normal transmit command to send the packet. During loopback the receiver checks for an address match and if CRC bit in the TCR is set, the receiver will also check the CRC. The last 8 bytes of the loopback packet are buffered and can be read out of the FIFO using the FIFO read port.

Loopback Modes

MODE 1: Loopback Through the Controller (LB1 = 0 , LB0 $= 1$.

If the loopback is through the NIC then the serializer is simply linked to the deserializer and the receive clock is derived from the transmit clock.

MODE 2: Loopback Through the SNI (LB1 = 1, LB0 = 0). If the loopback is to be performed through the SNI, the NIC provides a control (LPBK) that forces the SNI to loopback all signals.

MODE 3: Loopback to Coax (LB1 = 1, LB0 = 1).

Packets can be transmitted to the coax in loopback mode to check all of the transmit and receive paths and the coax itself.

Note: In MODE 1, CRS and COL lines are not indicated in any status register, but the NIC will still defer if these lines are active. In MODE 2, COL is masked and in MODE 3 CRS and COL are not masked. It is not possible to go directly between the loopback modes, it is necessary to return to normal operation (00H) when changing modes.

Reading the Loopback Packet

The last eight bytes of a received packet can be examined by 8 consecutive reads of the FIFO register. The FIFO pointer is incremented after the rising edge of the CPU's read strobe by internally synchronizing and advancing the pointer. This may take up to four bus clock cycles, if the pointer has not been incremented by the time the CPU reads the FIFO register again, the NIC will insert wait states Note: The FIFO may only be read during Loopback. Reading the FIFO at any other time will cause the NIC to malfunction.

12.0 Loopback Diagnostics (Continued) Alignment of the Received Packet in the FIFO

Reception of the packet in the FIFO begins at location zero, after the FIFO pointer reaches the last location in the FIFO, the pointer wraps to the top of the FIFO overwriting the previously received data. This process continues until the last byte is received. The NIC then appends the received byte count in the next two locations of the FIFO. The contents of the Upper Byte Count are also copied to the next FIFO location. The number of bytes used in the loopback packet determines the alignment of the packet in the FIFO.

For the following alignment in the FIFO the packet length should be ($N \times 8$) + 5 Bytes. Note that if the CRC bit in the TCR is set, CRC will not be appended by the transmitter. If the CRC is appended by the transmitter, the last four bytes, bytes N-3 to N, correspond to the CRC.
FIFO

FIFO LOCATION CONTENTS 0 BYTE N-4 \rightarrow First Byte Read 1 BYTE N-3 (CRC1) AR Second Byte Read 2 BYTE N-2 (CRC2) # 3 BYTE N-1 (CRC3) 4 BYTE N (CRC4) 5 LOWER BYTE COUNT 6 UPPER BYTE COUNT \rightarrow Last Byte Read UPPER BYTE COUNT The Nicolay of Bond and the Control of th

LOOPBACK TESTS

Loopback capabilities are provided to allow certain tests to be performed to validate operation of the DP8390D NIC prior to transmitting and receiving packets on a live network. Typically these tests may be performed during power up of a node. The diagnostic provides support to verify the following:

- 1) Verify integrity of data path. Received data is checked against transmitted data.
- 2) Verify CRC logic's capability to generate good CRC on transmit, verify CRC on receive (good or bad CRC).
- 3) Verify that the Address Recognition Logic can a) Recognize address match packets
	- b) Reject packets that fail to match an address

LOOPBACK OPERATION IN THE NIC

Loopback is a modified form of transmission using only half of the FIFO. This places certain restrictions on the use of loopback testing. When loopback mode is selected in the TCR, the FIFO is split. A packet should be assembled in memory with programming of TPSR and TBCR0,TBCR1 registers. When the transmit command is issued the following operations occur:

Transmitter Actions

- 1) Data is transferred from memory by the DMA until the FIFO is filled. For each transfer TBCR0 and TBCR1 are decremented. (Subsequent burst transfers are initiated when the number of bytes in the FIFO drops below the programmed threshold.)
- 2) The NIC generates 56 bits of preamble followed by an 8-bit synch pattern.
- 3) Data transferred from FIFO to serializer.
- 4) If $CRC = 1$ in TCR, no CRC calculated by NIC, the last byte transmitted is the last byte from the FIFO (Allows software CRC to be appended). If $CRC = 0$, NIC calculates and appends four bytes of CRC.
- 5) At end of Transmission PTX bit set in ISR.

Receiver Actions

- 1) Wait for synch, all preamble stripped.
- 2) Store packet in FIFO, increment receive byte count for each incoming byte.
- 3) If CRC=0 in TCR, receiver checks incoming packet for CRC errors. If $CRC = 1$ in TCR, receiver does not check CRC errors, CRC error bit always set in RSR (for address matching packets).
- 4) At end of receive, receive byte count written into FIFO, receive status register is updated. The PRX bit is typically set in the RSR even if the address does not match. If CRC errors are forced, the packet must match the address filters in order for the CRC error bit in the RS to be set.

EXAMPLES

The following examples show what results can be expected from a properly operating NIC during loopback. The restrictions and results of each type of loopback are listed for reference. The loopback tests are divided into two sets of tests. One to verify the data path, CRC generation and byte count through all three paths. The second set of tests uses internal loopback to verify the receiver's CRC checking and address recognition. For all of the tests the DCR was programmed to 40h.

Note 1: Since carrier sense and collision detect inputs are blocked during internal loopback, carrier and CD heartbeat are not seen and the CRS and CDH bits are set.

Note 2: CRC errors are always indicated by receiver if CRC is appended by the transmitter.

Note 3: Only the PTX bit in the ISR is set, the PRX bit is only set if status is written to memory. In loopback this action does not occur and the PRX bit remains 0 for all loopback modes.

Note 4: All values are hex.

12.0 Loopback Diagnostics (Continued)

Note 1: CDH is set, CRS is not set since it is generated by the external encoder/decoder.

Note 1: CDH and CRS should not be set. The TSR however, could also contain 01H,03H,07H and a variety of other values depending on whether collisions were encountered or the packet was deferred.

Note 2.Will contain 08H if packet is not transmittable.

Note 3: During external loopback the NIC is now exposed to network traffic, it is therefore possible for the contents of both the Receive portion of the FIFO and the RSR to be corrupted by any other packet on the network. Thus in a live network the contents of the FIFO and RSR should not be depended on. The NIC will still abide by the standard CSMA/CD protocol in external loopback mode. (i.e. The network will not be disturbed by the loopback packet).

Note 4: All values are hex.

CRC AND ADDRESS RECOGNITION

The next three tests exercise the address recognition logic and CRC. These tests should be performed using internal loopback only so that the NIC is isolated from interference from the network. These tests also require the capability to generate CRC in software.

The address recognition logic cannot be directly tested. The CRC and FAE bits in the RSR are only set if the address of the packet matches the address filters. If errors are expected to be set and they are not set, the packet has been rejected on the basis of an address mismatch. The following sequence of packets will test the address recognition logic. The DCR should be set to 40H, the TCR should be set to 03H with a software generated CRC.

Note 1: Status will read 21H if multicast address used.

Note 2: Status will read 22H if multicast address used.

Note 3: In test A, the RSR is set up. In test B the address is found to match since the CRC is flagged as bad. Test C proves that the address recognition logic can distinguish a bad address and does not notify the RSR of the bad CRC. The receiving CRC is proven to work in test A and test B.

Note 4: All values are hex.

NETWORK MANAGEMENT FUNCTIONS

Network management capabilities are required for maintenance and planning of a local area network. The NIC supports the minimum requirement for network management in hardware, the remaining requirements can be met with software counts. There are three events that software alone can not track during reception of packets: CRC errors, Frame Alignment errors, and missed packets.

Since errored packets can be rejected, the status associated with these packets is lost unless the CPU can access the Receive Status Register before the next packet arrives. In situations where another packet arrives very quickly, the CPU may have no opportunity to do this. The NIC counts the number of packets with CRC errors and Frame Alignment errors. 8-bit counters have been selected to reduce overhead. The counters will generate interrupts whenever their MSBs are set so that a software routine can accumulate the network statistics and reset the counters before overflow occurs. The counters are sticky so that when they reach a count of 192 (C0H) counting is halted. An additional counter is provided to count the number of packets NIC misses due to buffer overflow or being offline. The contents have not a content of the signal of the content of the signal

The structure of the counters is shown below:

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Additional information required for network management is available in the Receive and Transmit Status Registers. Transmit status is available after each transmission for information regarding events during transmission.

Typically, the following statistics might be gathered in soft-

Traffic: Frames Sent OK Frames Received OK Multicast Frames Received Packets Lost Due to Lack of Resources Retries/Packet Errors: CRC Errors Alignment Errors

Excessive Collisions Packet with Length Errors Heartbeat Failure

ware:

13.0 Bus Arbitration and Timing (Continued)

REMOTE READ TIMING

- 1) The DMA reads byte/word from local buffer memory and writes byte/word into latch, increments the DMA address and decrements the byte count (RBCR0,1).
- 2) A Request Line (PRQ) is asserted to inform the system that a byte is available.
- 3) The system reads the port, the read strobe $(\overline{\mathsf{RACK}})$ is used as an acknowledge by the Remote DMA and it goes back to step 1.

Steps 1–3 are repeated until the remote DMA is complete.

Note that in order for the Remote DMA to transfer a byte from memory to the latch, it must arbitrate access to the local bus via a BREQ, BACK handshake. After each byte or word is transferred to the latch, BREQ is dropped. If a Local DMA is in progress, the Remote DMA is held off until the local DMA is complete.

REMOTE WRITE TIMING

A Remote Write operation transfers data from the I/O port to the local buffer RAM. The NIC initiates a transfer by requesting a byte/word via the PRQ. The system transfers a byte/word to the latch via IOW, this write strobe is detected by the NIC and PRQ is removed. By removing the PRQ, the Remote DMA holds off further transfers into the latch until the current byte/word has been transferred from the latch, PRQ is reasserted and the next transfer can begin.

1) NIC asserts PRQ. System writes byte/word into latch. NIC removes PRQ.

2) Remote DMA reads contents of port and writes byte/word to local buffer memory, increments address and decrements byte count (RBCR0,1).

3) Go back to step 1.

Steps 1–3 are repeated until the remote DMA is complete.

13.0 Bus Arbitration and Timing (Continued) SLAVE MODE TIMING

When $\overline{\text{CS}}$ is low, the NIC becomes a bus slave. The CPU can then read or write any internal registers. All register access is byte wide. The timing for register access is shown below. The host CPU accesses internal registers with four address lines, RA0–RA3, SRD and SWR strobes.

ADS0 is used to latch the address when interfacing to a multiplexed, address data bus. Since the NIC may be a local bus master when the host CPU attempts to read or write to the controller, an ACK line is used to hold off the CPU until the NIC leaves master mode. Some number of BSCK cycles is also required to allow the NIC to synchronize to the read or write cycle.

14.0 Preliminary Electrical Characteristics

Absolute Maximum Ratings

If Military/Aerospace specified devices are required, please contact the National Semiconductor Sales Office/Distributors for availability and specifications. Supply Voltage (V_{CC}) $-0.5V$ to $+7.0V$ DC Input Voltage (V_{IN}) -0.5 V to V_{CC} + 0.5V DC Output Voltage (V_{OUT}) -0.5 V to V_{CC} + 0.5V Storage Temperature Range (T_{STG}) $-65°C$ to $+150°C$

Preliminary DC Specifications $T_A = 0^\circ C$ to 70°C, $V_{CC} = 5V \pm 5$ %, unless otherwise specified

Note 1: These levels are tested dynamically using a limited amount of functional test patterns, please refer to AC Test Load.

Note 2: Limited functional test patterns are performed at these input levels. The majority of functional tests are performed at levels of 0V and 3V.

Note 3: This is measured with a 0.1 μ F bypass capacitor between V_{CC} and GND.

Note 4: The low drive CMOS compatible V_{OH} and V_{OL} limits are not tested directly. Detailed device characterization validates that this specification can be guaranteed by testing the high drive TTL compatible V_{OL} and V_{OH} specification.

Note 1: Cycles T1', T2', T3', T4' are only issued for the first transfer in a burst when 32-bit mode has been selected.

Note 2: The rate of bus clock must be high enough to support transfers to/from the FIFO at a rate greater than the serial network transfers from/to the FIFO. Note 3: These limits include the RC delay inherent in our test method. These signals typically turn off within 15 ns, enabling other devices to drive these lines with no contention.

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